

The Density Formula

One Lemma to Bound them All

Michael Kaufmann
Universität Tübingen

Boris Klemz
Universität Würzburg

Kristin Knorr
Freie Universität Berlin

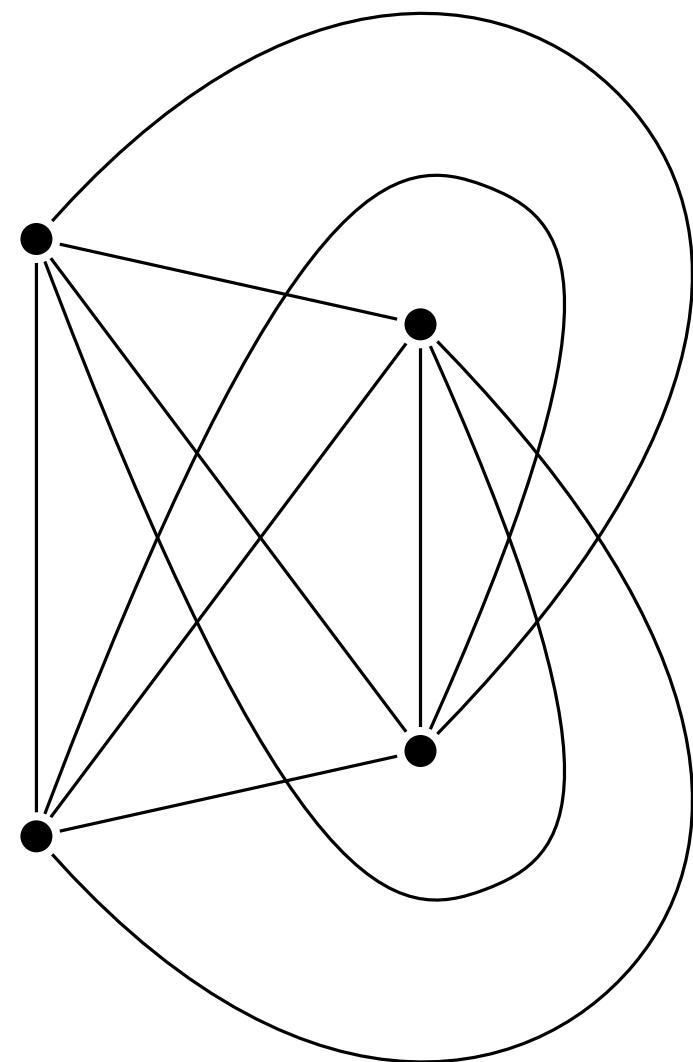
Meghana M. Reddy
ETH Zürich

Felix Schröder
Charles University Prague

Torsten Ueckerdt
Karlsruhe Institute of Technology

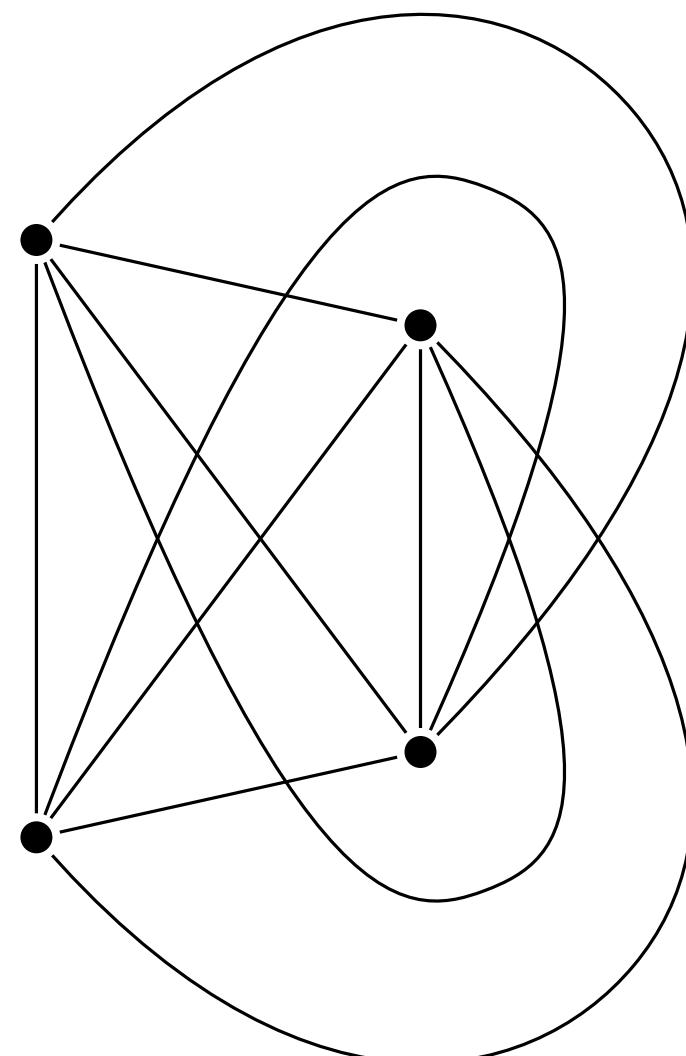
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Vienna, September 18, 2024

Drawings



- ▷ no loops
- ▷ parallel edges allowed
- ▷ connected, $n \geq 3$ vertices
- ▷ no self-crossings
- ▷ finite number of crossings
- ▷ no three edges through 1 point
- ▷ no touchings

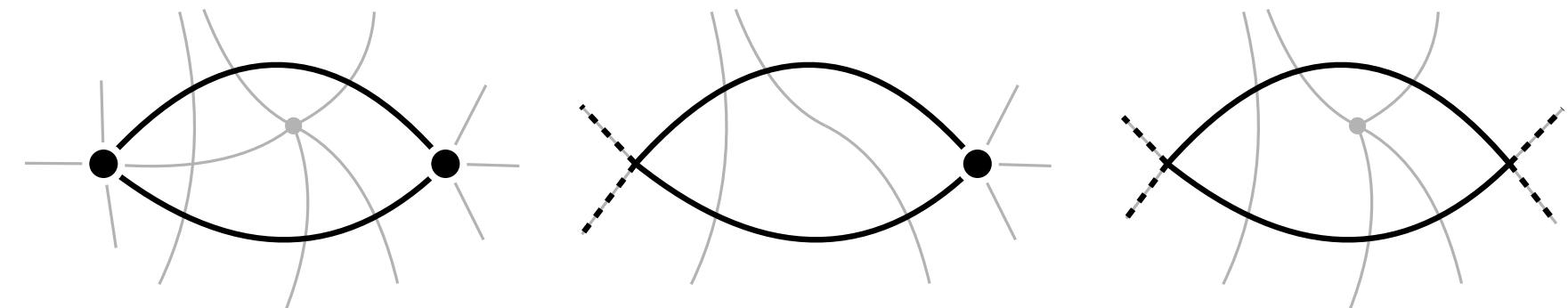
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non-homotopic and simple drawings:

lens = region bounded by two edge parts



non-homotopic drawing

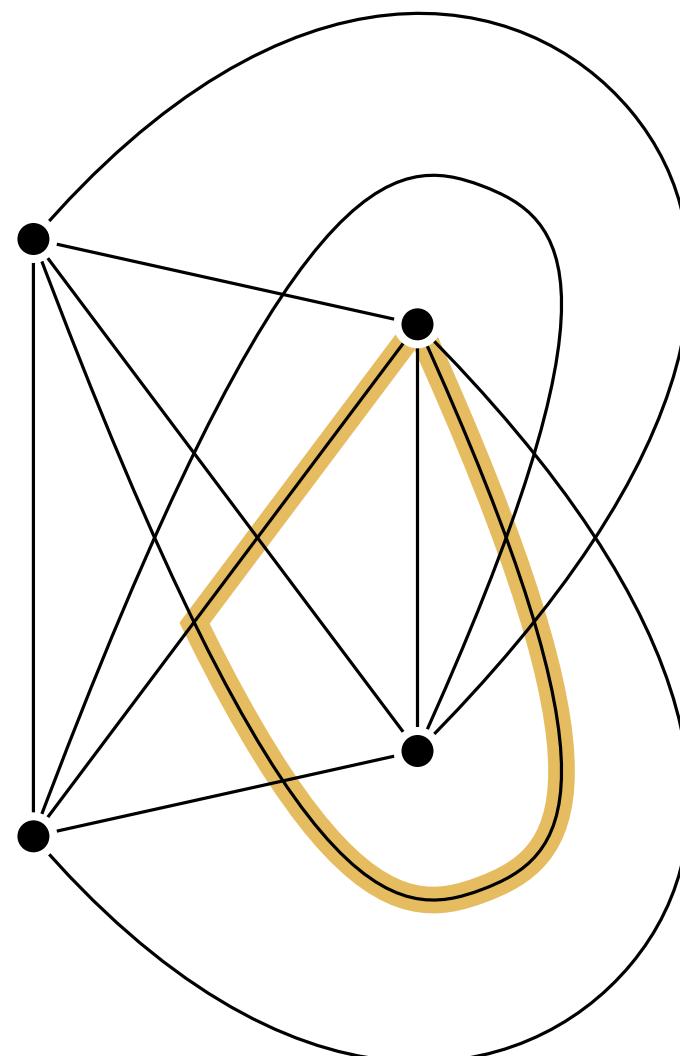
Every lens contains
a vertex or crossing.

simple drawing

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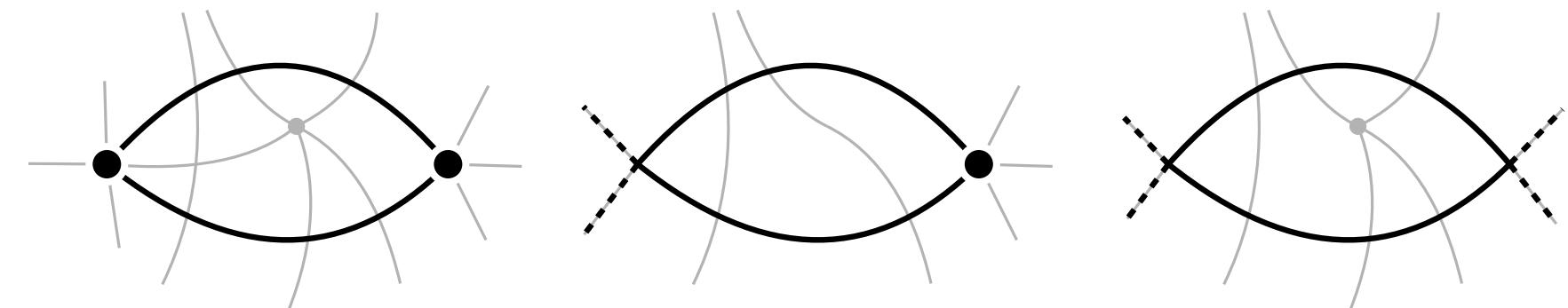
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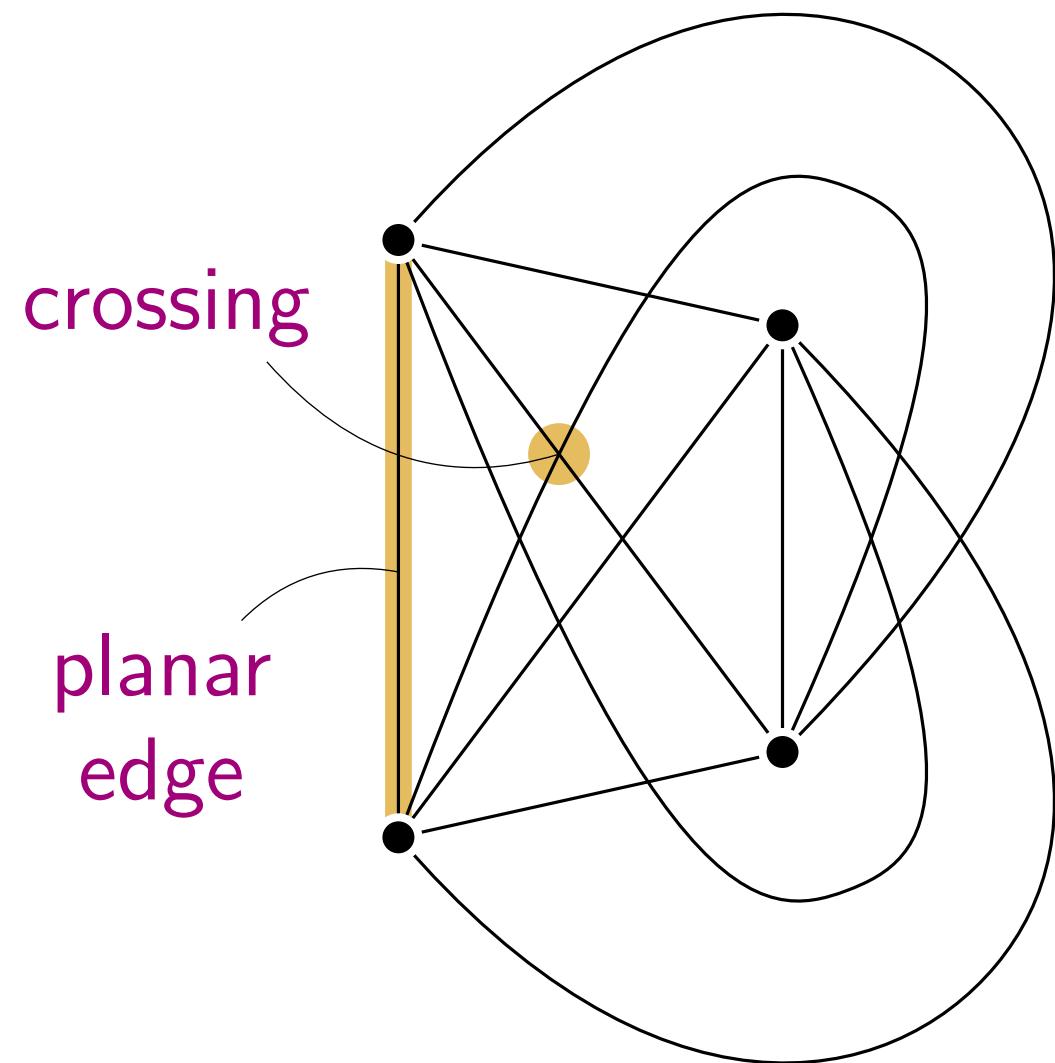
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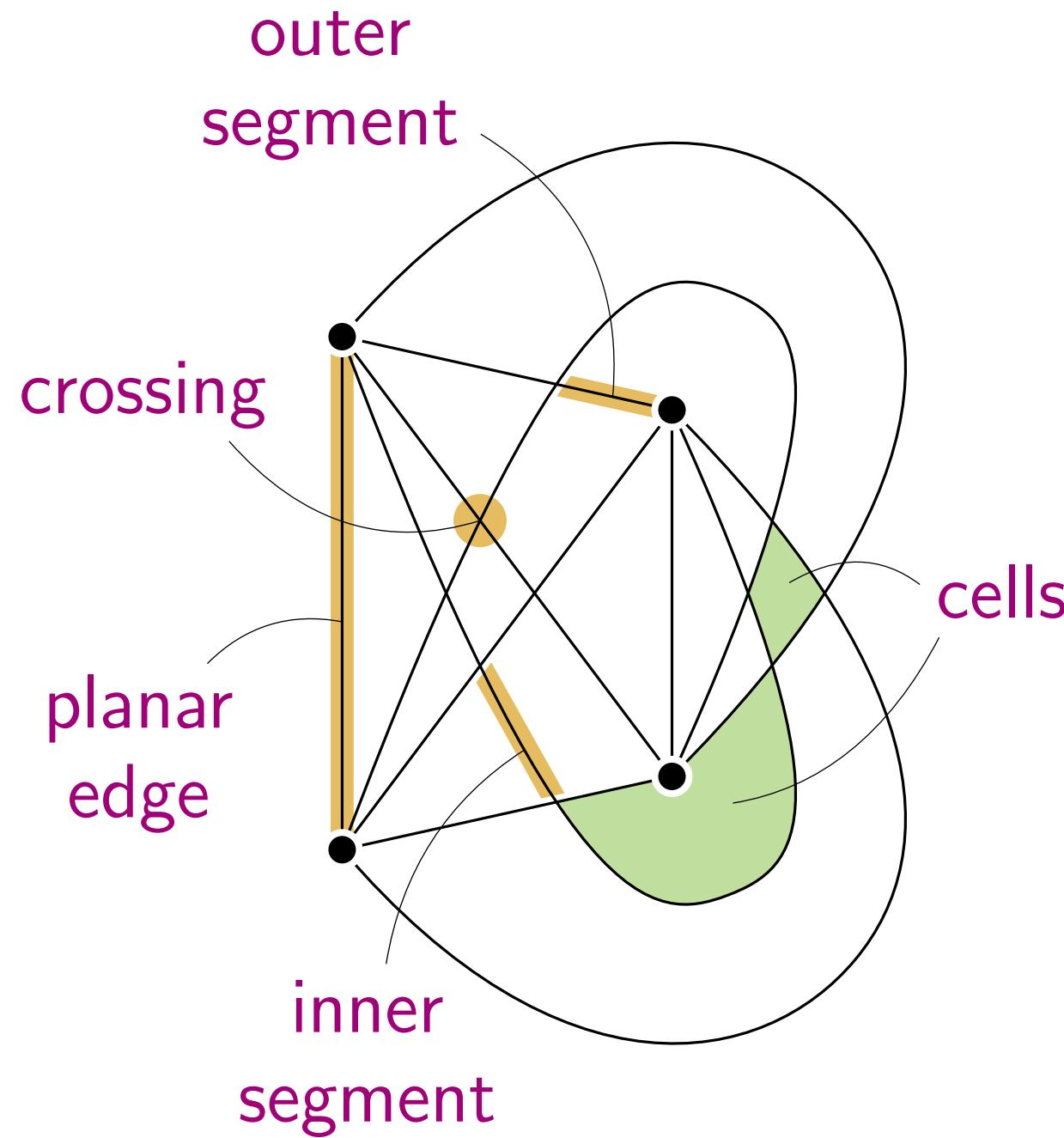
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$$E = \{ \text{edges} \} = E_x \cup E_p$$

crossed
planar

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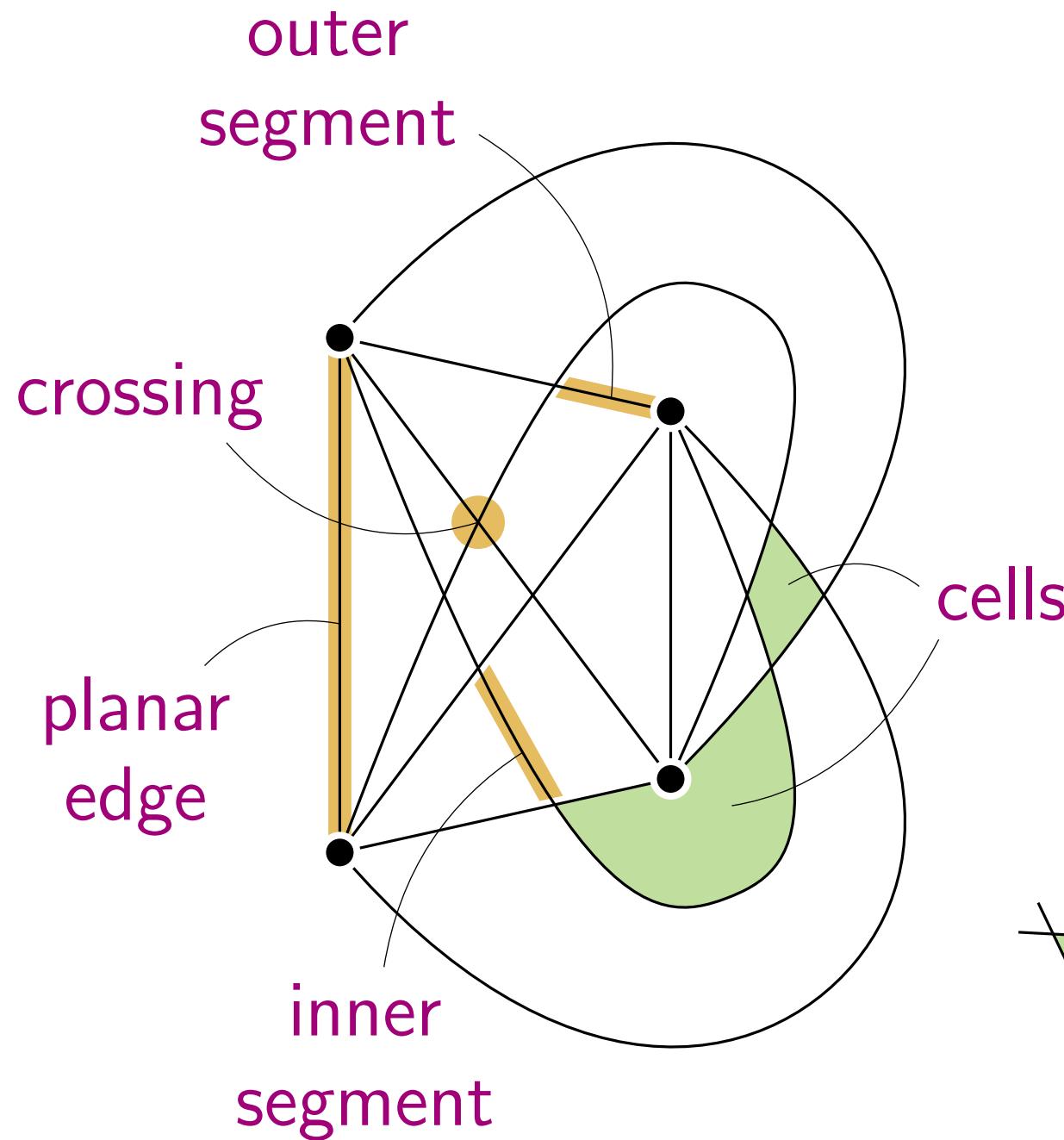
$$E = \{ \text{edges} \} = E_x \cup E_p$$

$$\mathcal{S} = \{ \text{segments} \} = \mathcal{S}_{\text{in}} \cup \mathcal{S}_{\text{out}}$$

$$\mathcal{C} = \{ \text{cells} \}$$

crossed
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inner
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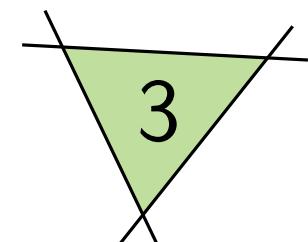
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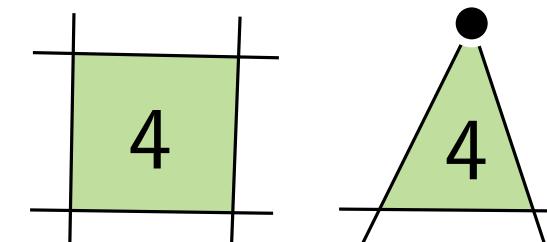
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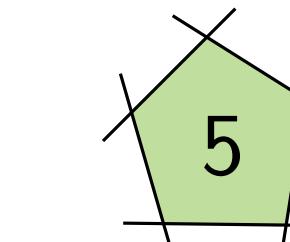
$\|c\| = \text{size of cell } c = \# \text{vertices} + \# \text{ segments}$



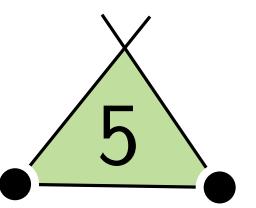
\mathcal{C}_3



\mathcal{C}_4



\mathcal{C}_5



Density Formula

Theorem. (Density Formula)

For any $t \in \mathbb{R}$ and any connected drawing
of any graph $G = (V, E)$, $|E| \geq 1$, we have

$$|E| = t(|V| - 2) - \sum_{c \in \mathcal{C}} \left(\frac{t-1}{4} ||c|| - t \right) - |\mathcal{X}|$$

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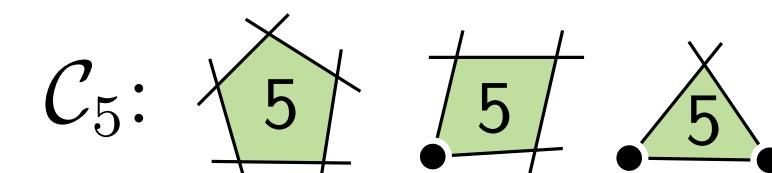
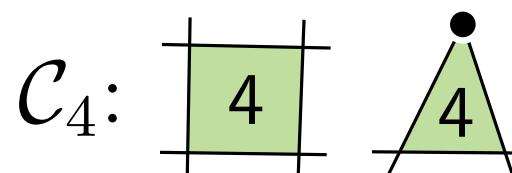
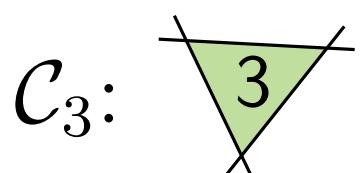
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special cases

$$t = 4 \quad |E| = 4(|V| - 2) + \frac{7}{4} |\mathcal{C}_3| + |\mathcal{C}_4| + \frac{1}{4} |\mathcal{C}_5| - \frac{1}{2} |\mathcal{C}_6| - \cdots - |\mathcal{X}|$$

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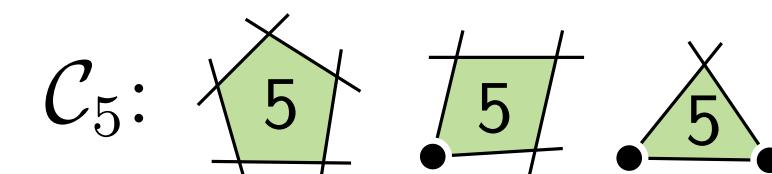
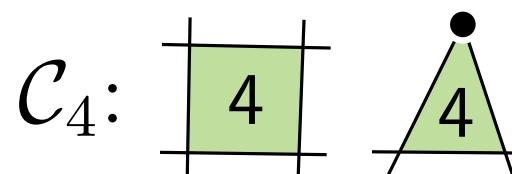
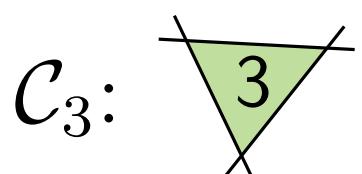
General approach:

- ▷ apply the formula with a specific t
- ▷ upper bounds on $|\mathcal{C}_3|, |\mathcal{C}_4|, \dots$ in terms of $|\mathcal{X}|$ and/or $|V|, |E|$

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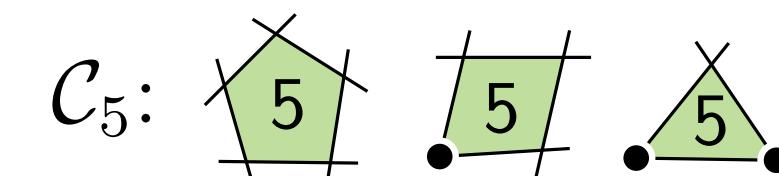
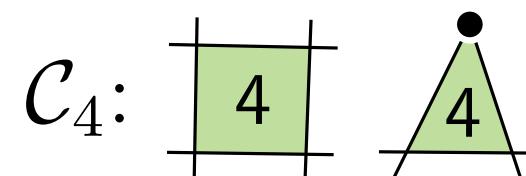
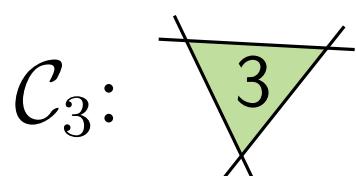
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1-planar graphs

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Theorem (Pach + Tóth, 1997)

n -vertex 1-planar graphs have $\leq 4n - 8$ edges.

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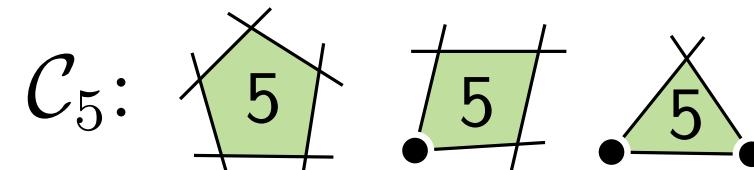
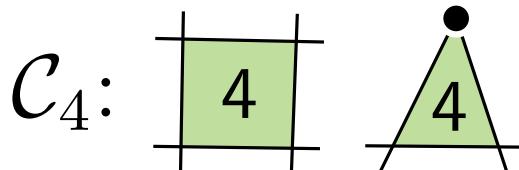
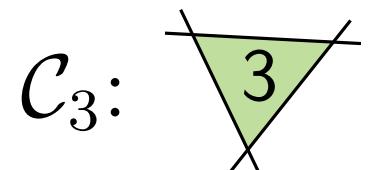
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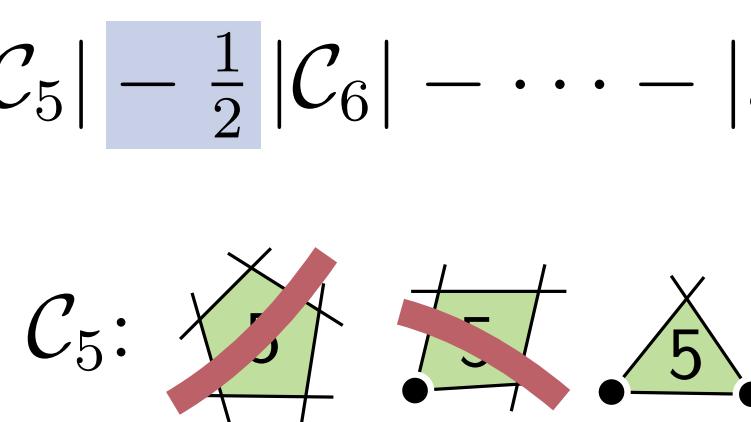
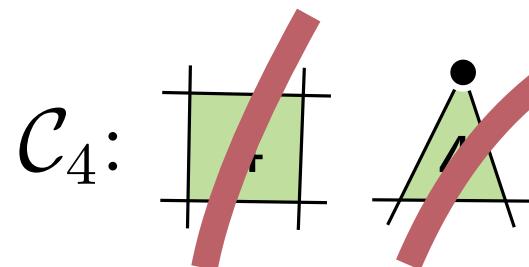
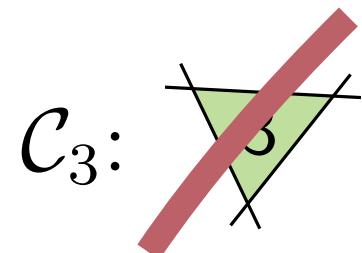
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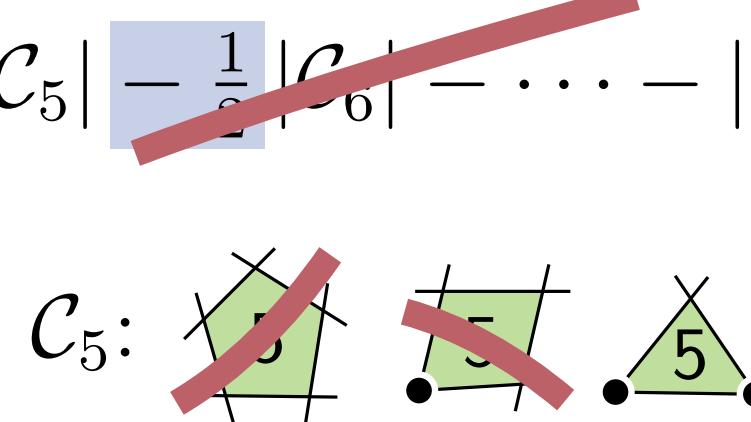
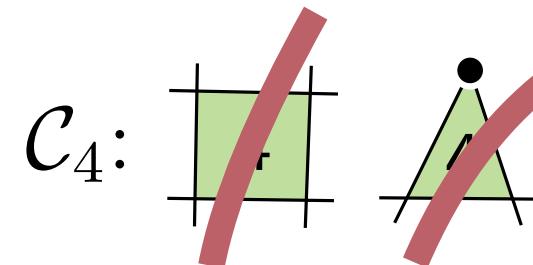
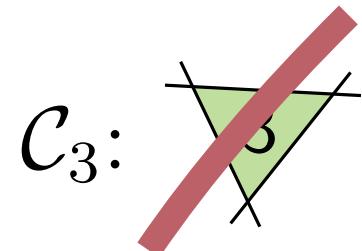
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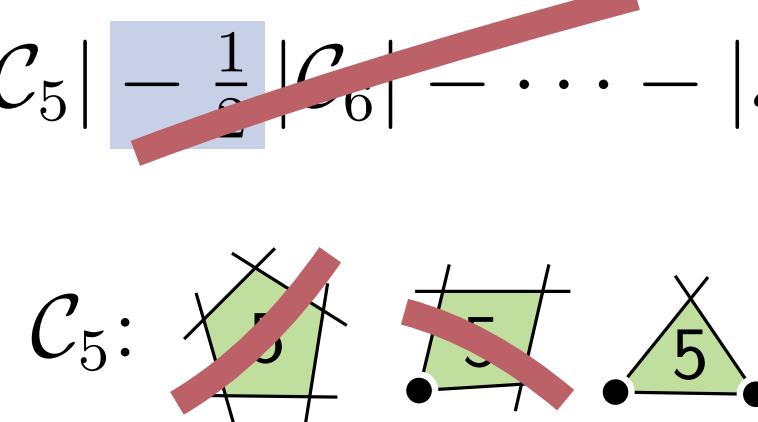
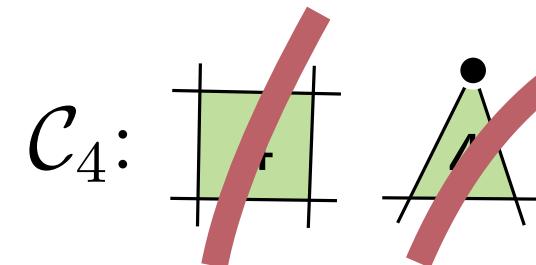
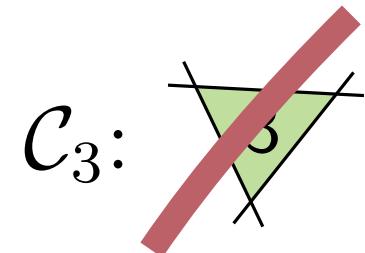
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$$\# \text{---} 5\text{-cells} \leq 4|\mathcal{X}| \implies \frac{1}{4}|\mathcal{C}_5| \leq |\mathcal{X}| \implies |E| \leq 4n - 8$$

□

If $|E| = 4n - 8$, all cells must be



Overview

	density =	max. # edges for n vertices
1-planar	$4n - 8$	Pach-Tóth '97
2-planar	$5n - 10$	Pach-Tóth '97
3-planar	$5.5n - \Theta(1)$	Pach-Radoičić-Tardos-Tóth '06
4-planar	$6n - \Theta(1)$	Ackerman '15
simple quasiplanar	$6.5n - 20$	Ackerman-Tardos '07
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0-bend RAC	$4n - 10$	Didimo-Eades-Liotta '11
1-bend RAC	$5n - 10$	Kaufmann-Klemz-Knorr-Reddy-S.-Ueckerdt '24
2-bend RAC	$10n - \Theta(1)$	Kaufmann-Klemz-Knorr-Reddy-S.-Ueckerdt '24
1^+ -real face	$5n - 10$	Binucci-Di Battista-Didimo-Hong-Kaufmann-Liotta-Morin-Tappini '23
2^+ -real face	$4n - 8$	Binucci-Di Battista-Didimo-Hong-Kaufmann-Liotta-Morin-Tappini '23
k^+ -real face	$\frac{k}{k-2}(n - 2)$	Binucci-Di Battista-Didimo-Hong-Kaufmann-Liotta-Morin-Tappini '23
fan-planar	$5n - 10$	(Kaufmann-Ueckerdt '22) Ackerman-Keszegh '23

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Proof.

- ▷ $|\mathcal{S}| = |E| + 2|\mathcal{X}|$
- ▷ $\sum_{c \in \mathcal{C}} ||c|| = 2|\mathcal{S}| + \sum_{v \in V} \deg(v) = 4|E| + 4|\mathcal{X}|$

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 $\implies 0 = |V| - 2 - (|E| + |\mathcal{X}| - |\mathcal{C}|) \implies 0 = |V| - 2 - \sum_{c \in \mathcal{C}} \left(\frac{||c||}{4} - 1 \right)$

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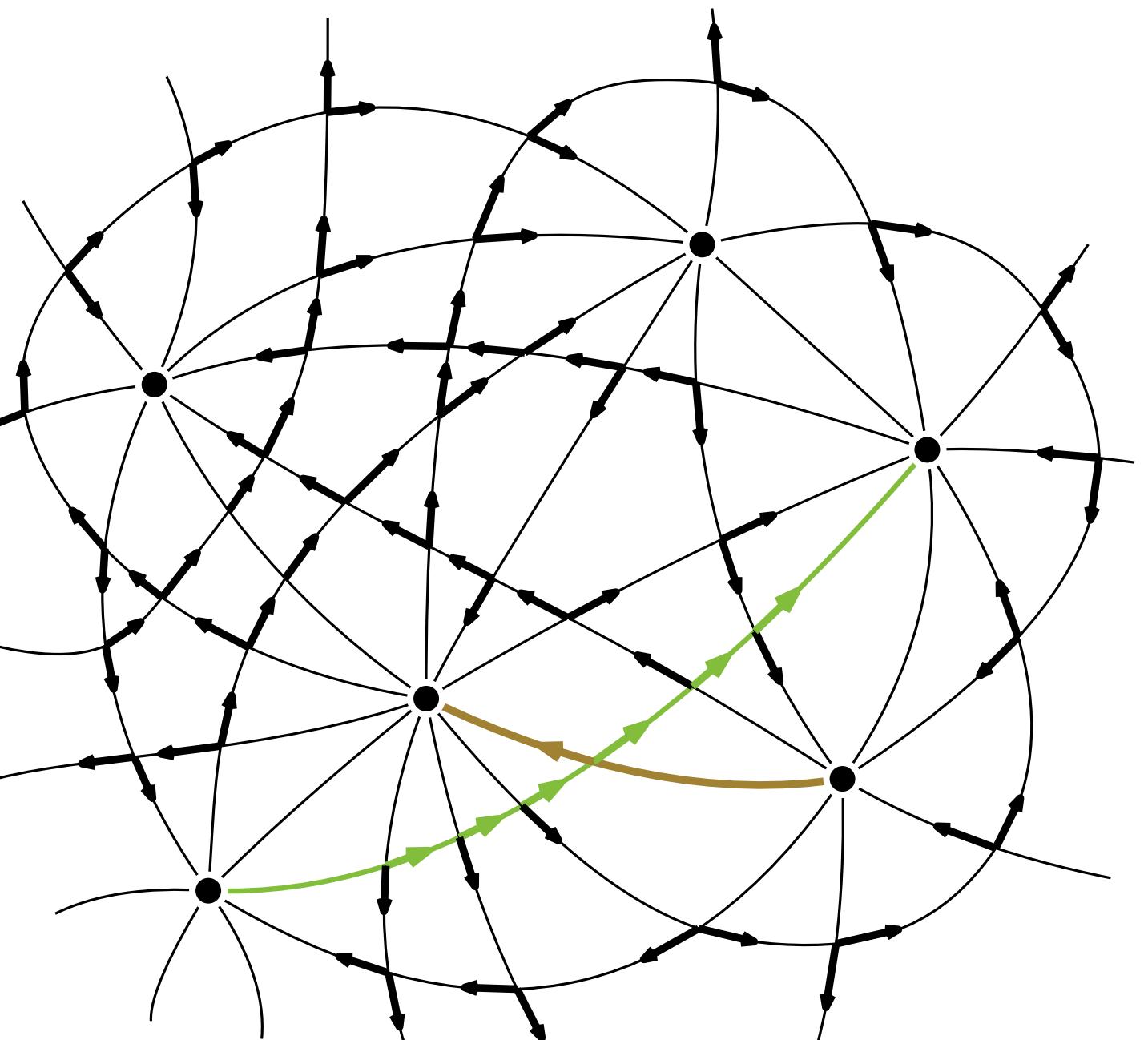
Proof.

- ▷ $|\mathcal{S}| = |E| + 2|\mathcal{X}|$
- ▷ $\sum_{c \in \mathcal{C}} ||c|| = 2|\mathcal{S}| + \sum_{v \in V} \deg(v) = 4|E| + 4|\mathcal{X}| \implies \sum_{c \in \mathcal{C}} \left(\frac{||c||}{4} - 1 \right) = |E| + |\mathcal{X}| - |\mathcal{C}|$
- ▷ **planarization** has $|V| + |\mathcal{X}|$ vertices, $|\mathcal{S}|$ edges, $|\mathcal{C}|$ faces
- ▷ **Euler's Formula:** $|V| + |\mathcal{X}| - |\mathcal{S}| + |\mathcal{C}| = 2 \implies |E| = |V| - 2 + |\mathcal{C}| - |\mathcal{X}| \quad | \times 1$
 $\implies 0 = |V| - 2 - (|E| + |\mathcal{X}| - |\mathcal{C}|) \implies 0 = |V| - 2 - \sum_{c \in \mathcal{C}} \left(\frac{||c||}{4} - 1 \right) \quad | \times (t-1)$

$$|E| = t(|V| - 2) - \sum_{c \in \mathcal{C}} \left(\frac{t-1}{4} ||c|| - t \right) - |\mathcal{X}|$$

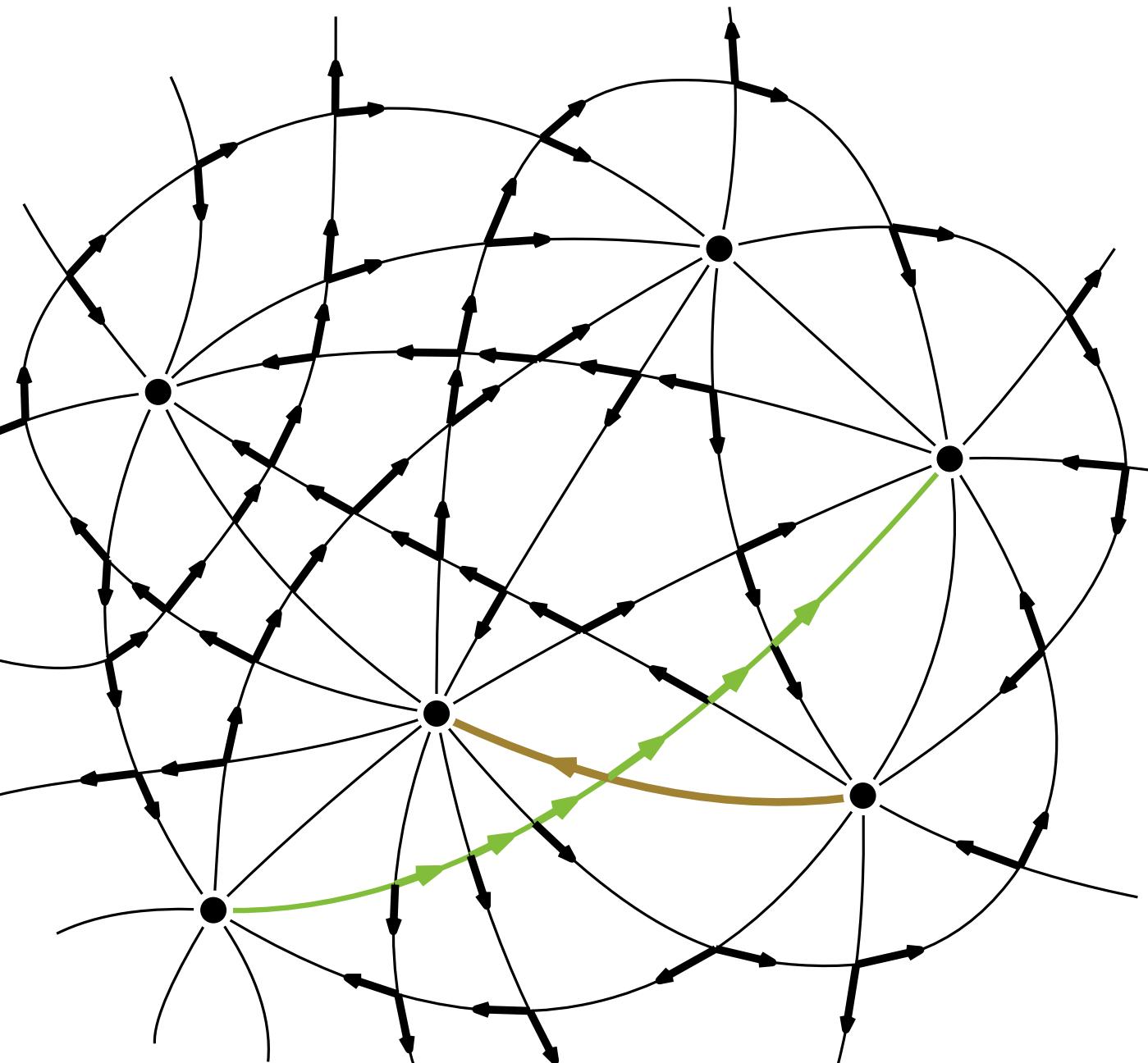
Useful Lemmas

$$|\mathcal{S}_{\text{in}}| = 2|\mathcal{X}| - |E_x|$$



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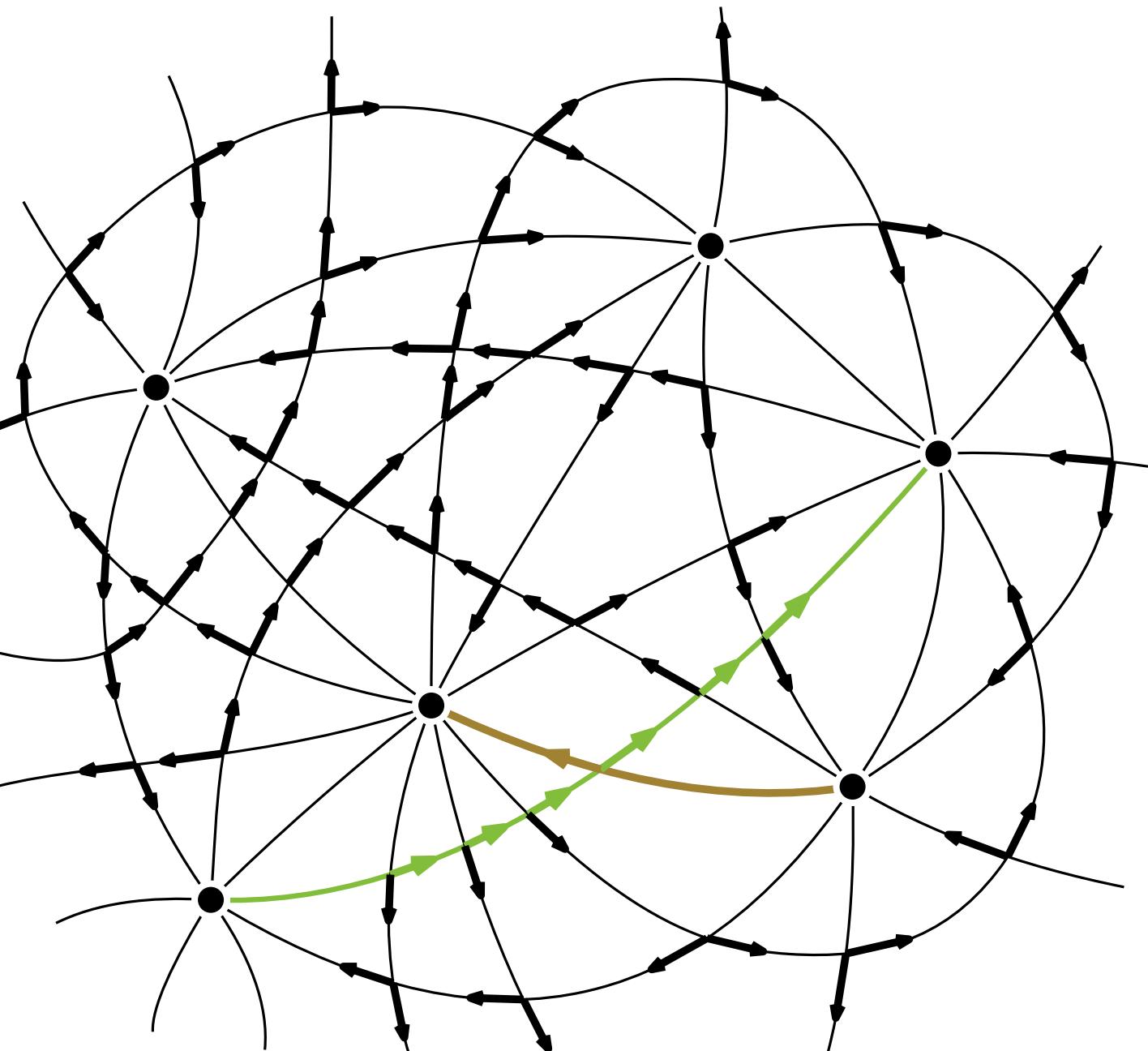
Lemma

In any non-homotopic drawing we have

$$|\mathcal{S}_{\text{in}}| \geq 3\# \begin{array}{c} \diagup \\ \backslash \end{array}^3\text{-cells} + 2\# \begin{array}{|c|c|} \hline & 4 \\ \hline \end{array}^4\text{-cells} + \# \begin{array}{c} \diagdown \\ / \end{array}^4\text{-cells}$$

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Lemma

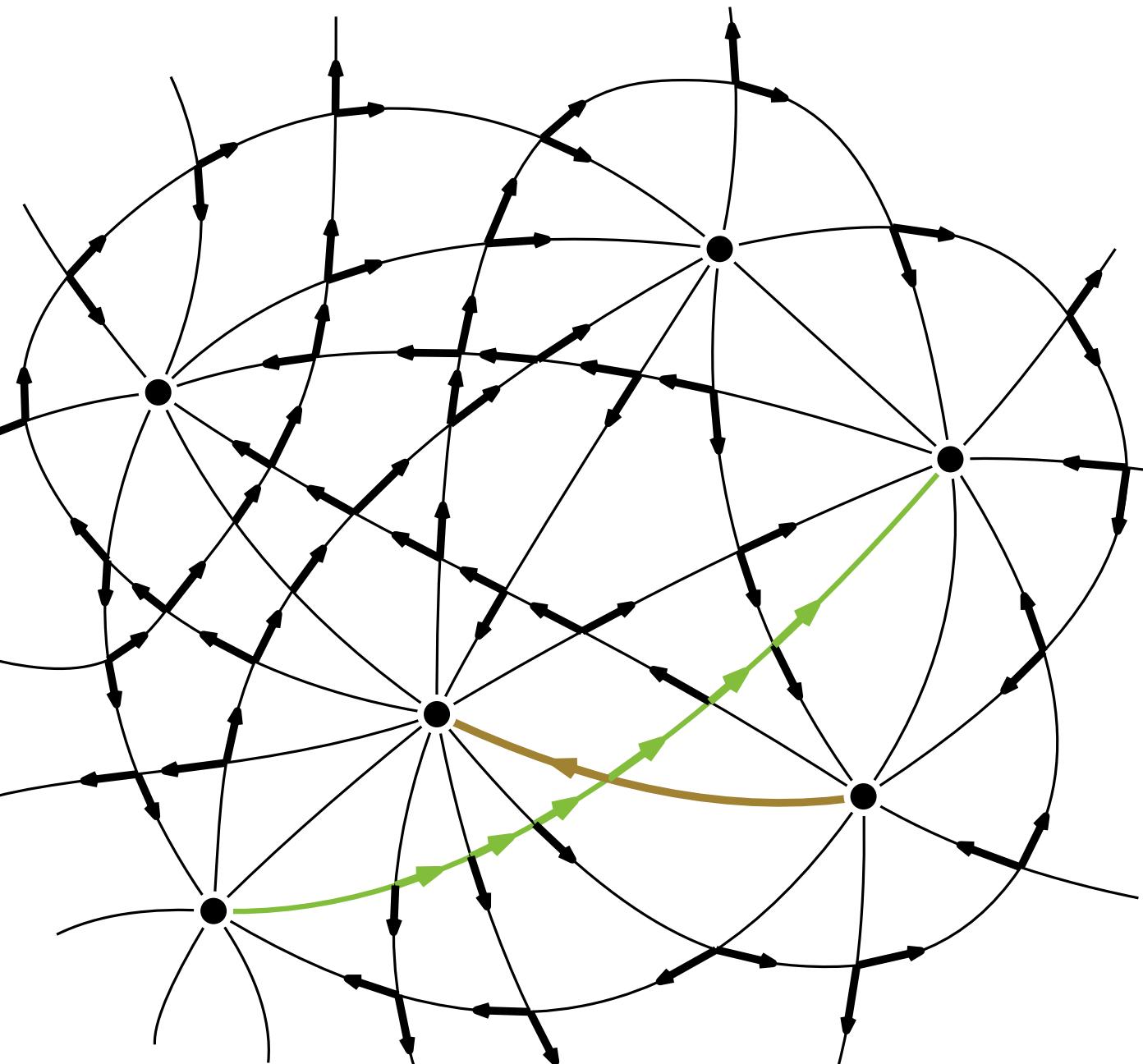
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Proof idea

Useful Lemmas

$$|\mathcal{S}_{\text{in}}| = 2|\mathcal{X}| - |E_x|$$

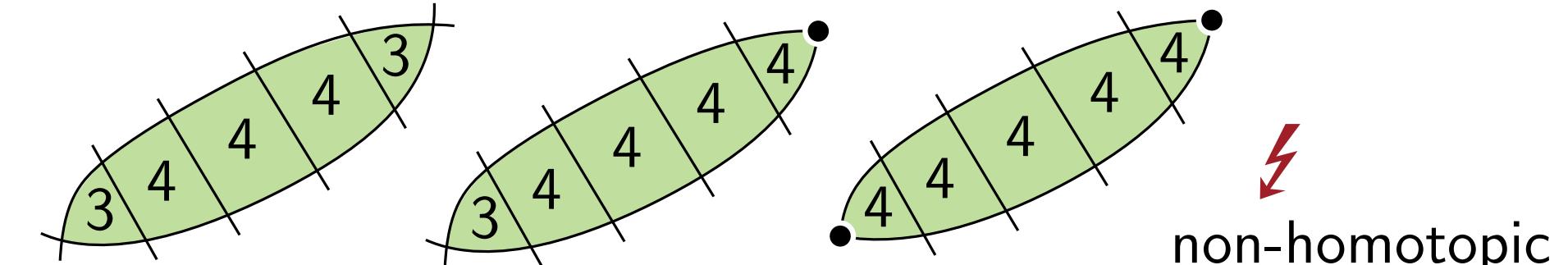


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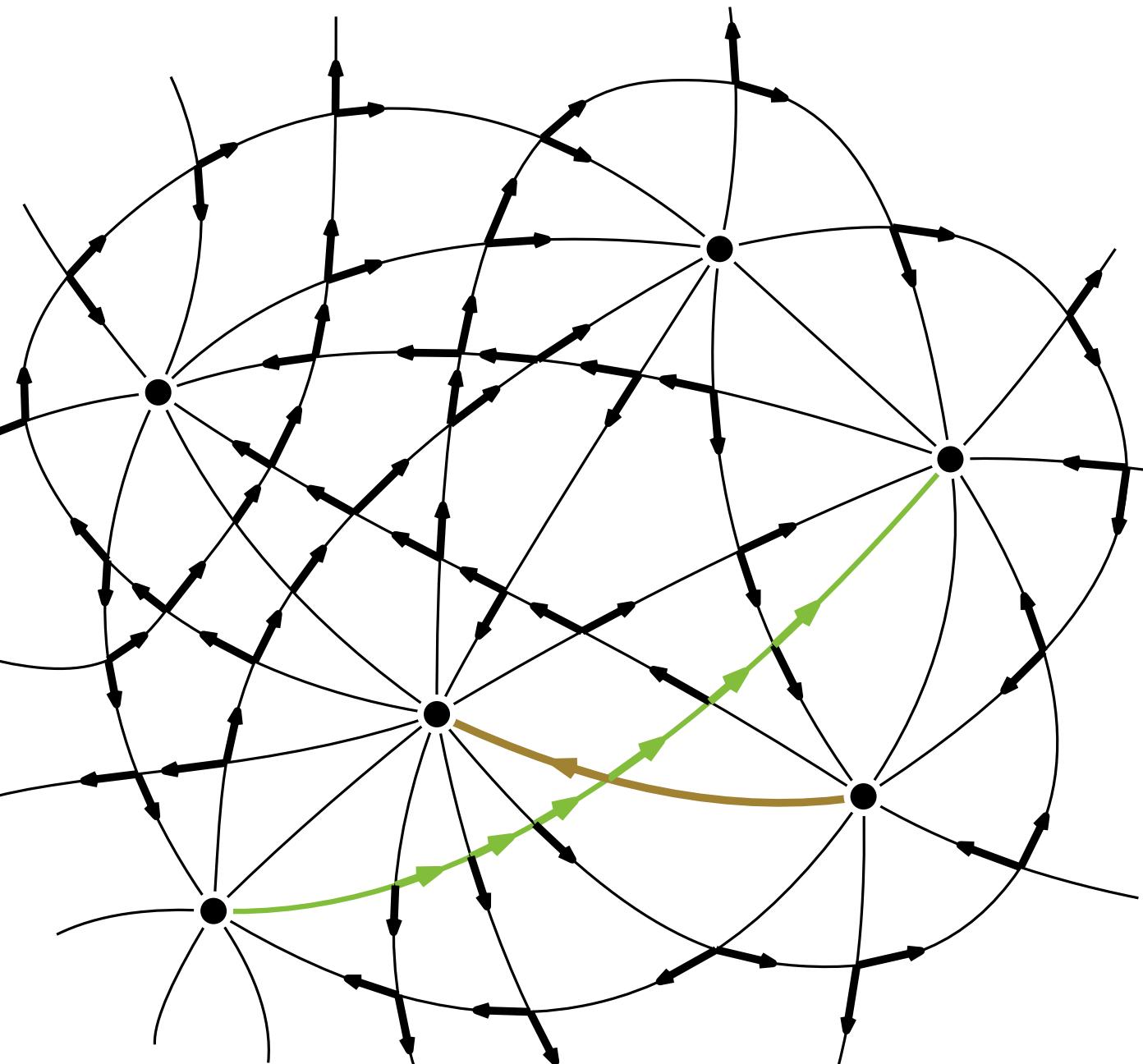
Proof idea



non-homotopic

Useful Lemmas

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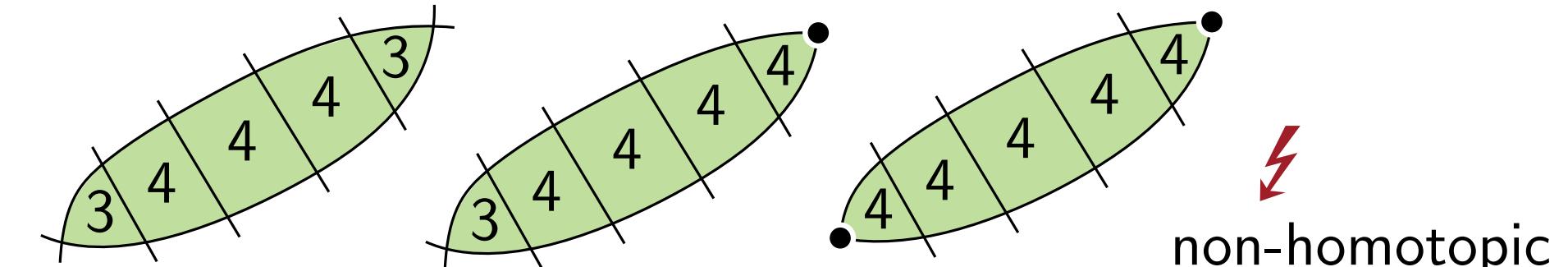


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Proof idea



non-homotopic

2-planar graphs

$$|\mathcal{S}_{\text{in}}| \geq 3\# \begin{array}{c} \diagup \\ \diagdown \end{array}^3\text{-cells} + 2\# \begin{array}{c} \diagup \diagdown \\ \diagup \diagdown \end{array}^4\text{-cells} + \# \begin{array}{c} \diagup \\ \diagdown \end{array}^4\text{-cells}$$

$$|E| = t(|V| - 2) - \sum_{c \in \mathcal{C}} \left(\frac{t-1}{4} ||c|| - t \right) - |\mathcal{X}|$$

Theorem (Pach + Tóth, 1997)

n -vertex 2-planar graphs have $\leq 5n - 10$ edges

2-planar graphs

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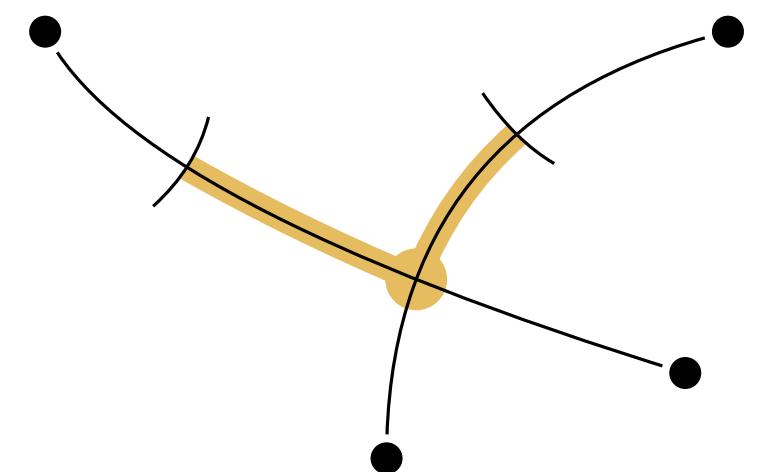
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n -vertex 2-planar graphs have $\leq 5n - 10$ edges

Proof using the Density Formula:

2 crossings at each inner segment
 ≤ 2 inner segments at each crossing

$$\implies |\mathcal{X}| \geq |\mathcal{S}_{\text{in}}|$$



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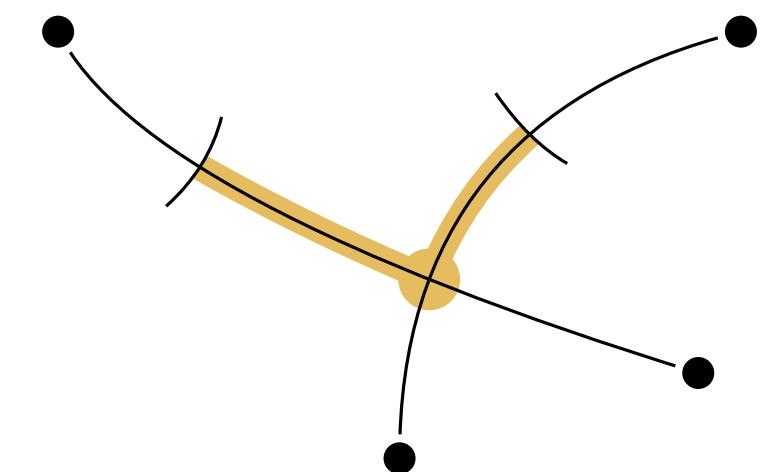
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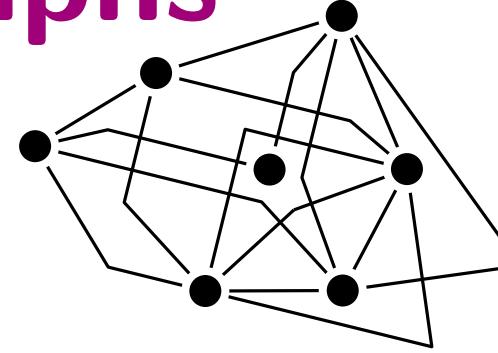
$$t = 5 \quad |E| = 5(n - 2) + 2|\mathcal{C}_3| + |\mathcal{C}_4| + 0|\mathcal{C}_5| - 1|\mathcal{C}_6| - \dots - |\mathcal{X}|$$

$$\implies |E| \leq 5(n - 2) + |\mathcal{S}_{\text{in}}| - |\mathcal{X}| \leq 5(n - 2)$$

□

k-bend RAC graphs

edges with $\leq k$ bends
Right-Angle Crossings



Theorem.

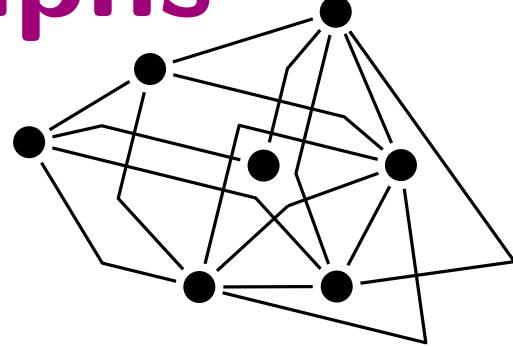
- n -vertex non-hom. 1-bend RAC graphs have $\leq 5n - 10$ edges.
- n -vertex non-hom. 2-bend RAC graphs have $\leq 10n - 19$ edges.

$$2|\mathcal{X}| - |E_x| = |\mathcal{S}_{\text{in}}| \geq 3\# \begin{array}{c} \diagup \\ \diagdown \end{array}^3 + 2\# \begin{array}{|c|c|} \hline \diagup & \diagdown \\ \hline \end{array}^4 + \# \begin{array}{c} \diagup \\ \diagdown \end{array}^4$$

$$t = 5 : \quad |E| \leq 5(n - 2) + 2|\mathcal{C}_3| + |\mathcal{C}_4| - |\mathcal{X}|$$

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Right-Angle Crossings



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$$t = 5 : \quad |E| \leq 5(n-2) + 2|\mathcal{C}_3| + |\mathcal{C}_4| - |\mathcal{X}|$$

Theorem.

n -vertex non-hom. 1-bend RAC graphs have $\leq 5n - 10$ edges.

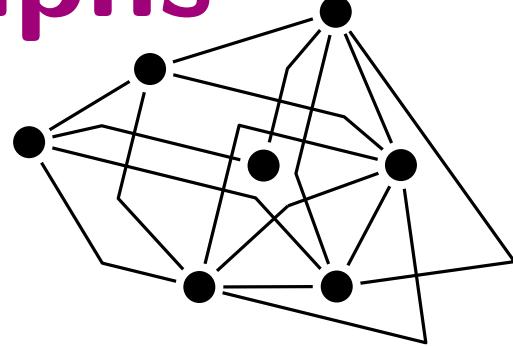
n -vertex non-hom. 2-bend RAC graphs have $\leq 10n - 19$ edges.

Proof.

≥ 1 convex bend at each* $\begin{array}{c} \diagup \\ \diagdown \end{array}^3$ -cell and $\begin{array}{c} \diagup \\ \diagdown \end{array}^4$ -cell $\implies k|E_x| \geq \# \begin{array}{c} \diagup \\ \diagdown \end{array}^3 + \# \begin{array}{c} \diagup \\ \diagdown \end{array}^4 - 1$

k-bend RAC graphs

edges with $\leq k$ bends
Right-Angle Crossings



Theorem.

- n -vertex non-hom. 1-bend RAC graphs have $\leq 5n - 10$ edges.
- n -vertex non-hom. 2-bend RAC graphs have $\leq 10n - 19$ edges.

Proof.

$$\geq 1 \text{ convex bend at each}^* \begin{array}{c} \diagup \\ \diagdown \end{array} \begin{array}{c} 3 \\ \times \end{array} \text{-cell and } \begin{array}{c} \diagup \\ \diagdown \end{array} \begin{array}{c} 4 \\ \square \end{array} \text{-cell} \implies k|E_x| \geq \# \begin{array}{c} \diagup \\ \diagdown \end{array} \begin{array}{c} 3 \\ \times \end{array} + \# \begin{array}{c} \diagup \\ \diagdown \end{array} \begin{array}{c} 4 \\ \square \end{array} - 1$$

$$\implies 2|\mathcal{X}| + (k-1)|E_x| \geq 4\# \begin{array}{c} \diagup \\ \diagdown \end{array} \begin{array}{c} 3 \\ \times \end{array} + 2\# \begin{array}{c} \diagup \\ \diagdown \end{array} \begin{array}{c} 4 \\ \square \end{array} + 2\# \begin{array}{c} \diagup \\ \diagdown \end{array} \begin{array}{c} 4 \\ \triangle \end{array} - 1 = 4|\mathcal{C}_3| + 2|\mathcal{C}_4| - 1$$

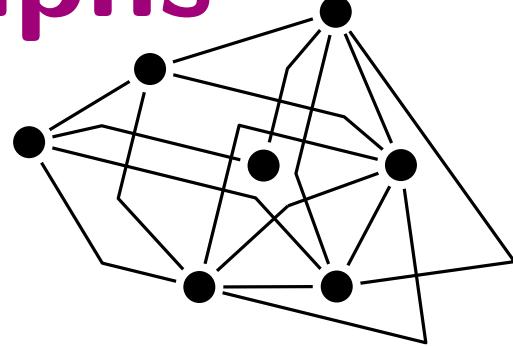
$$\implies 2|\mathcal{C}_3| + |\mathcal{C}_4| - |\mathcal{X}| \leq \frac{k-1}{2}|E_x| + \frac{1}{2}$$

$$2|\mathcal{X}| - |E_x| = |\mathcal{S}_{\text{in}}| \geq 3\# \begin{array}{c} \diagup \\ \diagdown \end{array} \begin{array}{c} 3 \\ \times \end{array} + 2\# \begin{array}{c} \diagup \\ \diagdown \end{array} \begin{array}{c} 4 \\ \square \end{array} + \# \begin{array}{c} \diagup \\ \diagdown \end{array} \begin{array}{c} 4 \\ \triangle \end{array}$$

$$t = 5 : \quad |E| \leq 5(n-2) + 2|\mathcal{C}_3| + |\mathcal{C}_4| - |\mathcal{X}|$$

k-bend RAC graphs

edges with $\leq k$ bends
Right-Angle Crossings



Theorem.

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 n -vertex non-hom. 2-bend RAC graphs have $\leq 10n - 19$ edges.

Proof.

$$\geq 1 \text{ convex bend at each* } \begin{array}{c} \diagup \\ \diagdown \end{array} \text{-cell and } \begin{array}{c} \diagup \\ \diagdown \end{array} \text{-cell} \implies k|E_x| \geq \# \begin{array}{c} \diagup \\ \diagdown \end{array} + \# \begin{array}{c} \diagup \\ \diagdown \end{array} - 1$$

$$\implies 2|\mathcal{X}| + (k-1)|E_x| \geq 4\# \begin{array}{c} \diagup \\ \diagdown \end{array} + 2\# \begin{array}{|c|c|} \hline \diagup & \diagdown \\ \hline \end{array} + 2\# \begin{array}{c} \diagup \\ \diagdown \end{array} - 1 = 4|\mathcal{C}_3| + 2|\mathcal{C}_4| - 1$$

$$\implies 2|\mathcal{C}_3| + |\mathcal{C}_4| - |\mathcal{X}| \leq \frac{k-1}{2}|E_x| + \frac{1}{2}$$

$$|E| \leq 5(n-2) + 2|\mathcal{C}_3| + |\mathcal{C}_4| - |\mathcal{X}| \leq 5(n-2) + \frac{k-1}{2}|E| + \frac{1}{2}$$

$$k=1 : |E| \leq 5(n-2) + \frac{1}{2}$$

$$k=2 : \frac{|E|}{2} \leq 5(n-2) + \frac{1}{2} \implies |E| \leq 10(n-2) + 1$$

□

$$2|\mathcal{X}| - |E_x| = |\mathcal{S}_{\text{in}}| \geq 3\# \begin{array}{c} \diagup \\ \diagdown \end{array} + 2\# \begin{array}{|c|c|} \hline \diagup & \diagdown \\ \hline \end{array} + \# \begin{array}{c} \diagup \\ \diagdown \end{array}$$

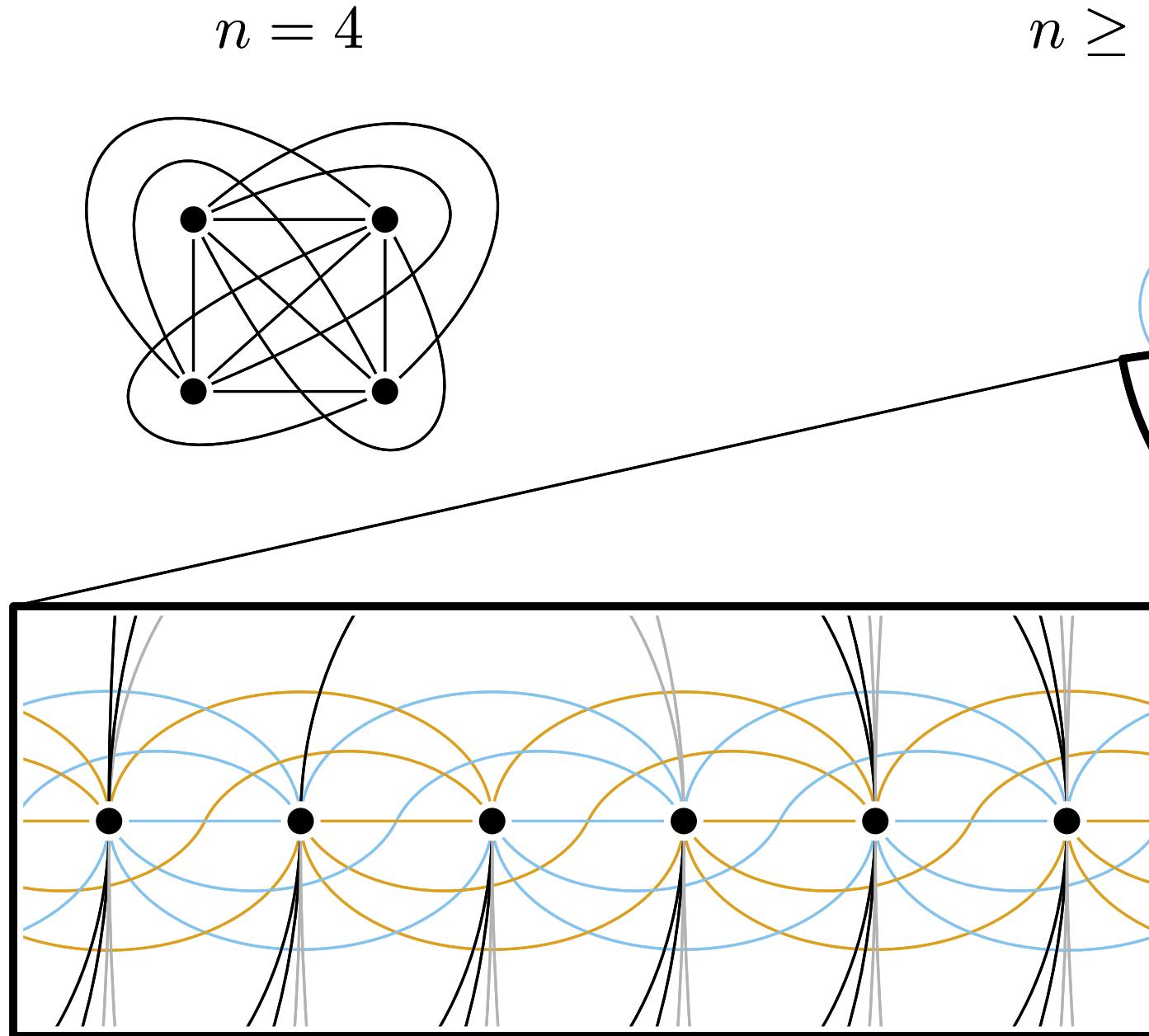
$$t=5 : |E| \leq 5(n-2) + 2|\mathcal{C}_3| + |\mathcal{C}_4| - |\mathcal{X}|$$

Overview

	density =	max. # edges for n vertices	
1-planar	$4n - 8$	Pach-Tóth '97	
2-planar	$5n - 10$	Pach-Tóth '97	
3-planar	$5.5n - \Theta(1)$	Pach-Radoičić-Tardos-Tóth '06	
4-planar	$6n - \Theta(1)$	Ackerman '15	
simple quasiplanar	$6.5n - 20$	Ackerman-Tardos '07	
non-hom. quasiplanar	$8n - 20$	Ackerman-Tardos '07	
0-bend RAC	$4n - 10$	Didimo-Eades-Liotta '11	
1-bend RAC	$5n - 10$	Kaufmann-Klemz-Knorr-Reddy-S.-Ueckerdt '24	
2-bend RAC	$10n - \Theta(1)$	Kaufmann-Klemz-Knorr-Reddy-S.-Ueckerdt '24	
1^+ -real face	$5n - 10$	Binucci-Di Battista-Didimo-Hong-Kaufmann-Liotta-Morin-Tappini '23	
2^+ -real face	$4n - 8$	Binucci-Di Battista-Didimo-Hong-Kaufmann-Liotta-Morin-Tappini '23	
k^+ -real face	$\frac{k}{k-2}(n - 2)$	Binucci-Di Battista-Didimo-Hong-Kaufmann-Liotta-Morin-Tappini '23	
fan-planar	$5n - 10$	(Kaufmann-Ueckerdt '22) Ackerman-Keszegh '23	

Thank You!

Lower Bound for Quasiplanar Graphs



- ▷ cycle C (n edges)
- ▷ 2-hops inside C (n edges)
- ▷ 2-hops outside C (n edges)
- ▷ 3-hops along C (n edges)
- ▷ 2 zig-zag paths inside C ($2(n - 5)$ edges)
- ▷ 2 zig-zag paths outside C ($2(n - 5)$ edges)