

**Jack Spalding-Jamieson (Jack S-J)**  
**jacksj@uwaterloo.ca**

**Graph Drawing 2024**

## **Morphing Planar Graph Drawings via Orthogonal Box Drawings**

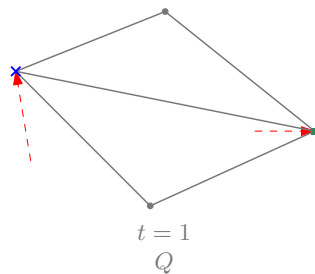
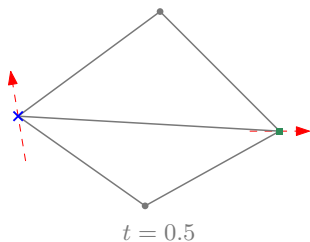
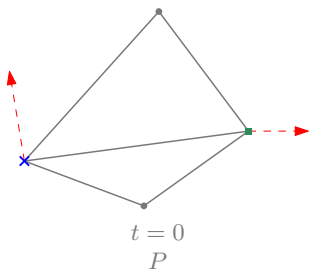
Joint work with Therese Biedl and Anna Lubiw

# Graph Morphing

Morph: Continuously transform between drawings (with time  $t \in [0, 1]$ )

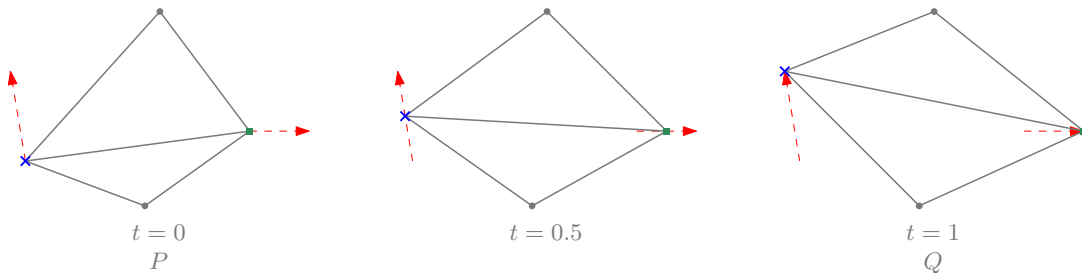
# Graph Morphing

Morph: Continuously transform between drawings (with time  $t \in [0, 1]$ )



# Graph Morphing

Morph: Continuously transform between drawings (with time  $t \in [0, 1]$ )



Linear morph: Linearly interpolate vertex (and other) locations

## Goals of morphing

Goal I: Make “nice” morphs.

- ❑ Simple paths of movement.
- ❑ Elementary steps.
- ❑ Few steps.

## Goals of morphing

Goal I: Make “nice” morphs.

- ❑ Simple paths of movement.
- ❑ Elementary steps.
- ❑ Few steps.

Goal II: Preserve drawing properties!

Planarity, few bend/straight line edges, orthogonality, drawing on a small grid, etc. Two types:

- ❑ Between elementary steps.
- ❑ During elementary steps/at all times.

## Goals of morphing

Goal I: Make “nice” morphs.

- ❑ Simple paths of movement.
- ❑ Elementary steps.
- ❑ Few steps.

Goal II: Preserve drawing properties!

Planarity, few bend/straight line edges, orthogonality, drawing on a small grid, etc. Two types:

- ❑ Between elementary steps.
- ❑ During elementary steps/at all times.

Goal III: Algorithmic properties

- ❑ Time complexity
- ❑ Computational Model

## Planarity-Preserving Morphs (preserved at all times)

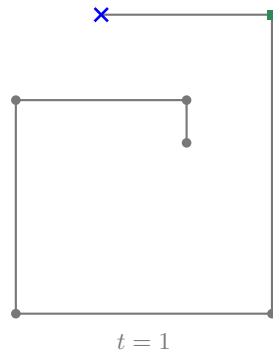
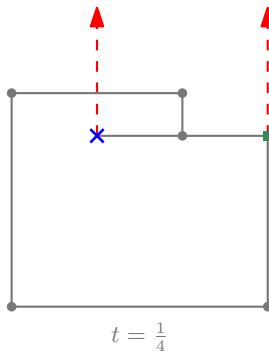
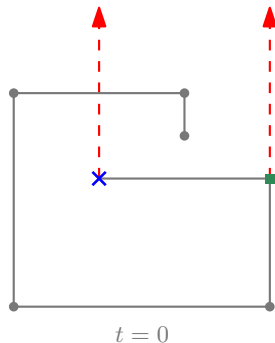
Planarity-Preserving Morph: At all times  $t$ , the “interpolated” drawing is also planar.



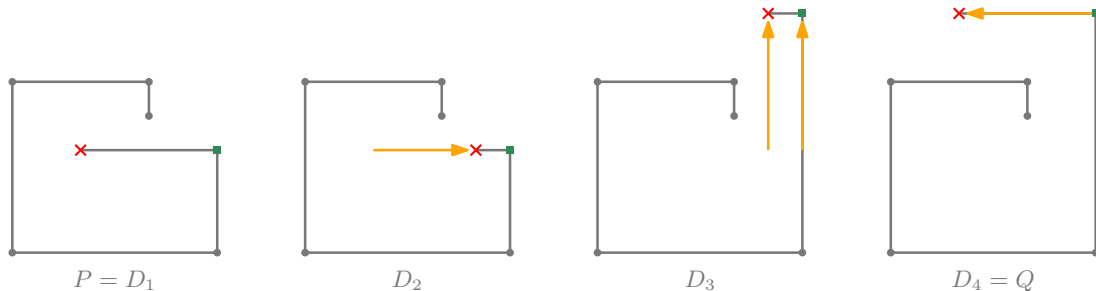
## Planarity-Preserving Morphs (preserved at all times)

Planarity-Preserving Morph: At all times  $t$ , the “interpolated” drawing is also planar.

Non-planarity-preserving morph:



## Linear Morphs Sequences



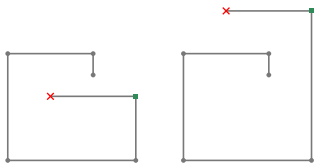
These are **explicit intermediate drawings**.

Entire morph represented by a finite sequence  $D_1, D_2, D_3, D_4$ .

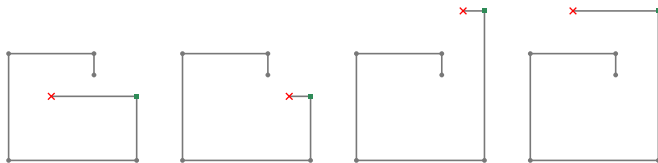
Can preserve extra properties on these drawings ("between steps").

# The Linear Morph Problem

Input:  
'Compatible' pair of drawings (labelled)

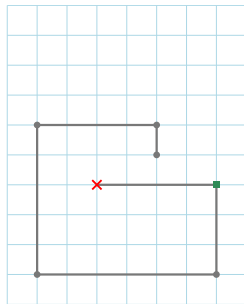
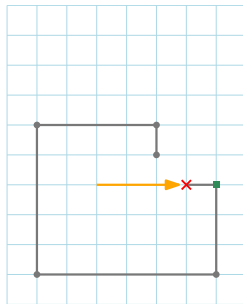
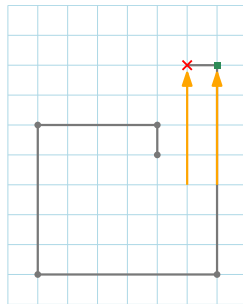
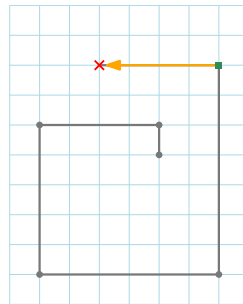


Output:  
Planarity-preserving linear morph sequence (list of drawings)



Objectives: Numerous!

## Linear Morph Sequences on a Grid


 $P = D_1$ 

 $D_2$ 

 $D_3$ 

 $D_4 = Q$ 

Explicit drawings are on a grid.

Implicit (interpolated) drawings are not.

## Open Problem in Morphing

### Open:

Input: Straight-line drawings each on an  $O(n) \times O(n)$  grid (same embedding of same graph).

Output: Linear morph sequence with properties:

- ❑ Explicit drawings: Polynomial-sized grid.
- ❑ Implicit+explicit drawings: Planar, straight-line edges
- ❑ Linear morph sequence length: Polynomial.

## Open Problem in Morphing

### Open:

Input: Straight-line drawings each on an  $O(n) \times O(n)$  grid (same embedding of same graph).

Output: Linear morph sequence with properties:

- ❖ Explicit drawings: Polynomial-sized grid.
- ❖ Implicit+explicit drawings: Planar, straight-line edges
- ❖ Linear morph sequence length: Polynomial.

**Conjecture:** Even stronger properties can be obtained.

Input: Straight-line drawings each on an  $O(n) \times O(n)$  grid (same embedding of same graph).

Output: Linear morph sequence with properties:

- ❖ Explicit drawings:  $O(n) \times O(n)$  grid.
- ❖ Implicit+explicit drawings: Planar, straight-line edges
- ❖ Linear morph sequence length:  $O(n)$

## Open Problem in Morphing

### Open:

Input: Straight-line drawings each on an  $O(n) \times O(n)$  grid (same embedding of same graph).

Output: Linear morph sequence with properties:

- ✦ Explicit drawings: Polynomial-sized grid.
- ✦ Implicit+explicit drawings: Planar, straight-line edges
- ✦ Linear morph sequence length: Polynomial.

**Conjecture:** Even stronger properties can be obtained.

Input: Straight-line drawings each on an  $O(n) \times O(n)$  grid (same embedding of same graph).

Output: Linear morph sequence with properties:

- ✦ Explicit drawings:  $O(n) \times O(n)$  grid.
- ✦ Implicit+explicit drawings: Planar, straight-line edges
- ✦ Linear morph sequence length:  $O(n)$

Various weakenings are known. We present a new one.

## Linear Morphs Sequences that Add/Remove Bends

Degenerate bend: Bend that “isn’t used” (coincident or  $180^\circ$  angle).

Equivalent drawings: Drawings that differ only by degenerate bends.





## Previous Results (Abridged)

Input: ('Compatible') pair of drawings

Output: Planarity-preserving linear morph sequence

	Graph/Drawing Class	Num linear morphs	Grid-size side-length	Bends per edge	Comput. Model	Time Complexity
Alamdari et al. (2017)	Straight-line	$O(n)$	Expo.	0	Powerful	$O(n^3)$
Klemz (2021)	Straight-line	$O(n)$	Expo.	0	Powerful	$O(n^2 \log n)$

## Previous Results (Abridged)

Input: ('Compatible') pair of drawings

Output: Planarity-preserving linear morph sequence

	Graph/Drawing Class	Num linear morphs	Grid-size side-length	Bends per edge	Comput. Model	Time Complexity
Alamdari et al. (2017)	Straight-line	$O(n)$	Expo.	0	Powerful	$O(n^3)$
Klemz (2021)	Straight-line	$O(n)$	Expo.	0	Powerful	$O(n^2 \log n)$
Klemz (2021)	2-connected	$O(n)$	Expo.	0	Powerful	$O(n^2)$

## Previous Results (Abridged)

Input: ('Compatible') pair of drawings

Output: Planarity-preserving linear morph sequence

	Graph/Drawing Class	Num linear morphs	Grid-size side-length	Bends per edge	Comput. Model	Time Complexity
Alamdari et al. (2017)	Straight-line	$O(n)$	Expo.	0	Powerful	$O(n^3)$
Klemz (2021)	Straight-line	$O(n)$	Expo.	0	Powerful	$O(n^2 \log n)$
Klemz (2021)	2-connected	$O(n)$	Expo.	0	Powerful	$O(n^2)$
Lubiw & Petrick (2011)	Straight-line	$O(n^6)$	$O(n^3)$	$O(n^5)$	Word RAM	Polynomial

## Previous Results (Abridged)

Input: ('Compatible') pair of drawings

Output: Planarity-preserving linear morph sequence

	Graph/Drawing Class	Num linear morphs	Grid-size side-length	Bends per edge	Comput. Model	Time Complexity
Alamdari et al. (2017)	Straight-line	$O(n)$	Expo.	0	Powerful	$O(n^3)$
Klemz (2021)	Straight-line	$O(n)$	Expo.	0	Powerful	$O(n^2 \log n)$
Klemz (2021)	2-connected	$O(n)$	Expo.	0	Powerful	$O(n^2)$
Lubiw & Petrick (2011)	Straight-line	$O(n^6)$	$O(n^3)$	$O(n^5)$	Word RAM	Polynomial
<b>This work (main result)</b>	Connected	$O(n)$	$O(n)$	$O(1)$	Word RAM	$O(n^2)$

## Previous Results (Abridged)

Input: ('Compatible') pair of drawings

Output: Planarity-preserving linear morph sequence

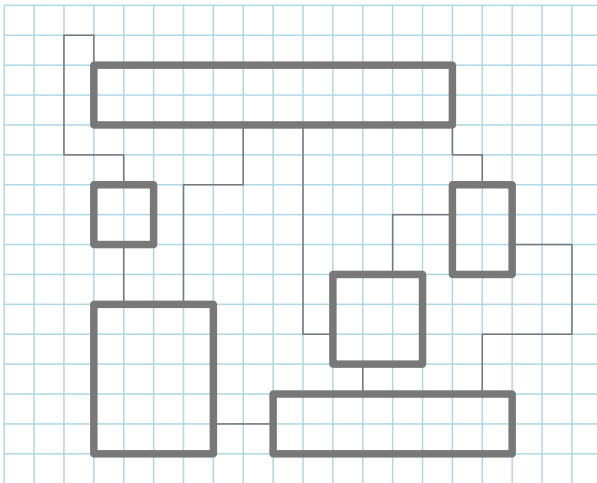
	Graph/Drawing Class	Num linear morphs	Grid-size side-length	Bends per edge	Comput. Model	Time Complexity
Alamdari et al. (2017)	Straight-line	$O(n)$	Expo.	0	Powerful	$O(n^3)$
Klemz (2021)	Straight-line	$O(n)$	Expo.	0	Powerful	$O(n^2 \log n)$
Klemz (2021)	2-connected	$O(n)$	Expo.	0	Powerful	$O(n^2)$
Lubiw & Petrick (2011)	Straight-line	$O(n^6)$	$O(n^3)$	$O(n^5)$	Word RAM	Polynomial
<b>This work (main result)</b>	Connected	$O(n)$	$O(n)$	$O(1)$	Word RAM	$O(n^2)$
Biedl et al. (2013)	Connected Orthogonal	$O(n^2)$	$O(n)$	$O(n)$	Word RAM	Polynomial
Van Goethem et al. (2022)	Orthogonal	$O(n)$	Polynomial	$O(1)$	Word RAM	Polynomial
<b>This work (main method)</b>	Connected Ortho-Box	$O(n)$	$O(n)$	$O(1)$	Word RAM	$O(n^2)$
<b>Open</b>	Many	Poly	Poly	0	Any	Any
<b>Lower Bounds</b>	Planar	$O(n)$	$O(n)$	0	Word RAM	$O(n^2)$

Grid size assumes input fits on the same grid.

Above table is not comprehensive.

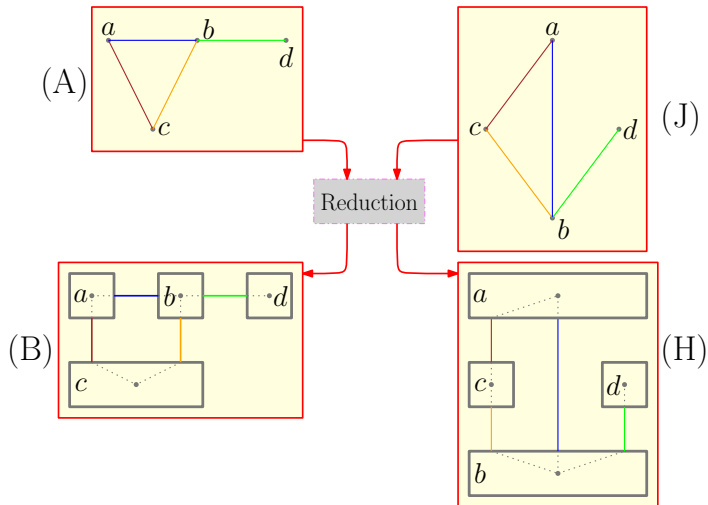
## High-Level Overview

First: Reduce problem to morphing **orthogonal box drawings**.



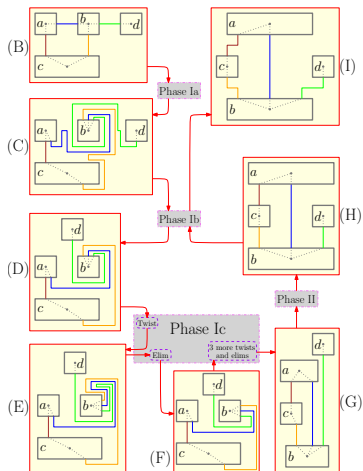
## High-Level Overview

First: Reduce problem to morphing “orthogonal box drawings”.



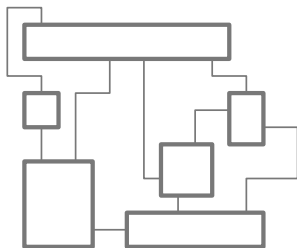
# High-Level Overview

Second: Morph orthogonal box drawings.

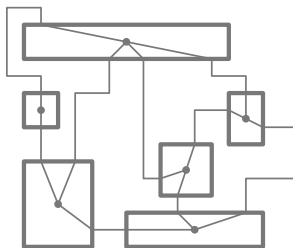




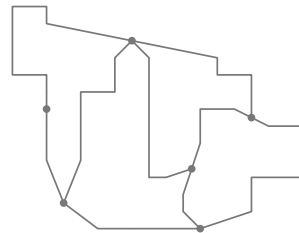
## Reduction: Admitted Drawings (1)



Orthogonal box drawing

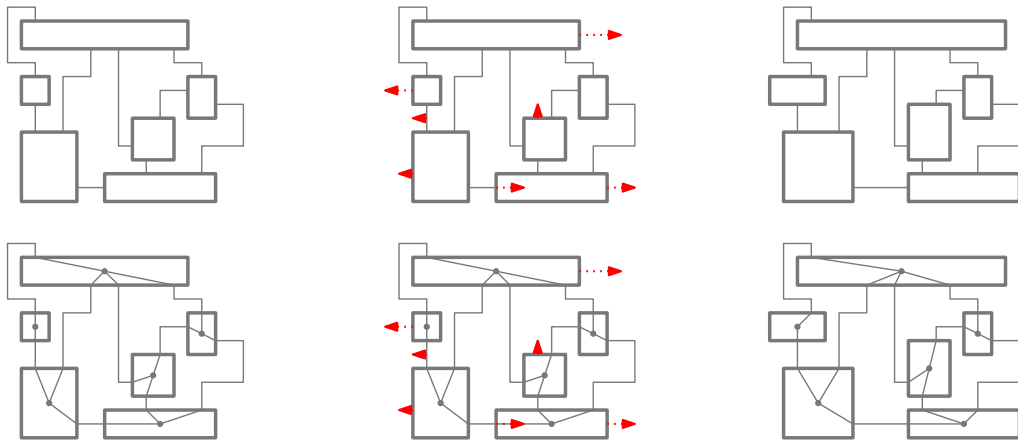


Both



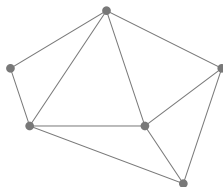
Admitted poly-line drawing

## Reduction: Admitted Drawings (2)

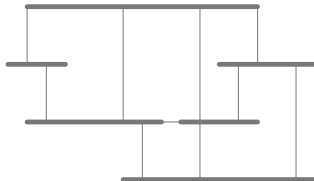


Morph of orthogonal box drawings  $\implies$  morph of admitted drawings

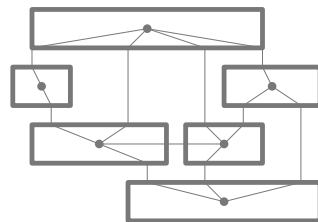
## Computing Box Drawings: Visibility Representations as an Intermediary



A planar straight-line drawing  $P$ .

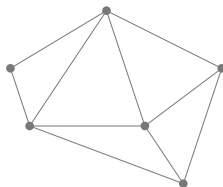


A visibility representation that can be computed from  $P$ .

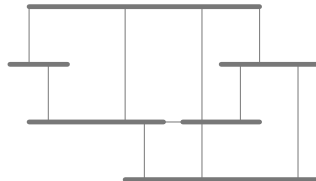


An orthogonal box drawing, and corresponding admitted drawing  $P'$ , which can both be computed from  $P$ .

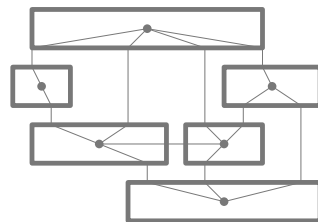
## Computing Box Drawings: Visibility Representations as an Intermediary



A planar straight-line drawing  $P$ .



A visibility representation that can be computed from  $P$ .



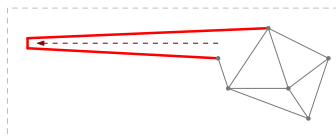
An orthogonal box drawing, and corresponding admitted drawing  $P'$ , which can both be computed from  $P$ .

How do we actually perform a morph?

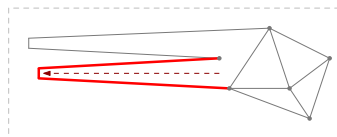
## Morphing from a straight-line to an admitted drawing: Method



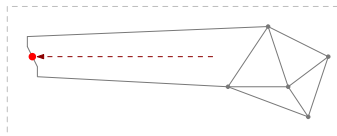
Step 0



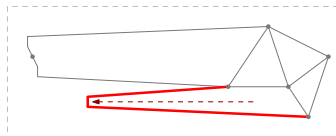
Step 1



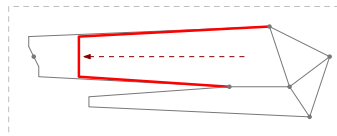
Step 2



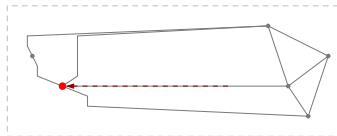
Step 3



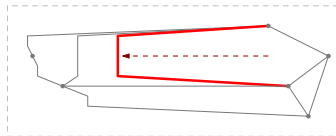
Step 4



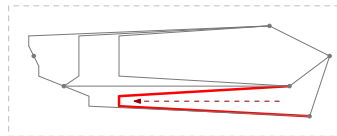
Step 5



Step 6

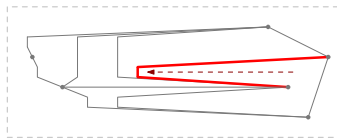


Step 7

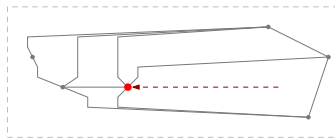


Step 8

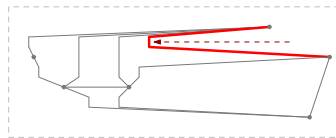
## Morphing from a straight-line to an admitted drawing: Method



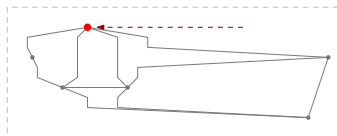
Step 9



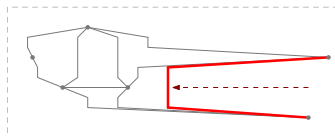
Step 10



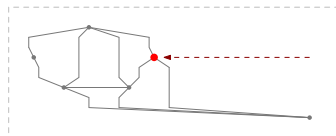
Step 11



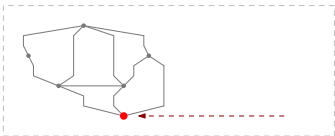
Step 12



Step 13



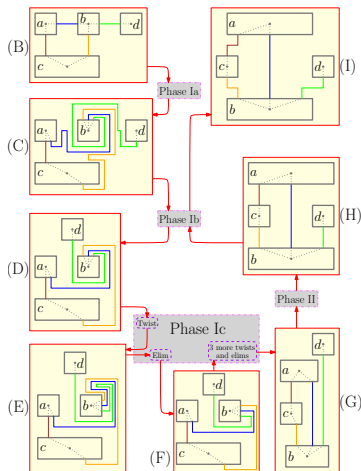
Step 14



Step 15

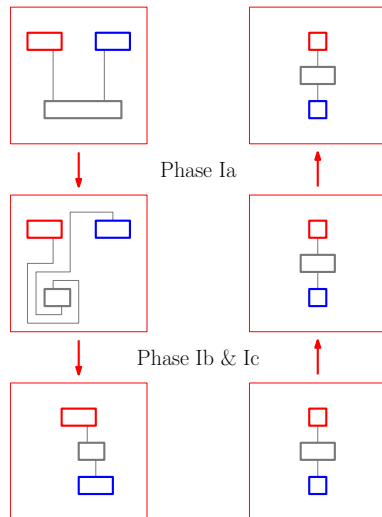
## Recall: High-Level Overview

Now have orthogonal box drawings, want to morph them.



# Phase I

Goal: Reduce to parallel box drawings.



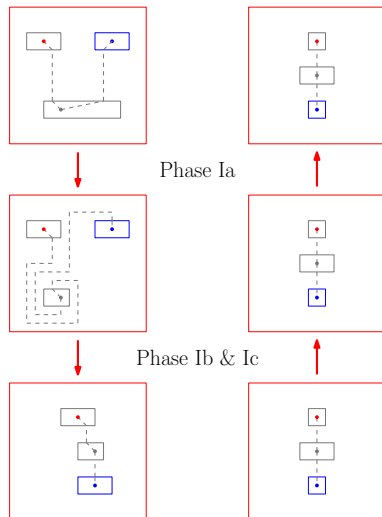


# Phase I

Goal: Reduce to parallel box drawings.

## Phase I Overview

- ❖ **Input:** Orthogonal box drawing pair
- ❖ **Output:** Parallel orthogonal box drawing pair (for each edge: same port locations, same sequence of turns)
- ❖ **Substeps:**
  - ❖ Phase Ia: Adjust port locations
  - ❖ Phase Ib: Initial zig-zag elimination
  - ❖ Phase Ic: Twists (plus more interspersed compaction/zig-zag elim)

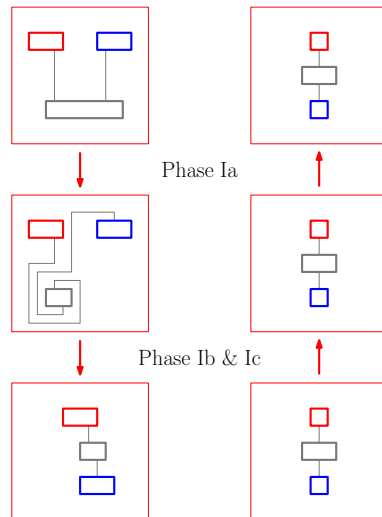


# Phase I

Goal: Reduce to parallel box drawings.

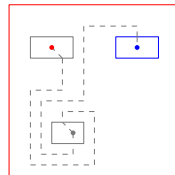
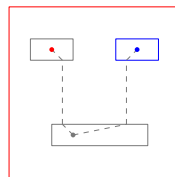
## Phase I Overview

- ❖ **Input:** Orthogonal box drawing pair
- ❖ **Output:** Parallel orthogonal box drawing pair (for each edge: same port locations, same sequence of turns)
- ❖ **Substeps:**
  - ❖ Phase Ia: Adjust port locations
  - ❖ Phase Ib: Initial zig-zag elimination
  - ❖ Phase Ic: Twists (plus more interspersed compaction/zig-zag elim)

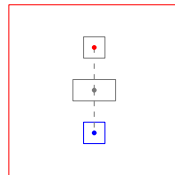
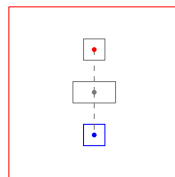


## Phase Ia

High-level: Move ports. Add bends to do so.



Phase Ia

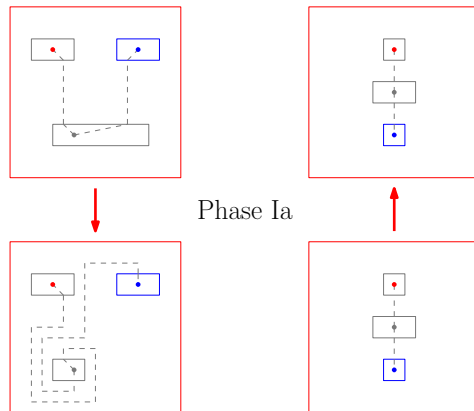


# Phase Ia

High-level: Move ports. Add bends to do so.

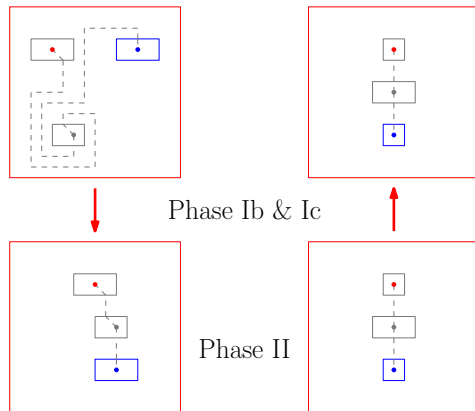
## Phase Ia Overview

- ❖ **Input:** Orthogonal box drawing pair
- ❖ **Output:** Port-aligned orthogonal box drawing pair (same relative port locations)



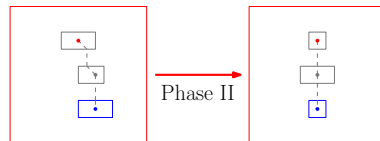
## Phase Ib & 1c

Get rid of all (extra) bends.

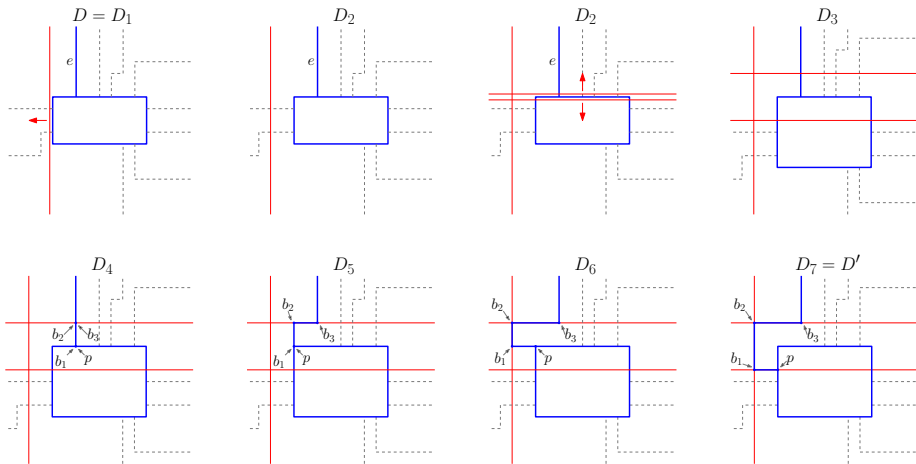


## Phase II

High-level: Use black-box result to morph parallel orthogonal box drawings (i.e., adjust lengths).



# Moving Ports around Corners

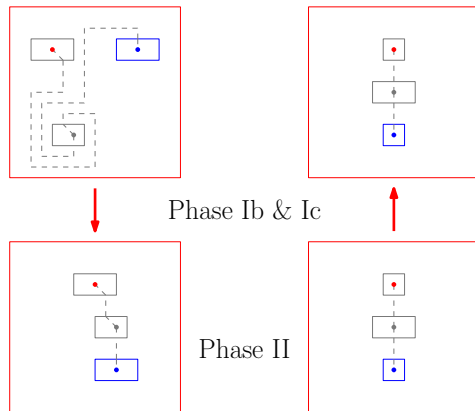


## Phase Ib & 1c

Get rid of all (extra) bends.

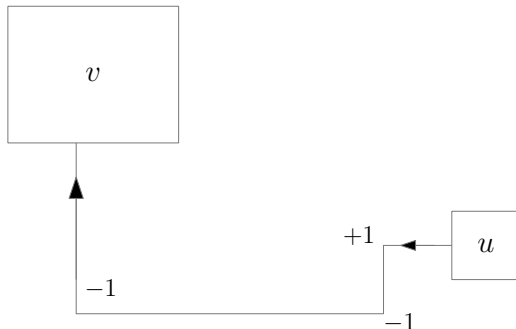
### Phase Ib Overview

- **Input:** Port-aligned orthogonal box drawing pair
- **Output:** Parallel orthogonal box drawing pair



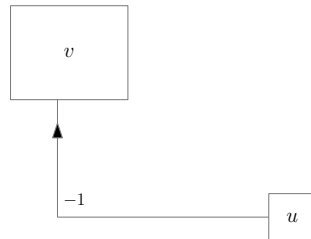
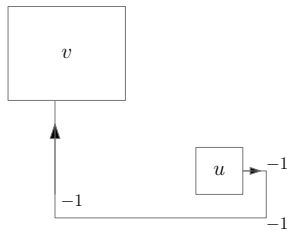


## Spirality



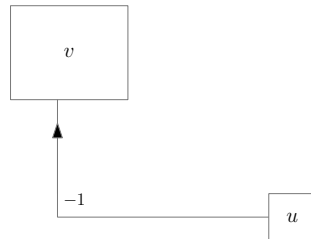
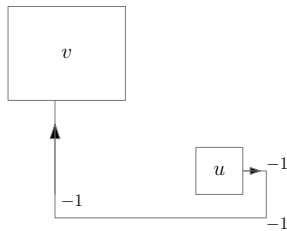
Spirality of the edge  $uv$  (oriented  $u$  to  $v$ ):  $-1$ .

## Difference in Spirality (1)



Difference in spirality of the edge  $uv$  (oriented  $u$  to  $v$ ):  $-2$ .

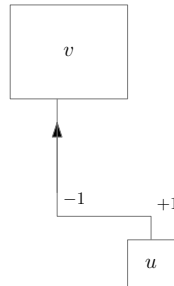
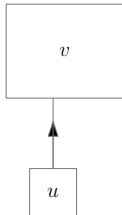
## Difference in Spirality (1)



Difference in spirality of the edge  $uv$  (oriented  $u$  to  $v$ ):  $-2$ .

Goal: Reduce this to zero.

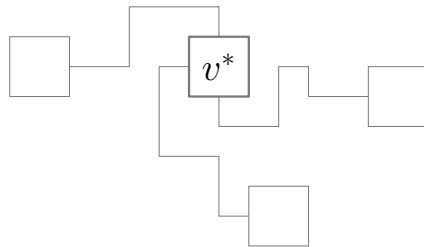
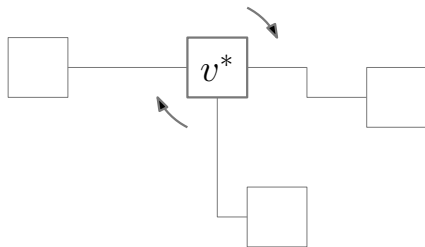
## Difference in Spirality (2)



Difference in spirality of the edge  $uv$  (oriented  $u$  to  $v$ ): 0.

Goal: Reduce this to zero.

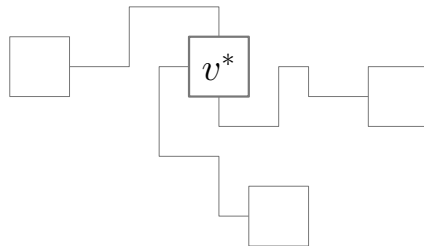
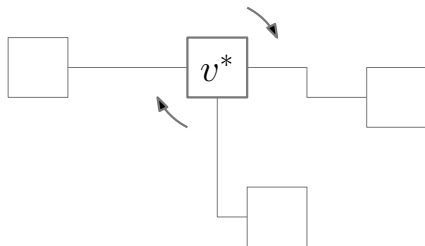
## Twists (High-Level)



Spirality changes! Net turns are added.

Similar to a result by Biedl et al.: Exists some number/direction of twists for each vertex so that difference in spirality becomes zero everywhere. This number is  $O(n)$  for each vertex.

## Twists (High-Level)



Spirality changes! Net turns are added.

Similar to a result by Biedl et al.: Exists some number/direction of twists for each vertex so that difference in spirality becomes zero everywhere. This number is  $O(n)$  for each vertex.

**Key difference/contribution:** We use simultaneous twists, so only  $O(n)$  operations needed.

## Twists (Implementation)

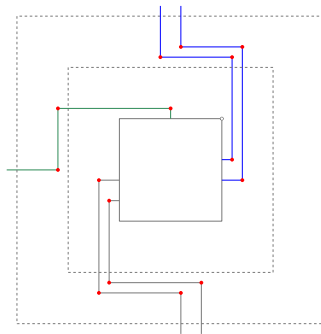
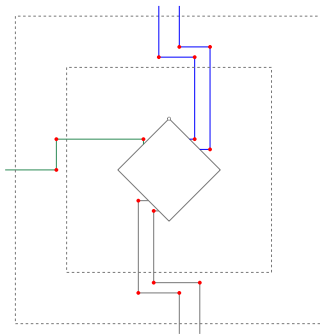
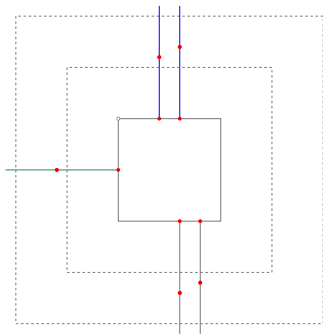
Two steps:

- ❑ “Prepare” drawing (make boxes square, well-spaced out)
- ❑ Twist everything simultaneously

## Twists (Implementation)

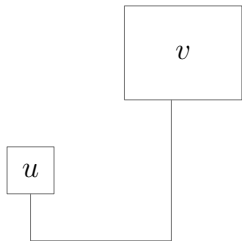
Two steps:

- ❑ “Prepare” drawing (make boxes square, well-spaced out)
- ❑ Twist everything simultaneously



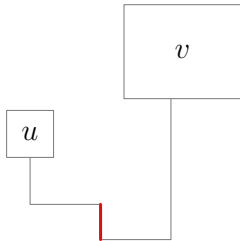


## Simplification/Canonical form: Zig-Zags

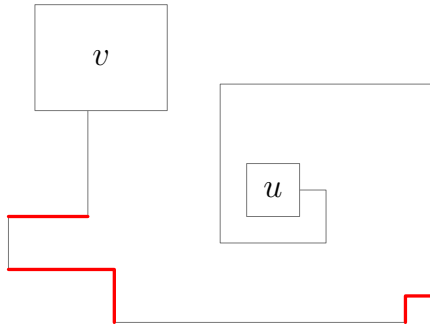


No zig-zags.

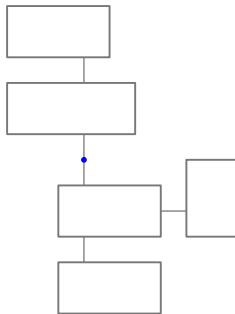
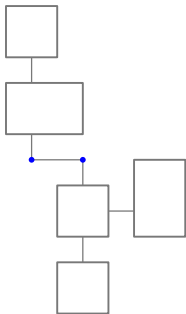
We want to remove zig-zags.

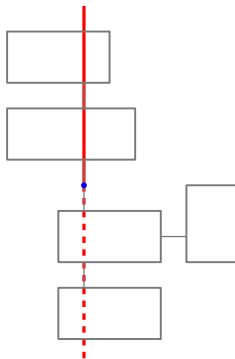
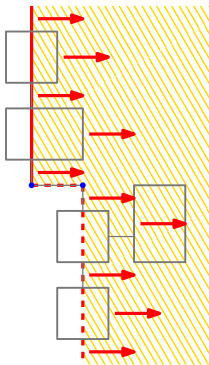


One (vertical) zig-zag.

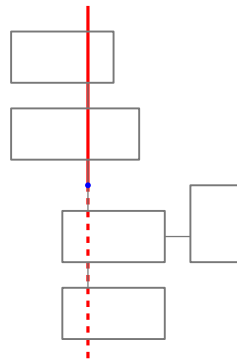
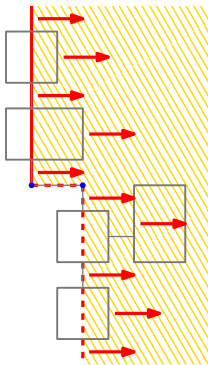


Five zig-zags, three horizontal and two vertical.





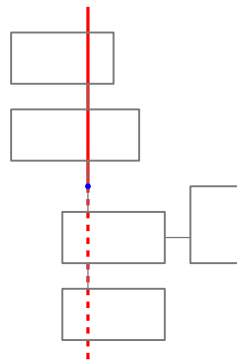
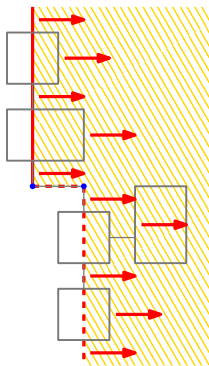
## Simplification/Canonical form: Removing a Single Zig-Zag (2)



Push each thing over if it lies to the right of the divider.

Problem: Requires a morph for each zig-zag (want  $O(1)$  morphs for all zig-zags).

## Simplification/Canonical form: Removing a Single Zig-Zag (2)



Push each thing over if it lies to the right of the divider.

Problem: Requires a morph for each zig-zag (want  $O(1)$  morphs for all zig-zags).

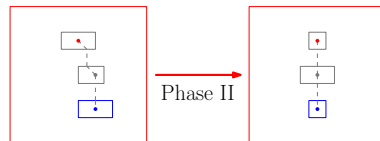
**Solution/new contribution:**  $O(1)$  morphs suffice, even on a grid (skipping details).

## Phase II

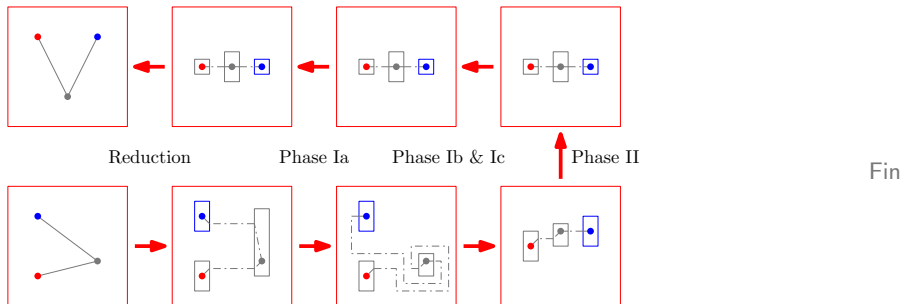
High-level: Use black-box result to morph parallel orthogonal box drawings (i.e., adjust lengths).

### Phase II Overview

- ❖ **Input:** Parallel orthogonal box drawing pair.
- ❖ **Output:** Linear morph sequence.
- ❖ **Methodology:** Appeal to black-box result by Biedl et al.. It requires connectivity.
  - ❖ Essentially, add edges to both drawings (and simplify again) until every face is a rectangle.

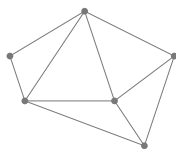
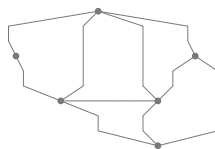


	Graph/Drawing Class	Num linear morphs	Grid-size side-length	Bends per edge	Comput. Model	Time Complexity
Main result	Connected	$O(n)$	$O(n)$	$O(1)$	Word RAM	$O(n^2)$
This work (main method)	Connected Ortho-Box	$O(n)$	$O(n)$	$O(1)$	Word RAM	$O(n^2)$
Open	Many	Poly	Poly	0	Any	Any
Lower Bounds	Planar	$O(n)$	$O(n)$	0	Word RAM	$O(n^2)$

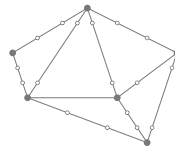


## Morphing from a straight-line to an admitted drawing: Brainstorming (1)

Have: a planar straight-line drawing  $P$ , an orthogonal box drawing  $D$  with an admitted drawing  $P'$ .  
Want: Morph from  $P$  to  $P'$ . Bends need to be added.

 $P$  $P'$ 

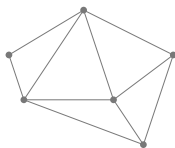
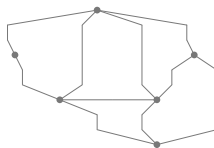
- ❖ Idea 1: Use same  $y$ -coordinate
- ❖ Problem: Not integer coordinates





## Morphing from a straight-line to an admitted drawing: Brainstorming (2)

Have: a planar straight-line drawing  $P$ , an orthogonal box drawing  $D$  with an admitted drawing  $P'$ .  
Want: Morph from  $P$  to  $P'$ . Bends need to be added.

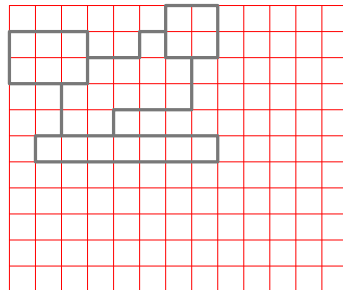
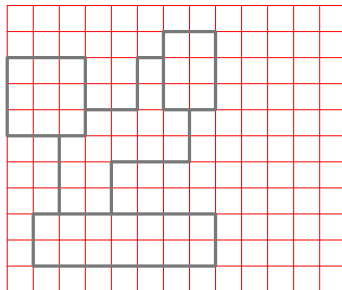
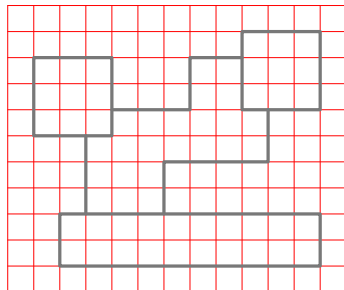
 $P$  $P'$ 

- ❖ ~~Idea 1: Use same  $y$ -coordinate~~
- ❖ Idea 2: Make them coincident with the vertex
- ❖ Possible problem: Not a unidirectional morph (complicated movement).
- ❖ Alleviation: Perform the morph on one vertex/edge at a time.

## Compressions

Want to be able to bring a drawing to an  $O(n) \times O(n)$  grid from an arbitrarily sized grid (where the constant is independent of the initial grid size).

Idea: Sort  $x$ -coordinates.



This is a unidirectional morph.

## Simplification—Removing all Horizontal Zig-Zags (High-level)

Each problem has a different solution:

- ❖ Requires a morph for each zig-zag (want  $O(1)$  morphs for all zig-zags).
  - ❖ Van Goethem et al.: A single morph suffices for many (disjoint) horizontal zig-zags.

## Simplification—Removing all Horizontal Zig-Zags (High-level)

Each problem has a different solution:

- ❖ Requires a morph for each zig-zag (want  $O(1)$  morphs for all zig-zags).
  - ❖ Van Goethem et al.: A single morph suffices for many (disjoint) horizontal zig-zags. Two issues with their solution:
    - ▶ Uses a large grid.
    - ▶ Slow time complexity.

## Simplification—Removing all Horizontal Zig-Zags (High-level)

Each problem has a different solution:

- ❖ Requires a morph for each zig-zag (want  $O(1)$  morphs for all zig-zags).
  - ❖ Van Goethem et al.: A single morph suffices for many (disjoint) horizontal zig-zags.  
Two issues with their solution:
    - ▶ Uses a large grid.
    - ▶ Slow time complexity.
- ❖ Requires  $O(n)$  time for each zig-zag (want  $O(n)$  time for all zig-zags).
  - ❖ Use circuit layout compaction!

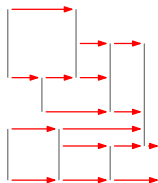
## Simplification—Circuit Compaction

**Goal:** Compress vertical line segments.

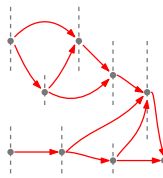
Solution by Doenhardt and Lengauer:



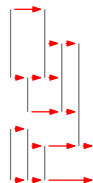
(1) Input



(2) Trapezoidal Map



(3) Trapezoidal Graph



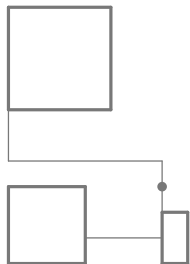
(4) Result from  
topological sort

Important note:

Last step of Doenhardt and Lengauer's algorithm only needs y-coordinates and trapezoidal graph.

## Simplification—Circuit Compaction for Box Drawings

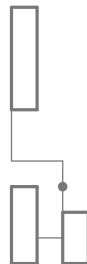
**Goal:** Compress a box drawing (again).



A box drawing  $C$



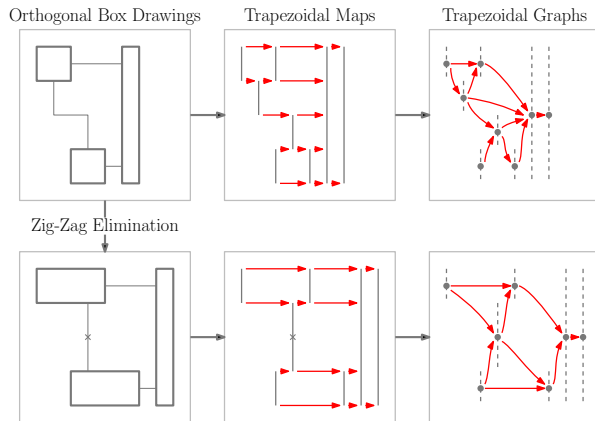
A set of maximal vertical line segments  $L(C)$  covering  $C$



The compressed drawing  $C'$

Side note: Doing this in  $O(n)$  time requires connectivity (via an algorithm by Chazelle for trapezoidal maps of simple polygons).

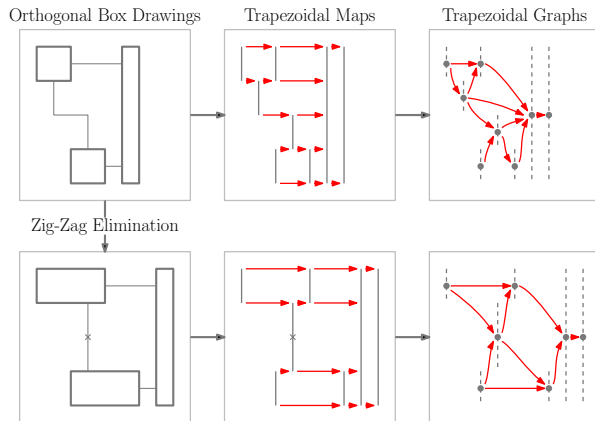
## Simplification—Zig-Zag Elimination and Circuit Compaction



Takeaway: The changes to the trapezoidal graph are local to the zig-zag being eliminated.



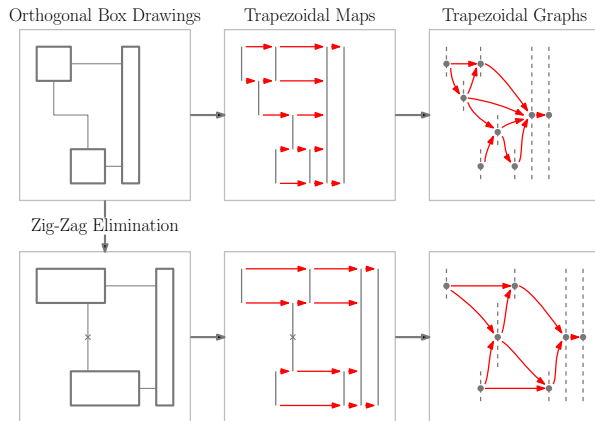
## Simplification—Zig-Zag Elimination and Circuit Compaction



Takeaway: The changes to the trapezoidal graph are local to the zig-zag being eliminated.

Recall: Last step of Doenhardt and Lengauer's algorithm only needs  $y$ -coordinates and trapezoidal graph.

## Simplification—Zig-Zag Elimination and Circuit Compaction

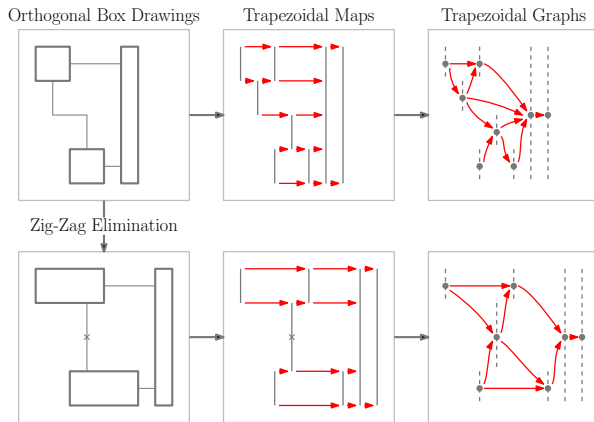


Takeaway: The changes to the trapezoidal graph are local to the zig-zag being eliminated.

Recall: Last step of Doenhardt and Lengauer's algorithm only needs y-coordinates and trapezoidal graph.

Idea: Compute only the trapezoidal graph after a sequence of zig-zag eliminations.

# Simplification—Zig-Zag Elimination and Circuit Compaction



Takeaway: The changes to the trapezoidal graph are local to the zig-zag being eliminated.

Recall: Last step of Doenhardt and Lengauer's algorithm only needs y-coordinates and trapezoidal graph.

Idea: Compute only the trapezoidal graph after a sequence of zig-zag eliminations.

Final result: Can remove all horizontal zig-zags in one linear morph, in  $O(n)$  time.



## Simplification—Eliminating All Zig-Zags

Eliminating all horizontal zig-zags  $\neq$  eliminating all zig-zags:



Eliminating all horizontal (and then vertical) zig-zags does reduce the number of bends per edge (unless there are no zig-zags).

## Simplification—Eliminating All Zig-Zags

Eliminating all horizontal zig-zags  $\neq$  eliminating all zig-zags:



Eliminating all horizontal (and then vertical) zig-zags does reduce the number of bends per edge (unless there are no zig-zags).

Idea: Since  $O(1)$  bends per edge is maintained, only need to do  $O(1)$  simultaneous eliminations to eliminate all zig-zags.

## Phase I High-Level

