UNIVERSITÄT
TRIER

ALGO

On *k*-Plane Insertion into Plane Drawings

Julia Katheder

Philipp Kindermann

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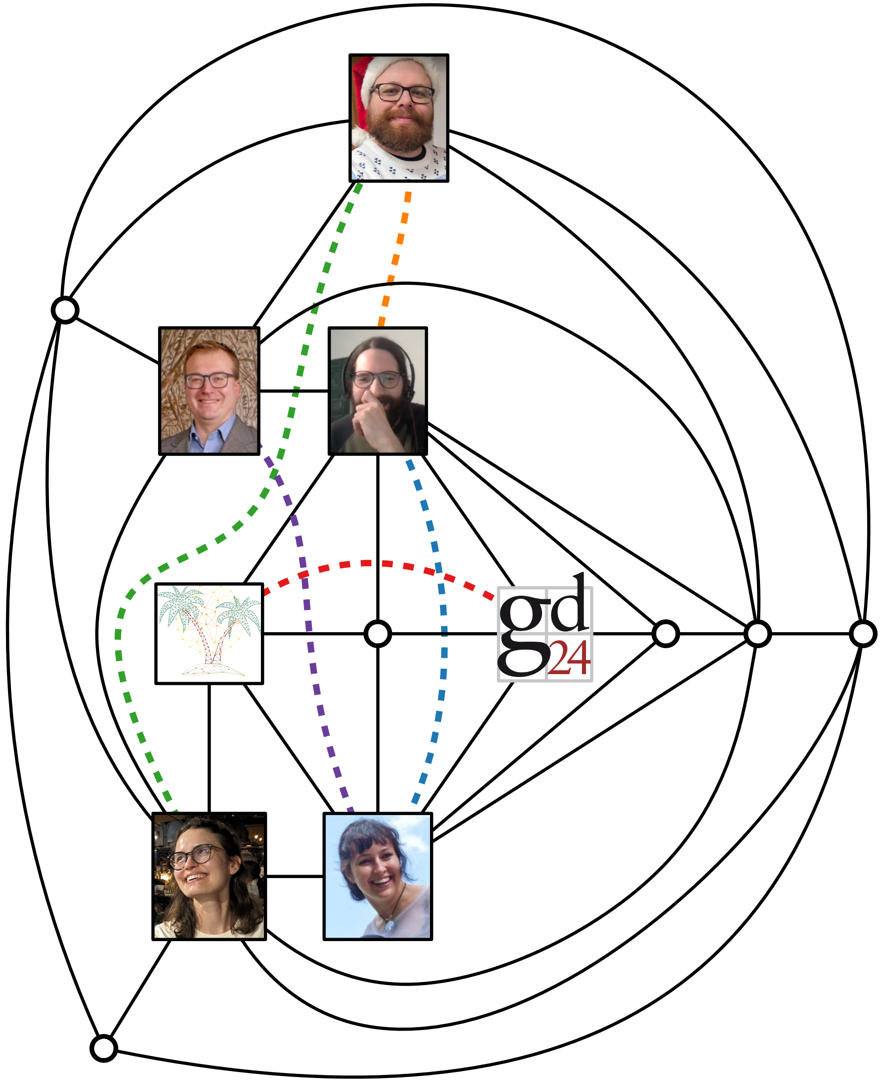
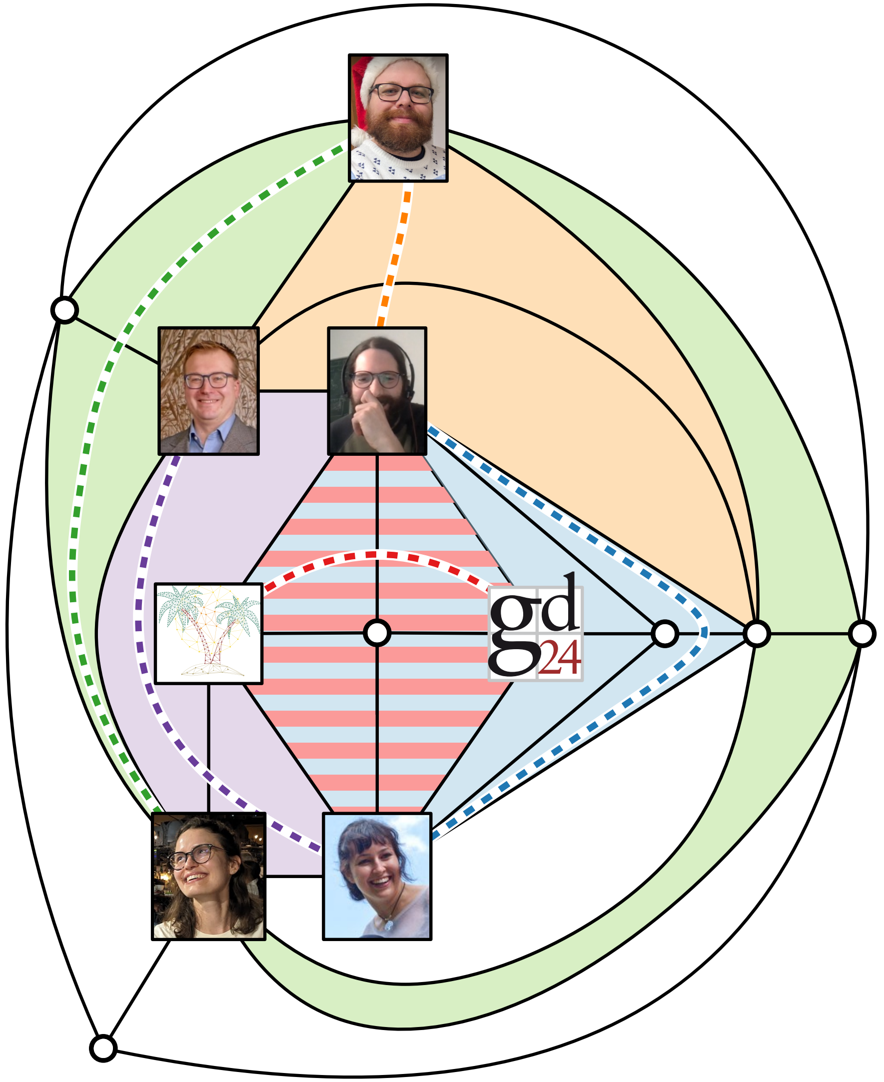
Ignaz Rutter

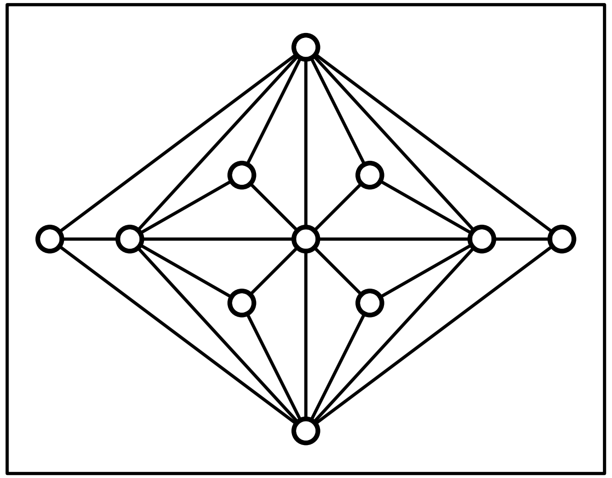




SWGD 2023
Summer Workshop
on Graph Drawing

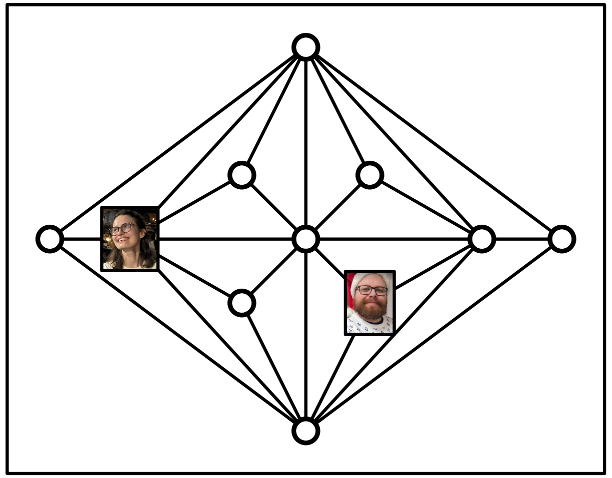
GD 2024
Graph Drawing &
Network Vis.

Inserting an Edge Into a Planar Graph



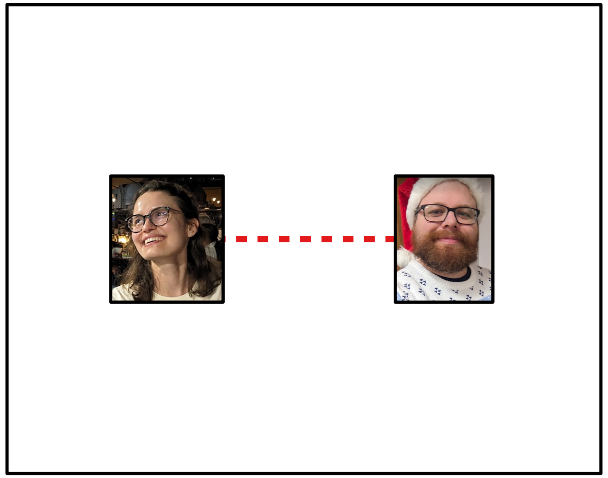
planar graph G

Inserting an Edge Into a Planar Graph



planar graph G

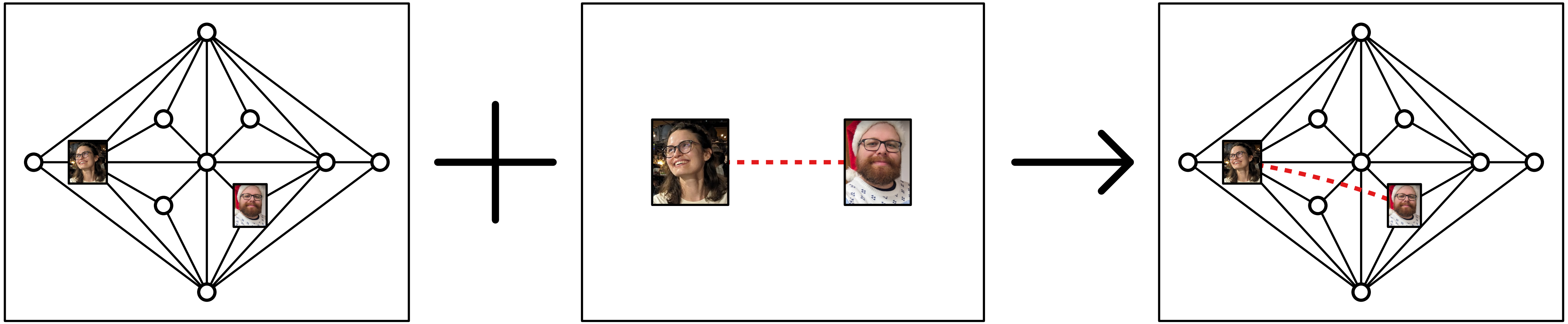
+



an edge e btw.
2 vertices of G

Inserting an Edge Into a Planar Graph

2 crossings

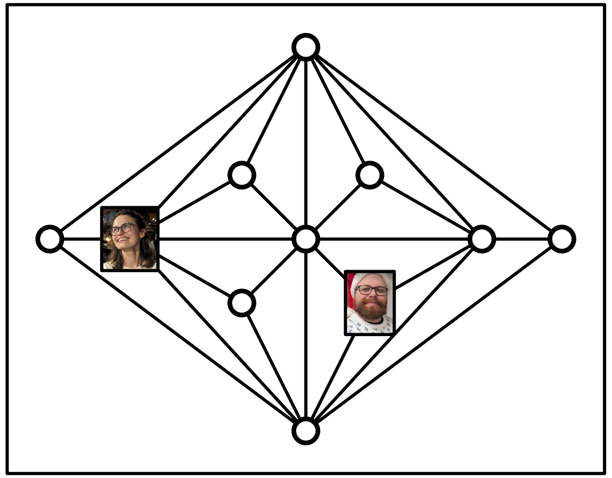


planar graph G

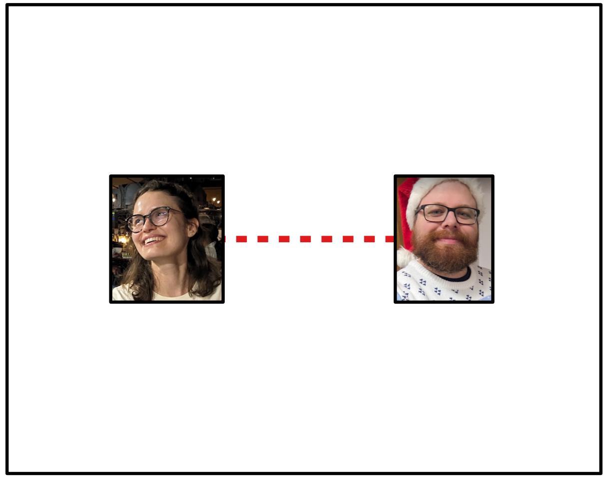
an edge e btw.
2 vertices of G

crossing-min. drawing of $G + e$

Inserting an Edge Into a Planar Graph



planar graph G



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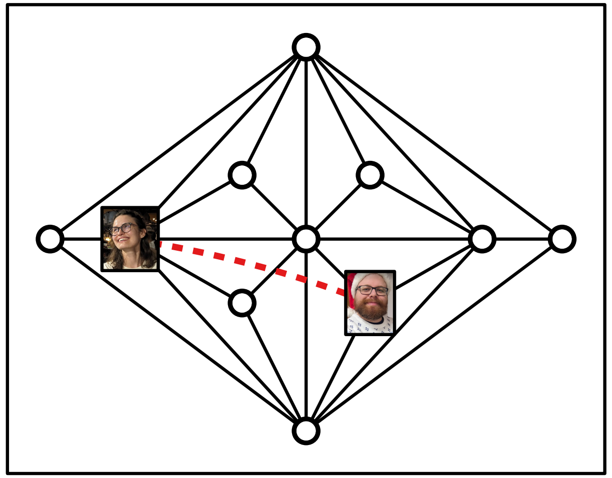
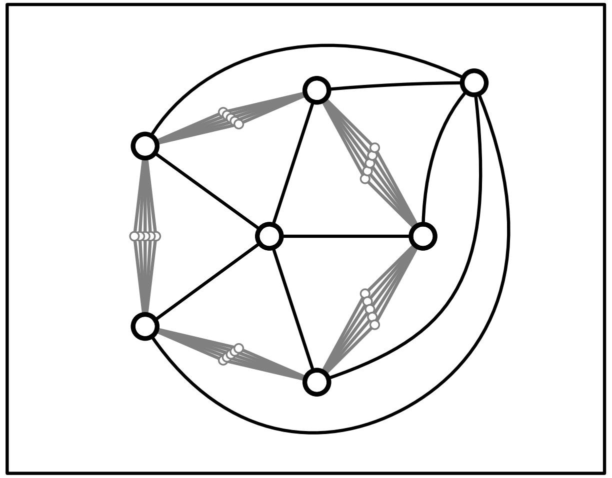
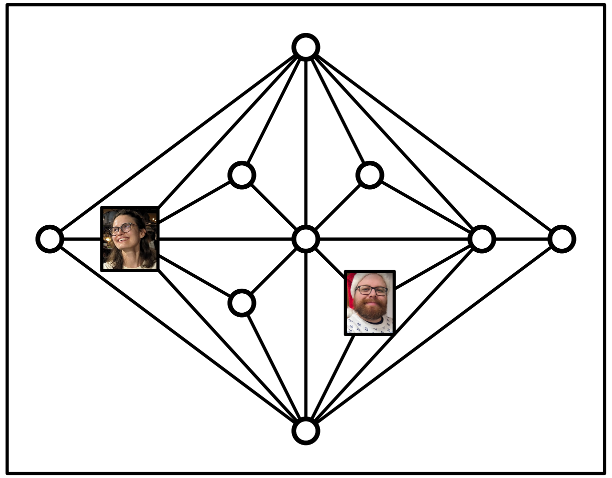


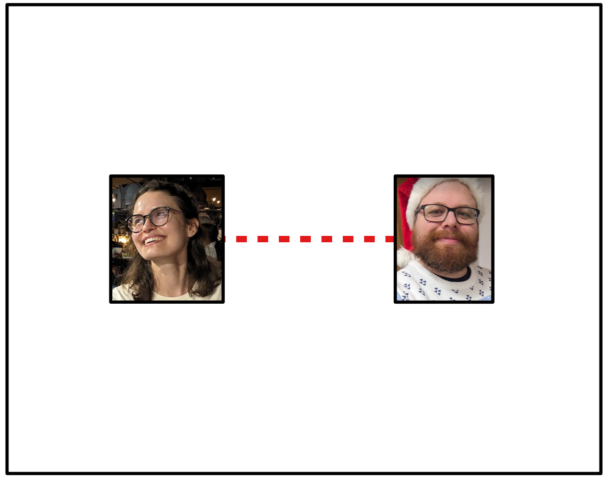
crossing-min. drawing of $G + e$



Inserting an Edge Into a Planar Graph



planar graph G



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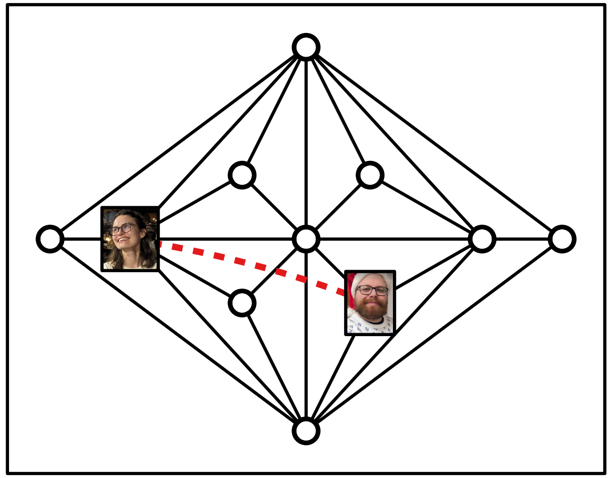
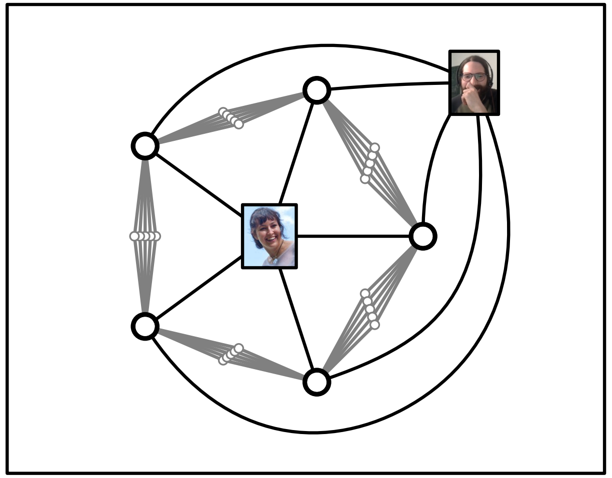
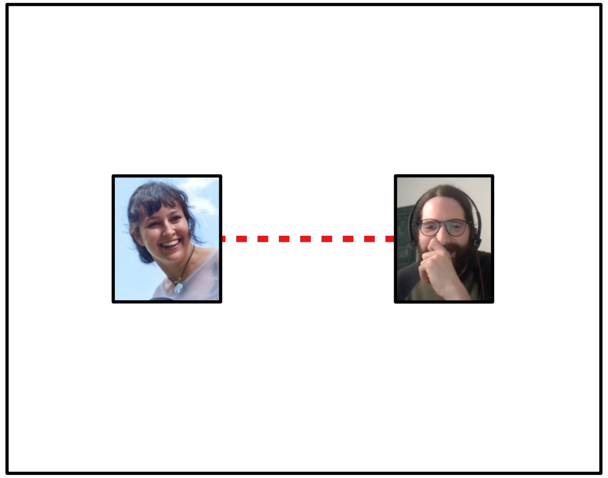


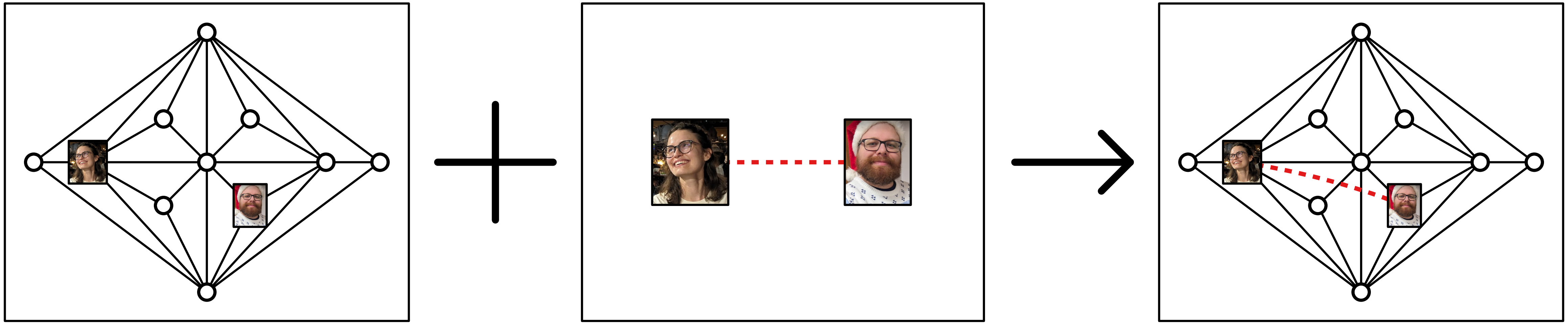
crossing-min. drawing of $G + e$





Inserting an Edge Into a Planar Graph

2 crossings



planar graph G

an edge e btw.
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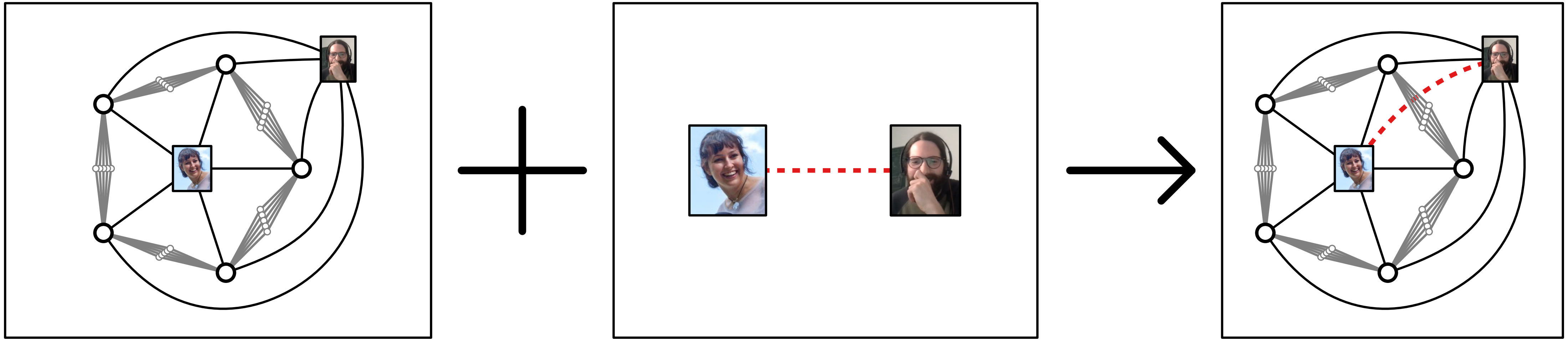
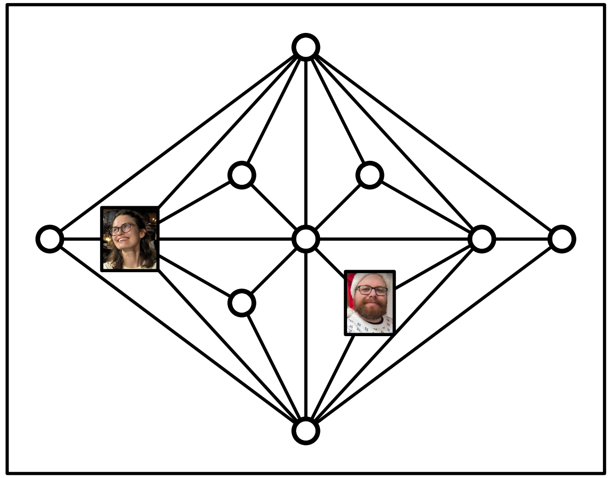
crossing-min. drawing of $G + e$

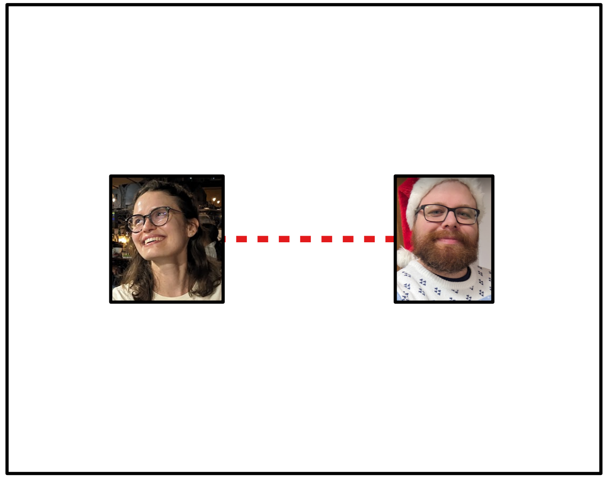


k crossings

Inserting an Edge Into a Planar Graph



planar graph G



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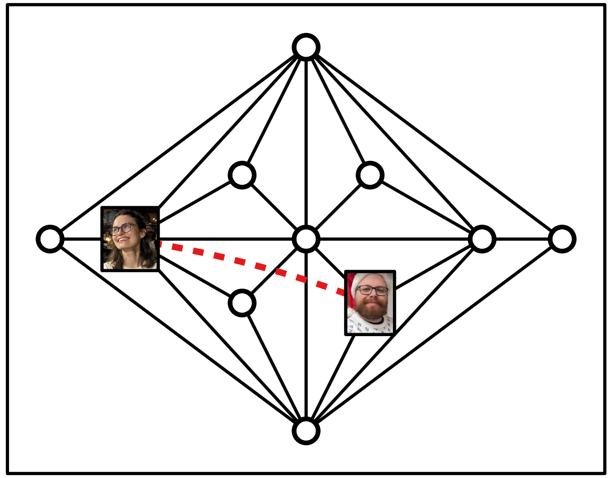
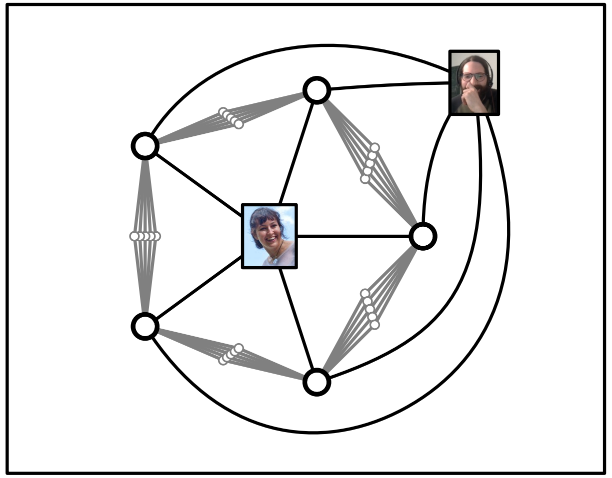
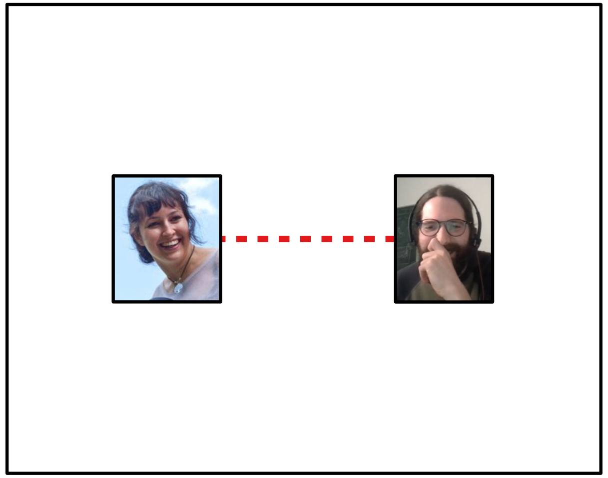
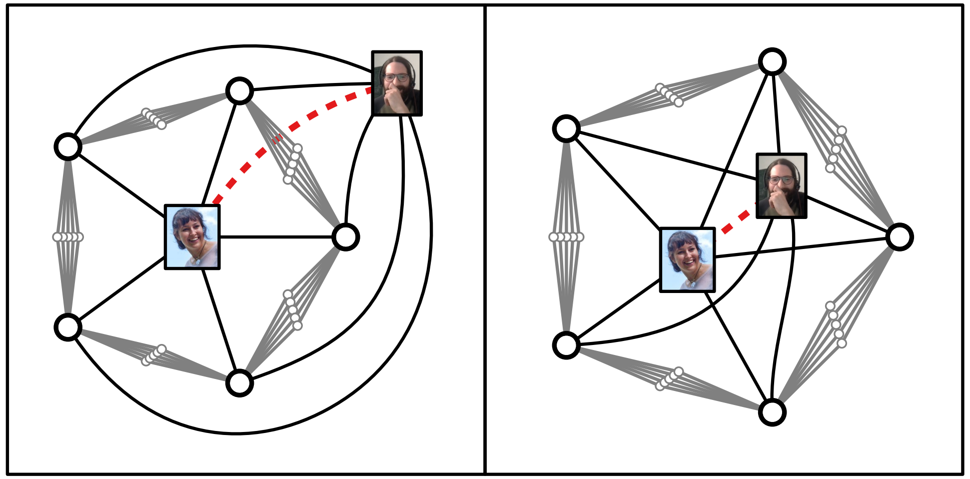
2 crossings



crossing-min. drawing of $G + e$





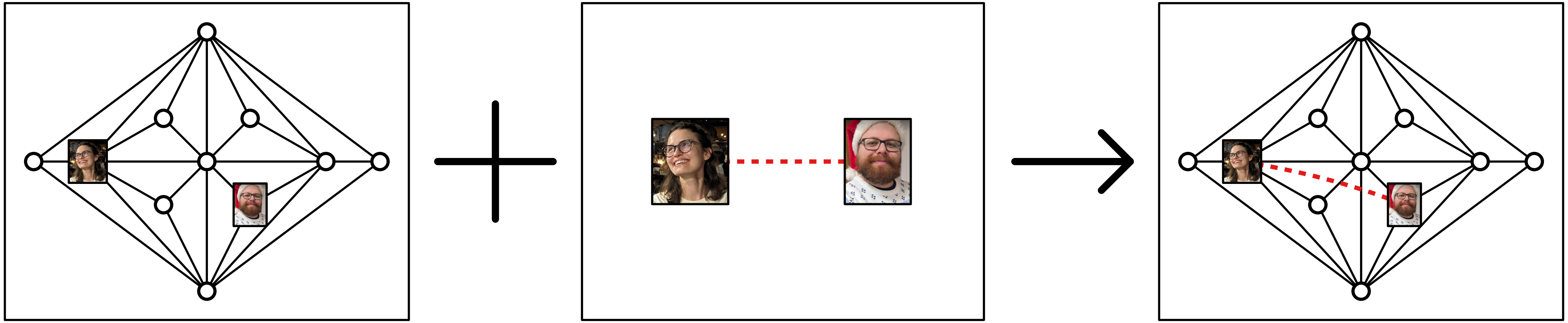


k crossings

4 crossings

Inserting an Edge Into a Planar Graph

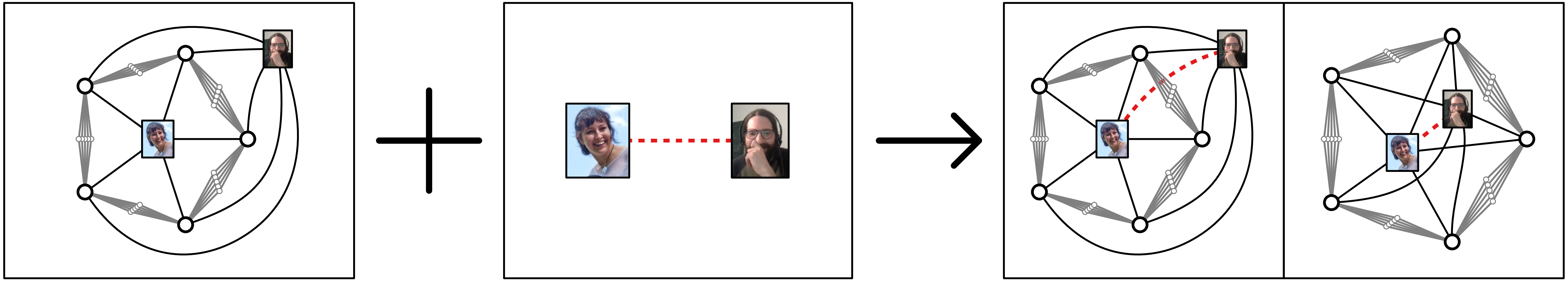
2 crossings



planar graph G

an edge e btw. 2 vertices of G

crossing-min. drawing of $G + e$

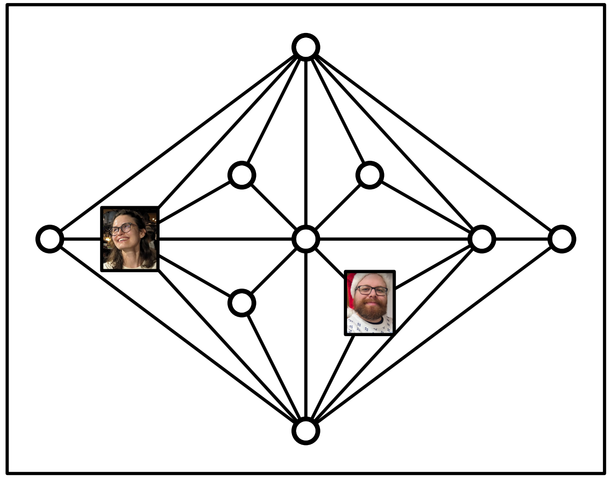


k crossings

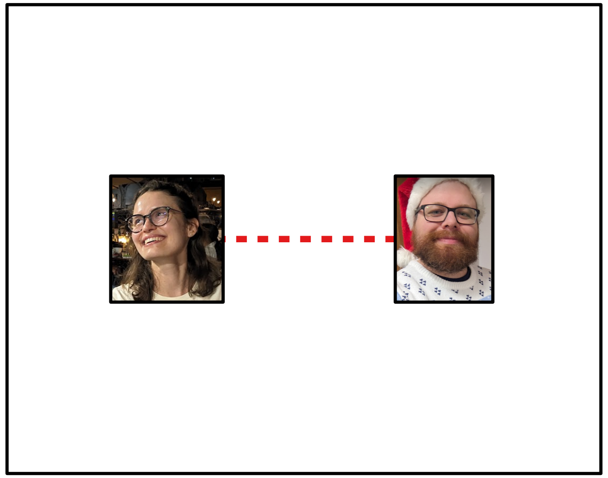
4 crossings

This problem is NP-hard. [Cabello & Mohar '08]

Inserting an Edge Into a Planar Graph II

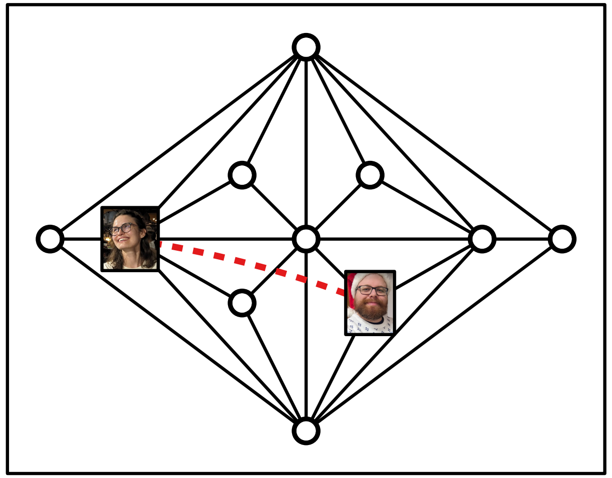


planar graph G

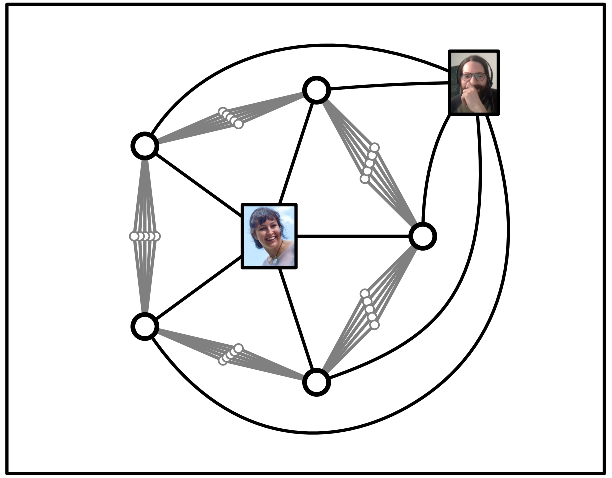


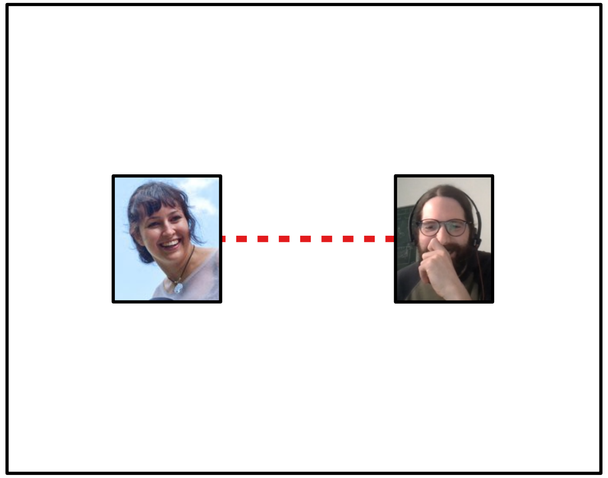
an edge e btw.
2 vertices of G

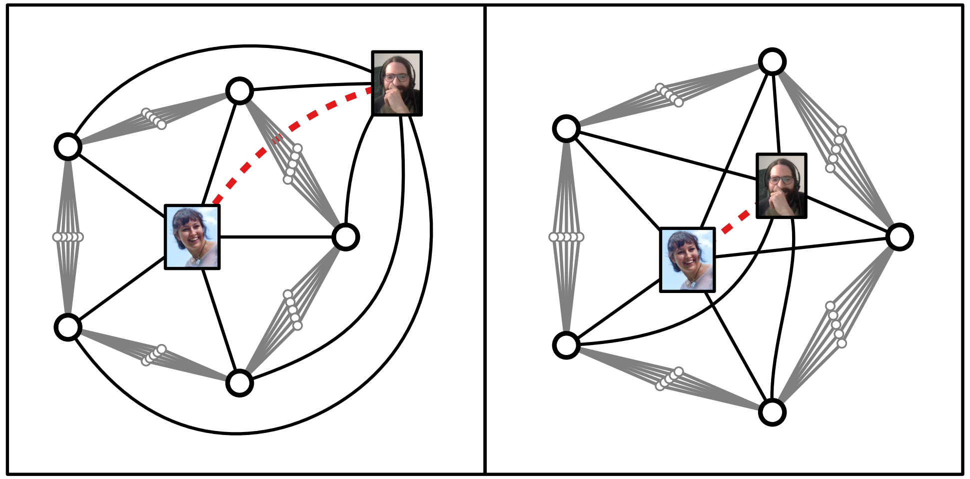
2 crossings



crossing-min. drawing of $G + e$
s.t. G is drawn planar

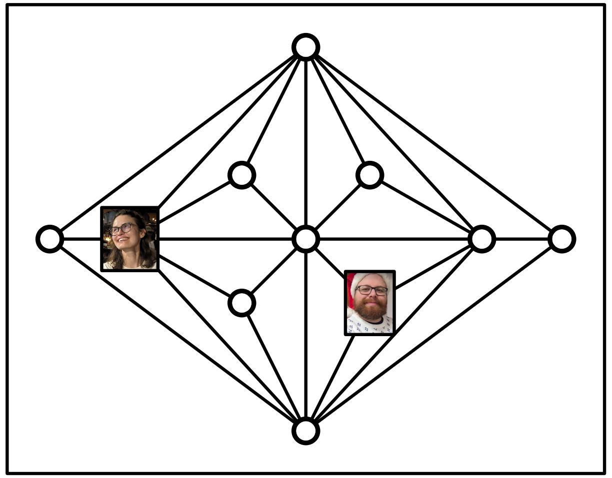






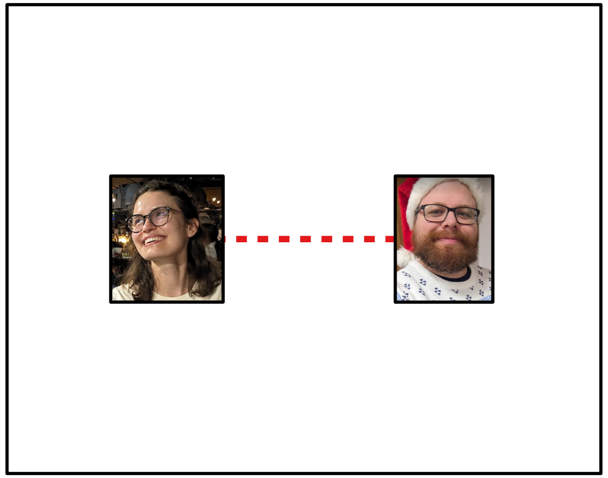
k crossings   4 crossings

Inserting an Edge Into a Planar Graph II



planar graph G

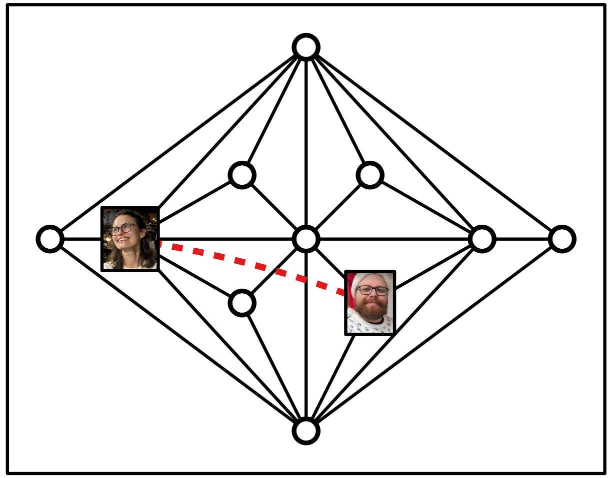
+



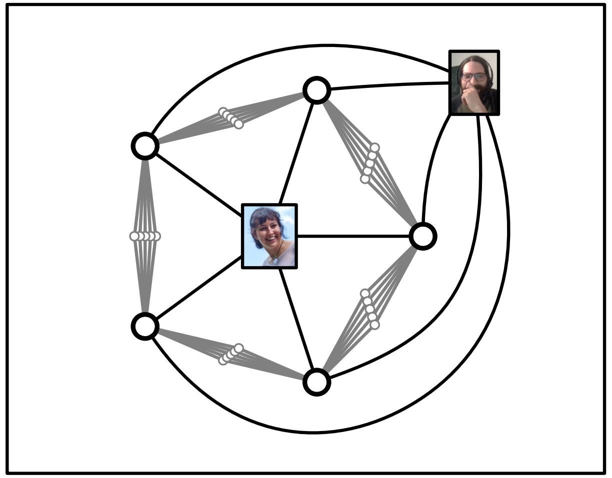
an edge e btw.
2 vertices of G

→

2 crossings



crossing-min. drawing of $G + e$
s.t. G is drawn planar



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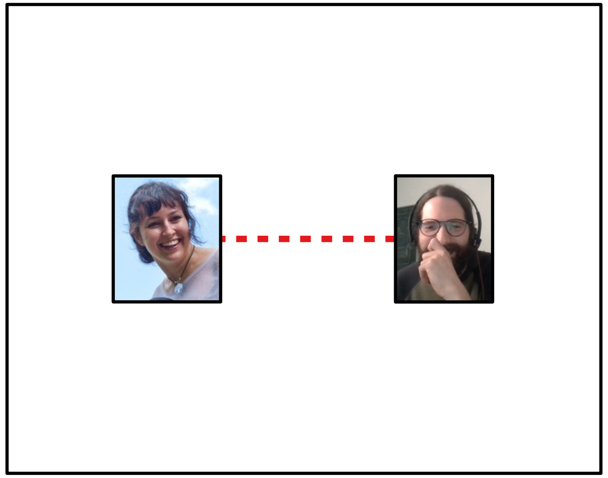

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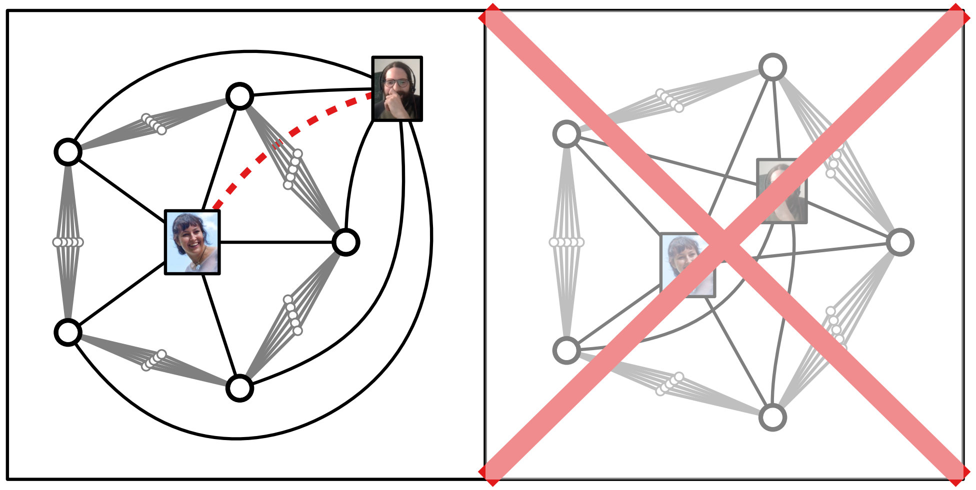
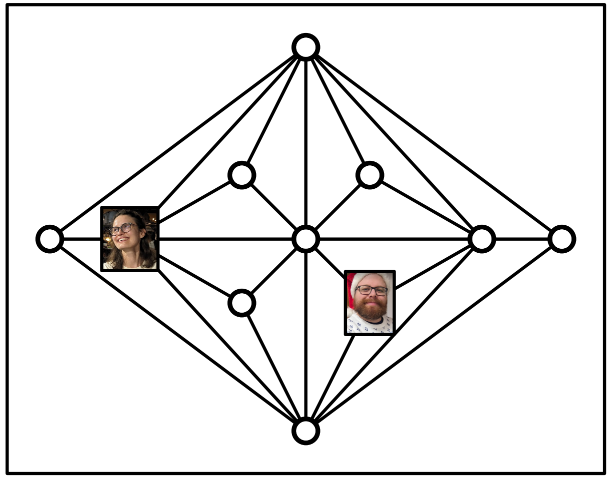


k crossings

4 crossings

Inserting an Edge Into a Planar Graph II



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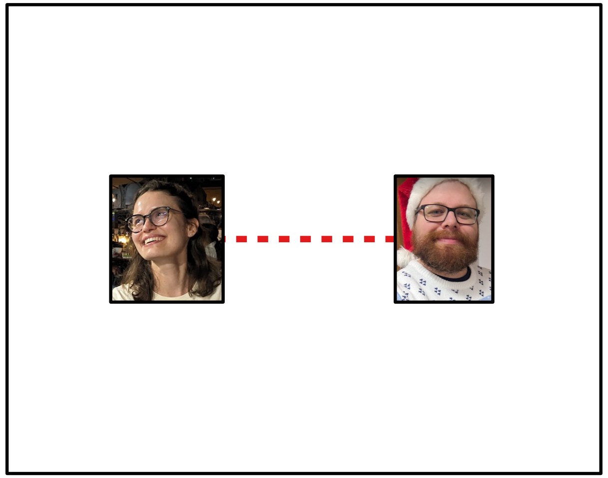

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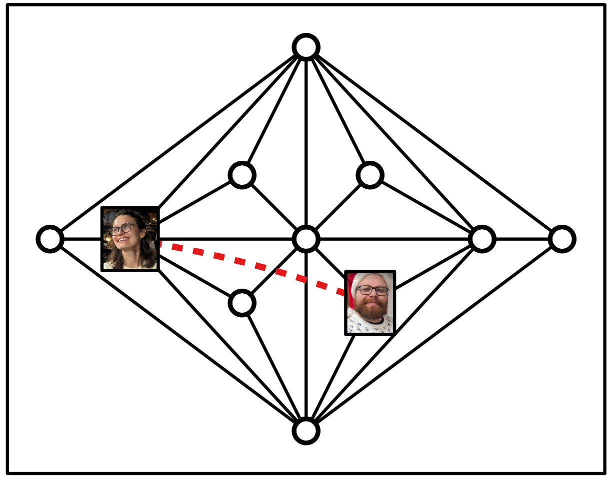
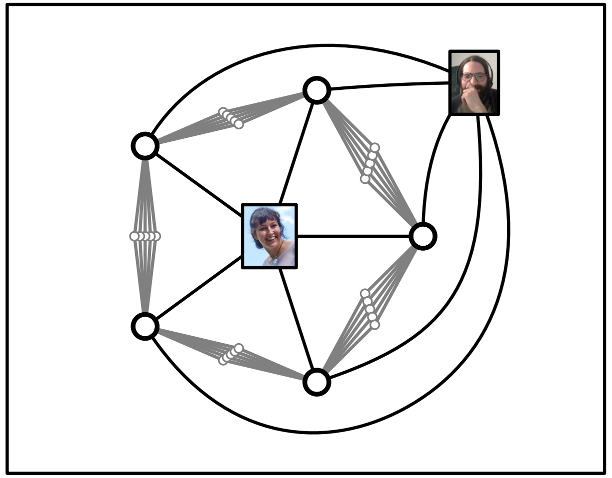
2 crossings



planar graph G

an edge e btw.
2 vertices of G

crossing-min. drawing of $G + e$
s.t. G is drawn planar



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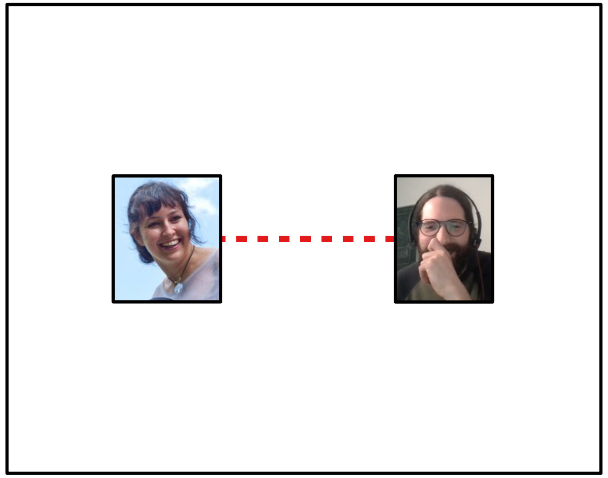

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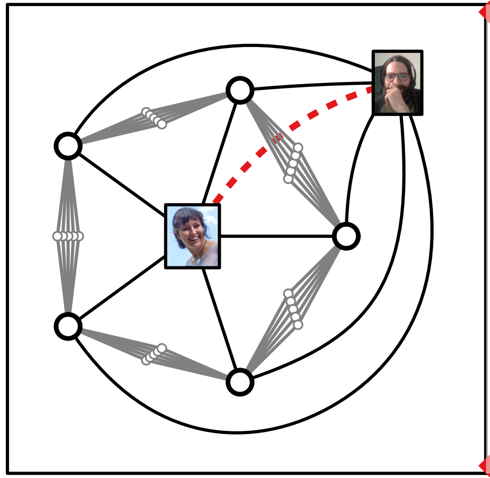
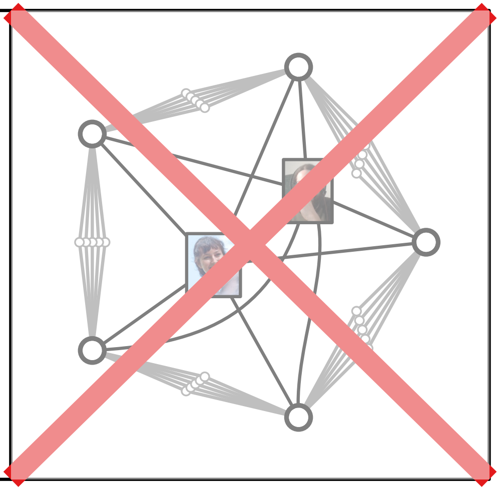


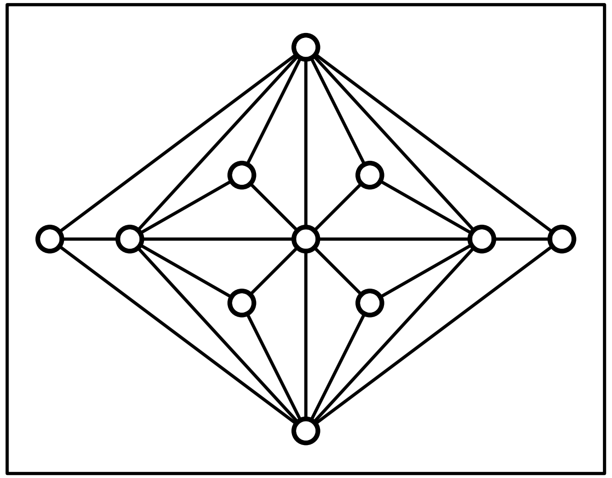
k crossings



4 crossings

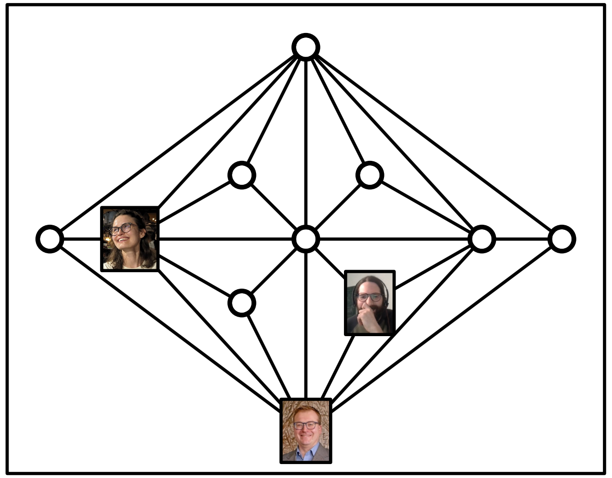
This problem can be solved in $\mathcal{O}(n)$ time. [Gutwenger, Mutzel & Weiskircher '05]

Inserting a **Vertex** Into a Planar Graph

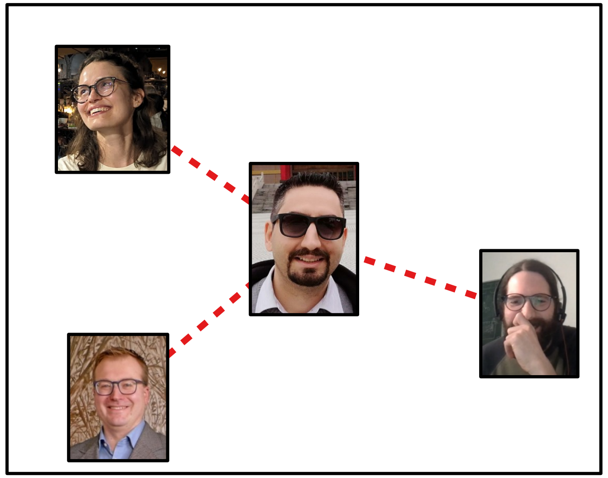


planar graph G

Inserting a **Vertex** Into a Planar Graph

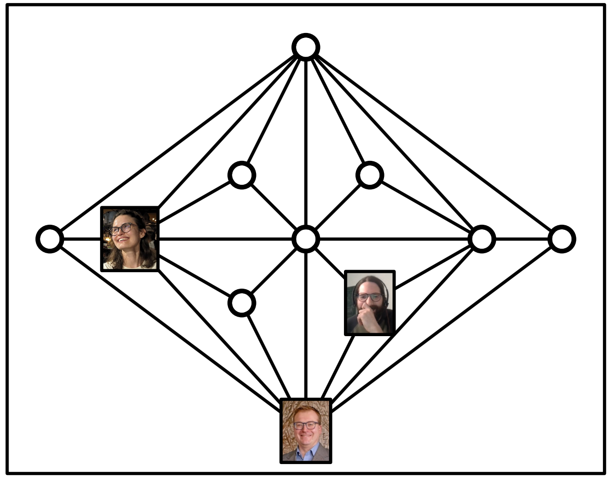


planar graph G

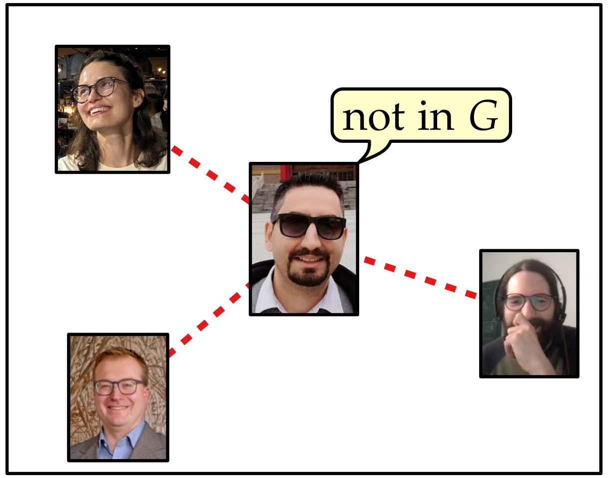


a star S with
its leaves in G

Inserting a **Vertex** Into a Planar Graph



planar graph G



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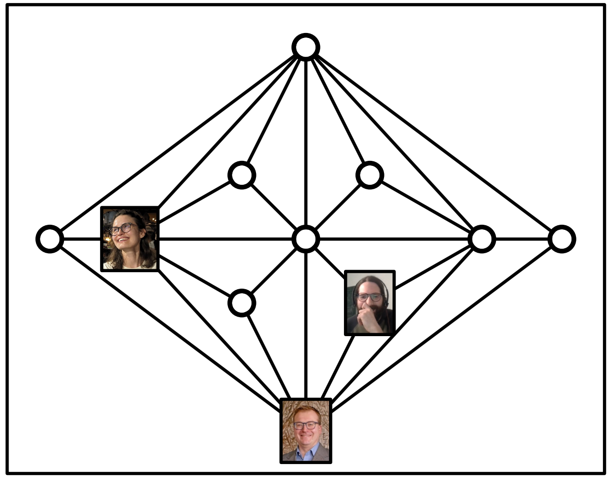
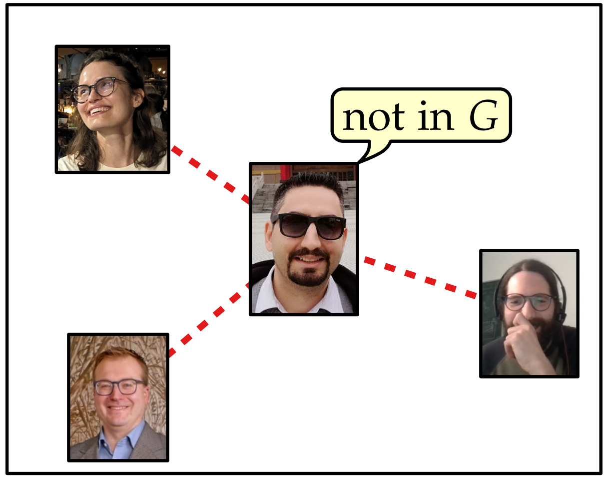


planar graph G



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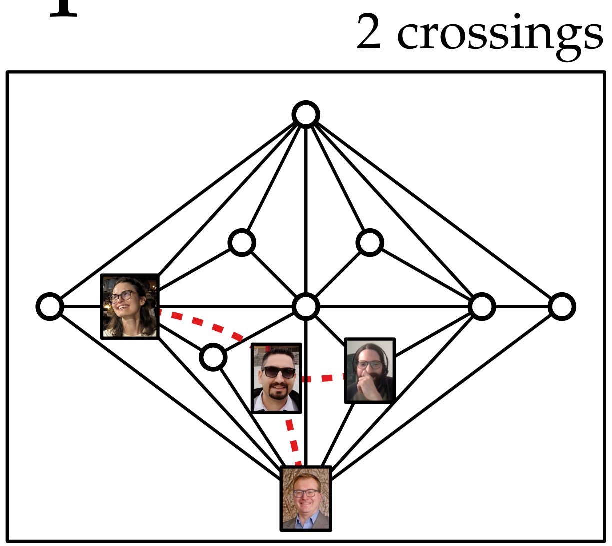
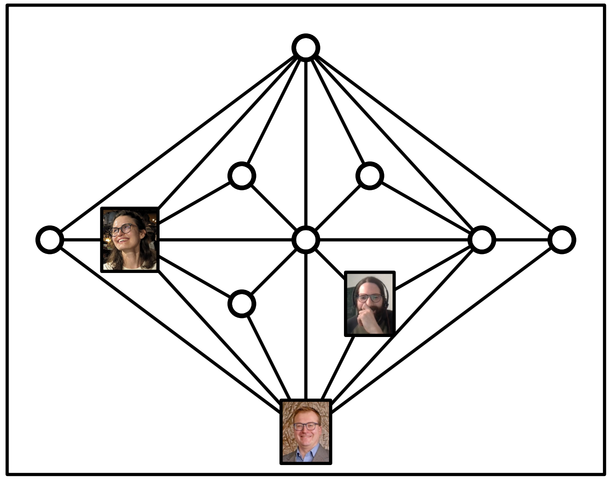
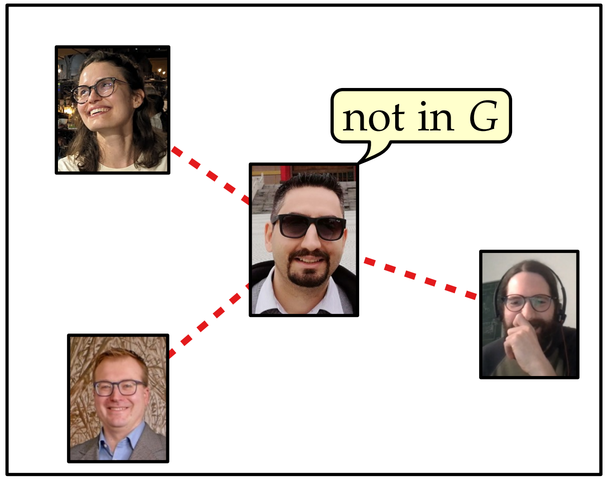


crossing-min. drawing of $G + S$
s.t. G is drawn planar

Inserting a **Vertex** Into a Planar Graph

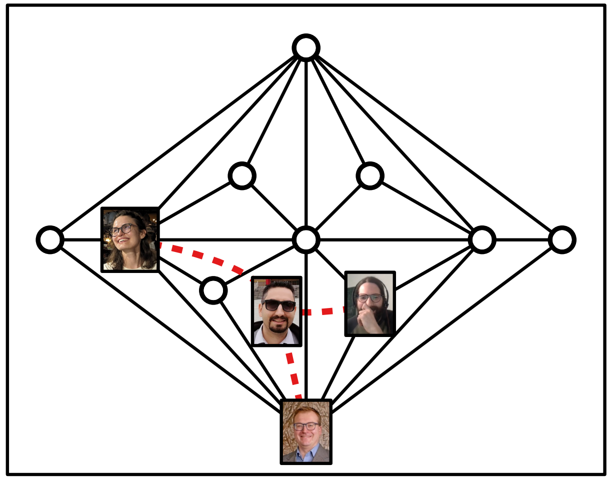


planar graph G



a star S with
its leaves in G

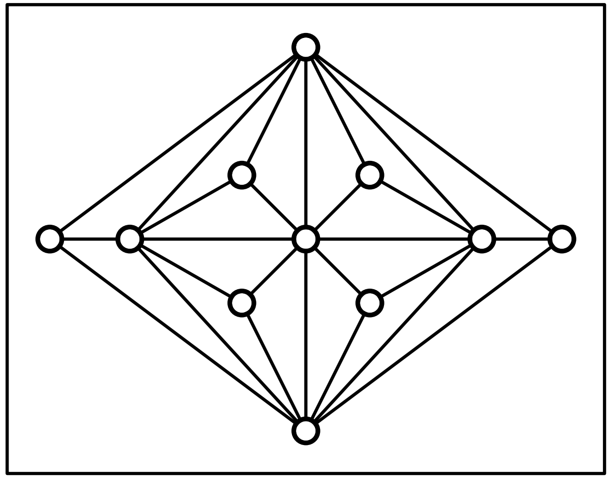
2 crossings



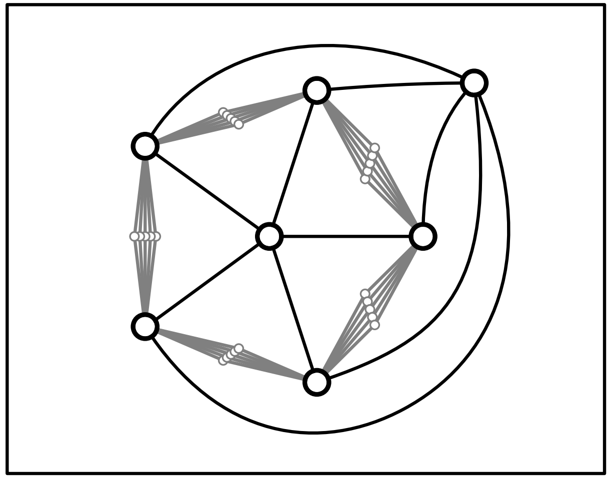
crossing-min. drawing of $G + S$
s.t. G is drawn planar

This problem can be solved in $\mathcal{O}(n^7)$ time. [Chimani, Gutwenger, Mutzel & Wolf '05]

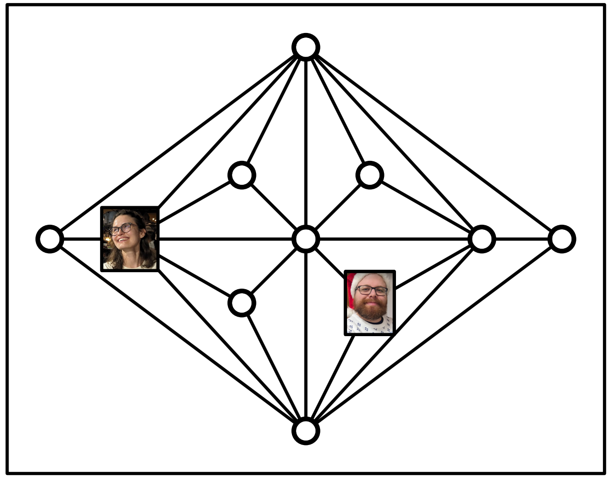
Inserting an Edge Into a **Plane** Graph



plane graph G
(planar graph
+ planar embedding)



Inserting an Edge Into a **Plane** Graph



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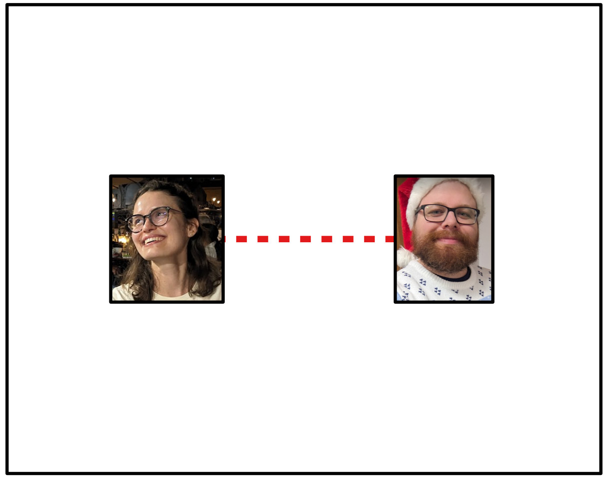
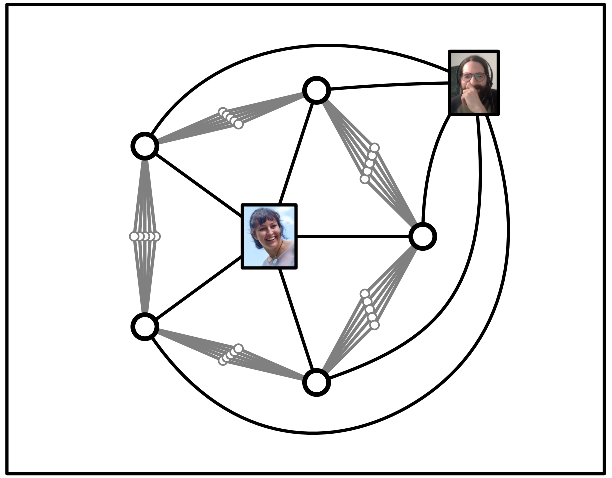


plane graph G
(planar graph
+ planar embedding)

an edge e btw.
2 vertices of G



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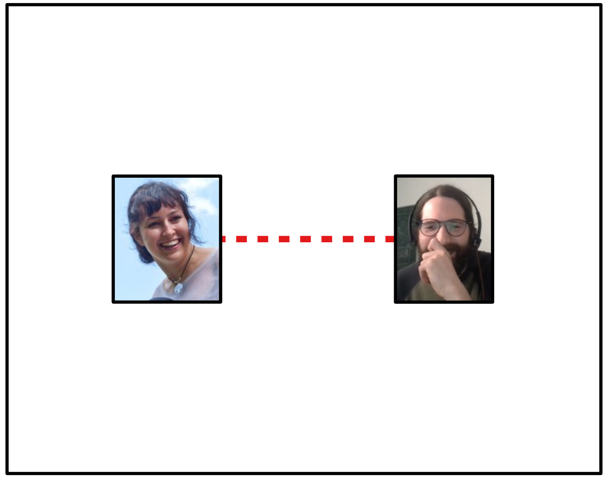
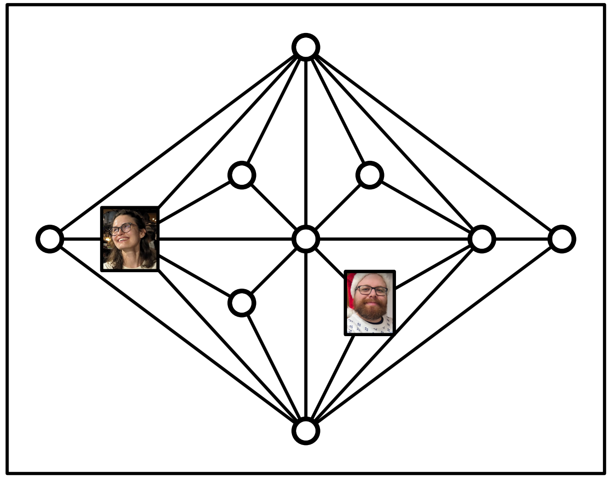


Inserting an Edge Into a **Plane** Graph

2 crossings



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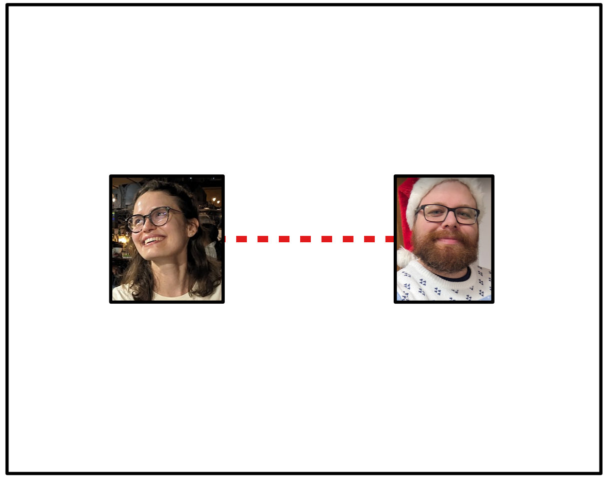

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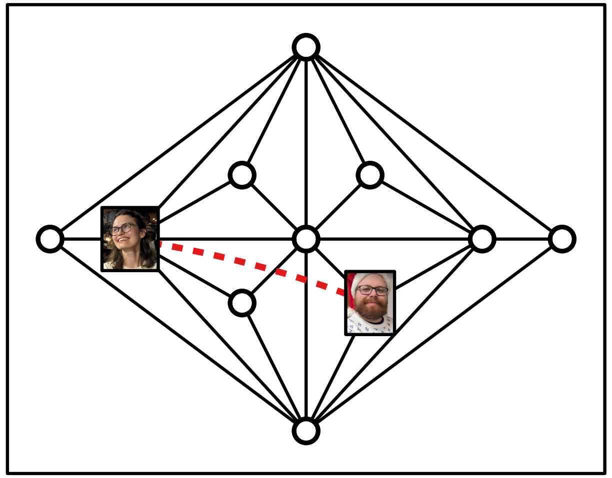
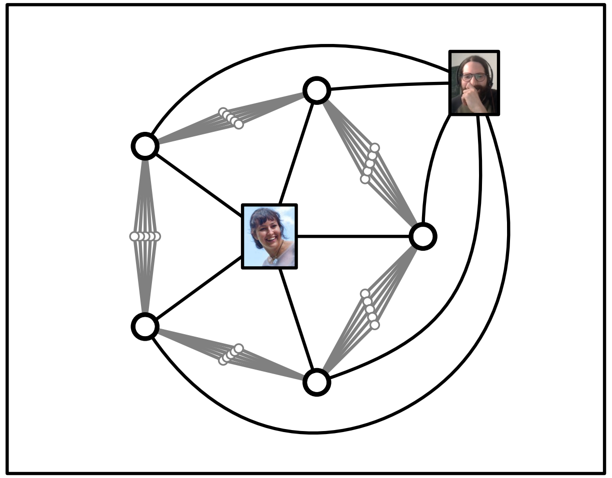


plane graph G (planar graph + planar embedding)

an edge e btw. 2 vertices of G

crossing-min. drawing of $G + e$ that keeps the embedding of G

k crossings



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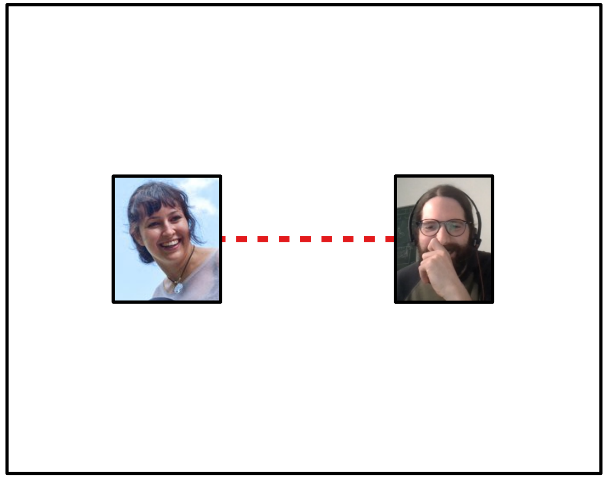

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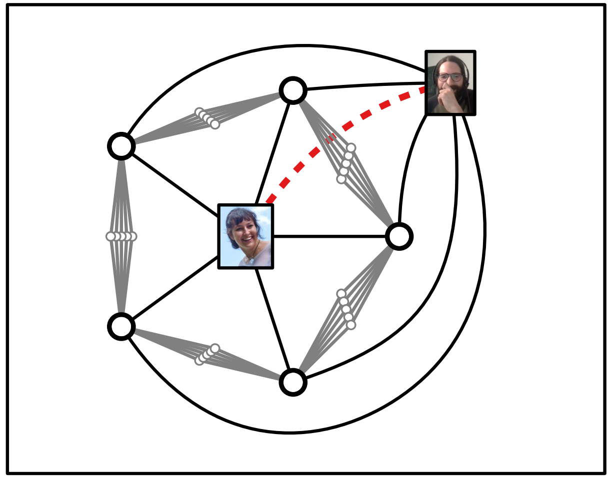
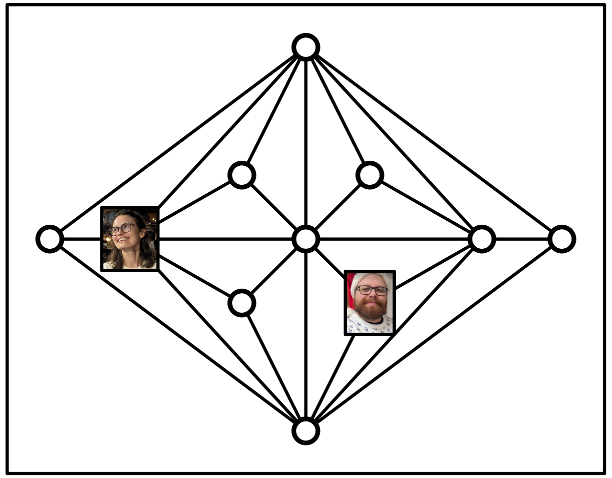
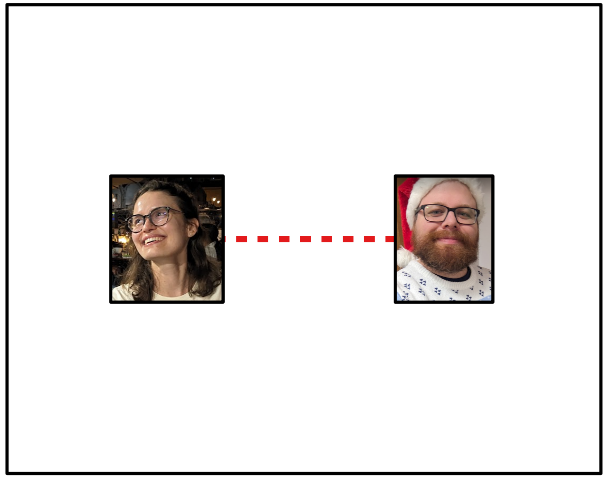
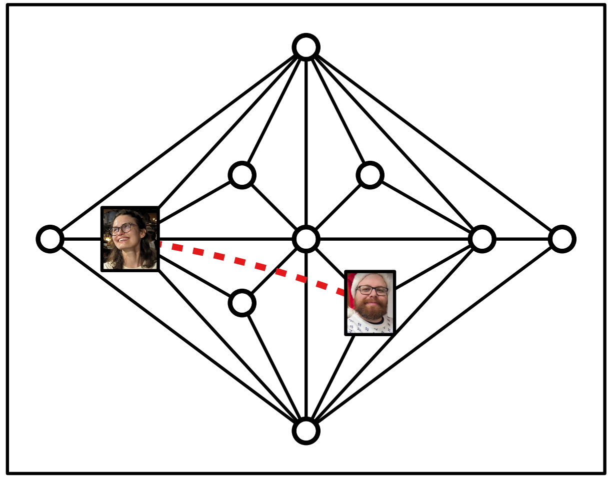


Inserting an Edge Into a **Plane** Graph





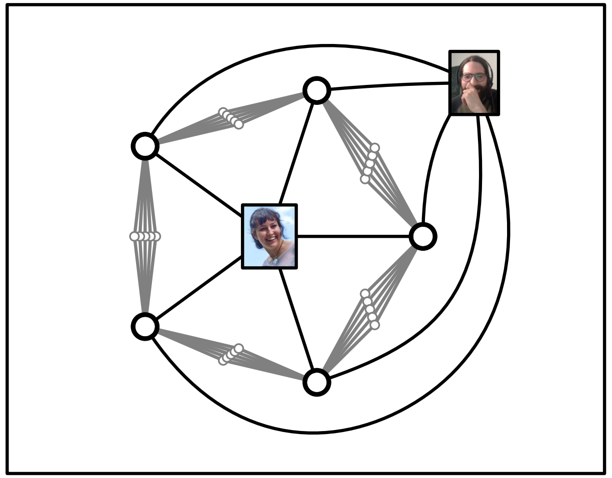
2 crossings

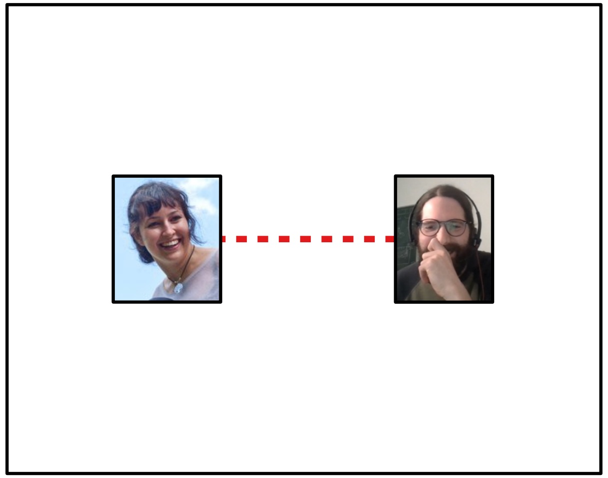


plane graph G
(planar graph
+ planar embedding)

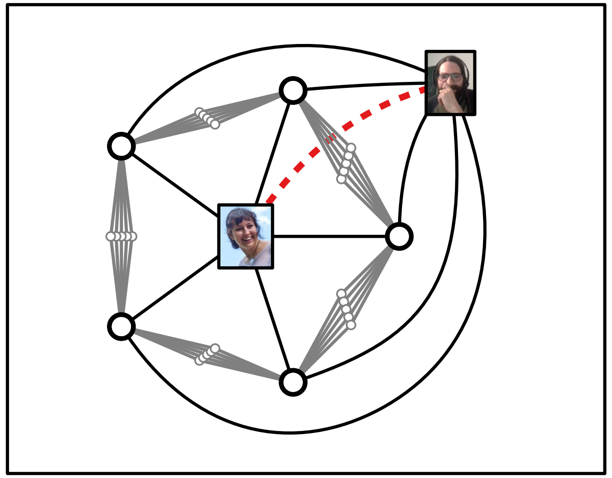
an edge e btw.
2 vertices of G

crossing-min. drawing of $G + e$
that keeps the embedding of G



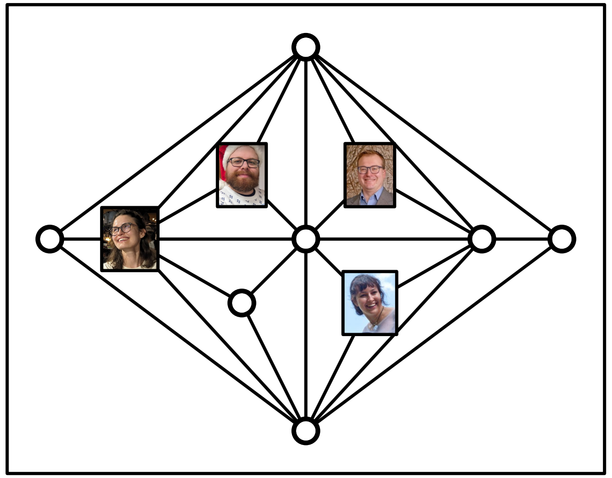


k crossings

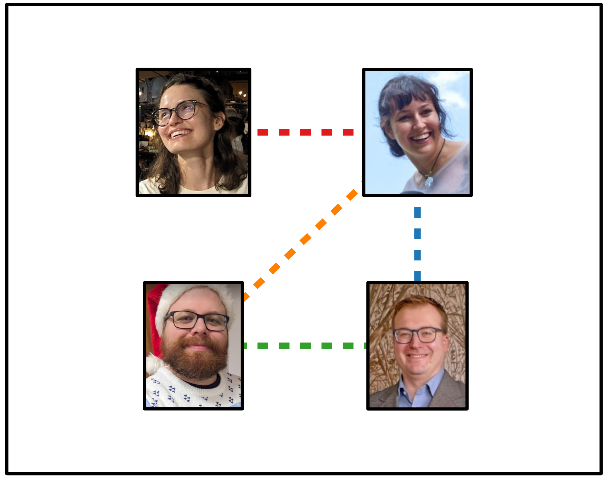


This problem can be solved in $\mathcal{O}(n)$ time. [BFS]

Inserting **Edges** Into a Plane Graph

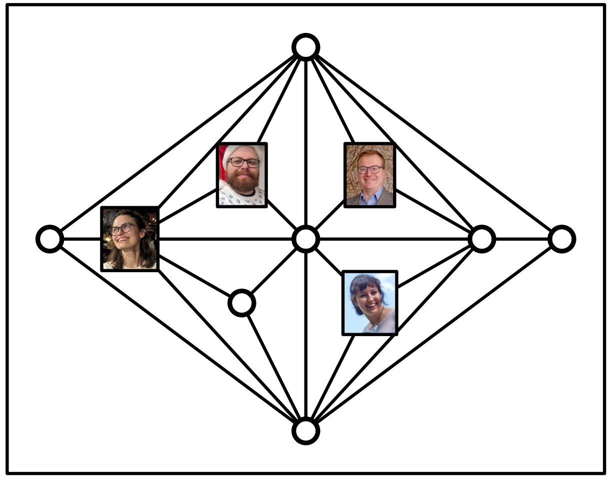


plane graph G



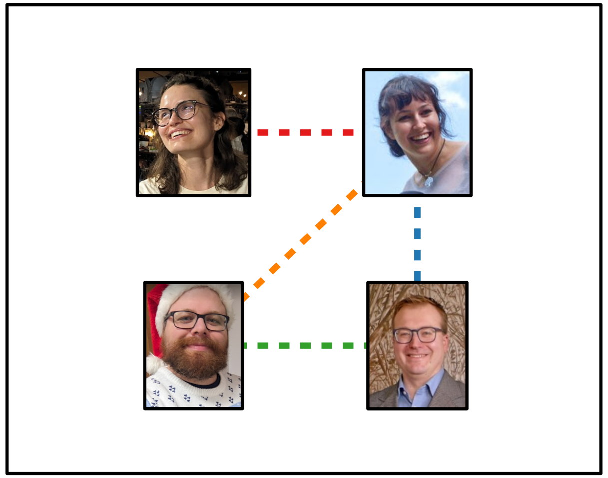
edges E' btw. vtcs in G

Inserting **Edges** Into a Plane Graph



plane graph G

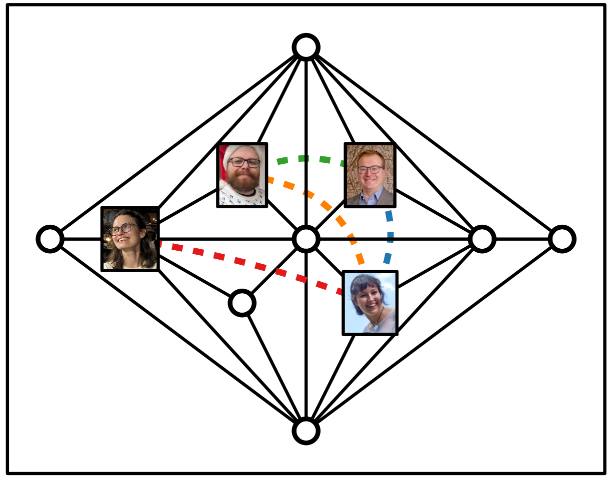
+



edges E' btw. vtcs in G

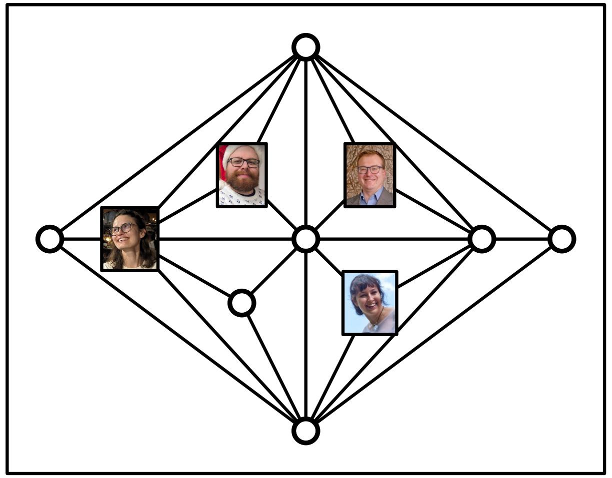
→

7 crossings

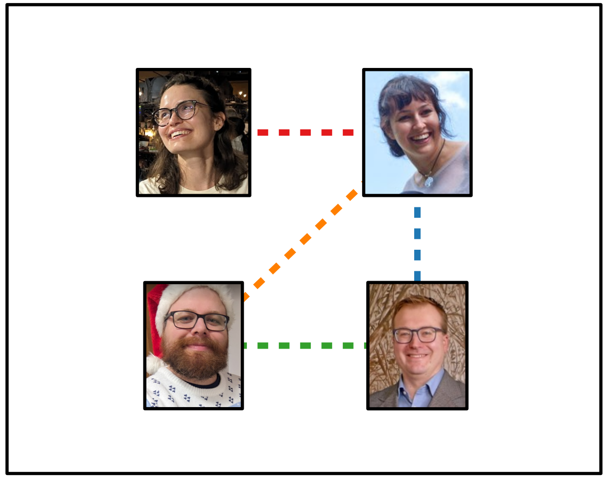


crossing-min. drawing of $G + E'$ that keeps the embedding of G

Inserting Edges Into a Plane Graph



plane graph G



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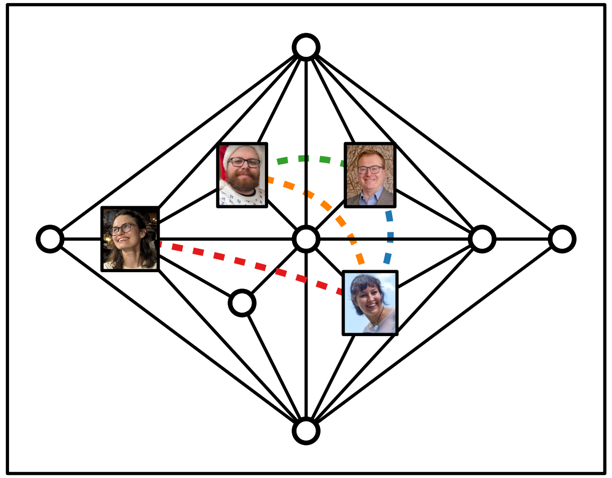
7 crossings



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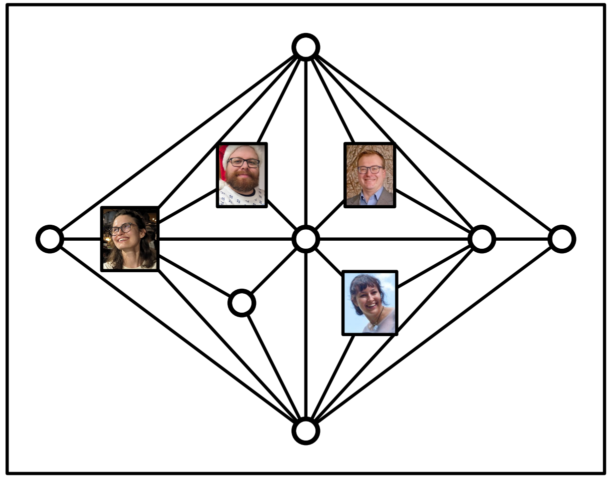
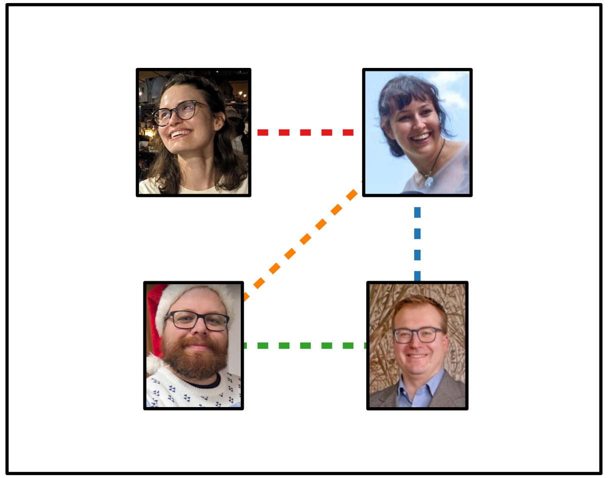
This problem is NP-hard. [Ziegler '01]

Inserting **Edges** Into a Plane Graph



plane graph G



edges E' btw. vtcs in G

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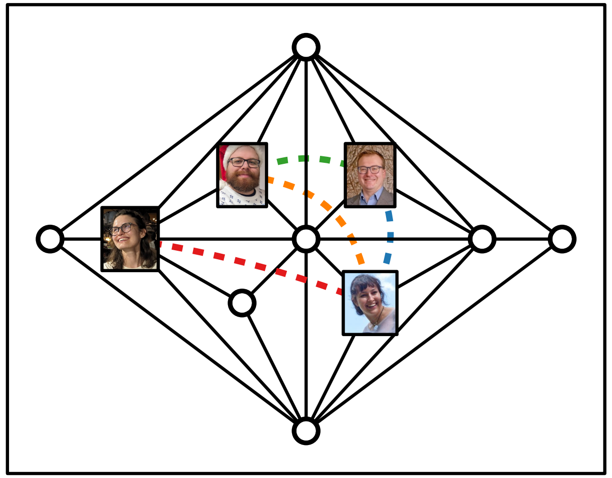

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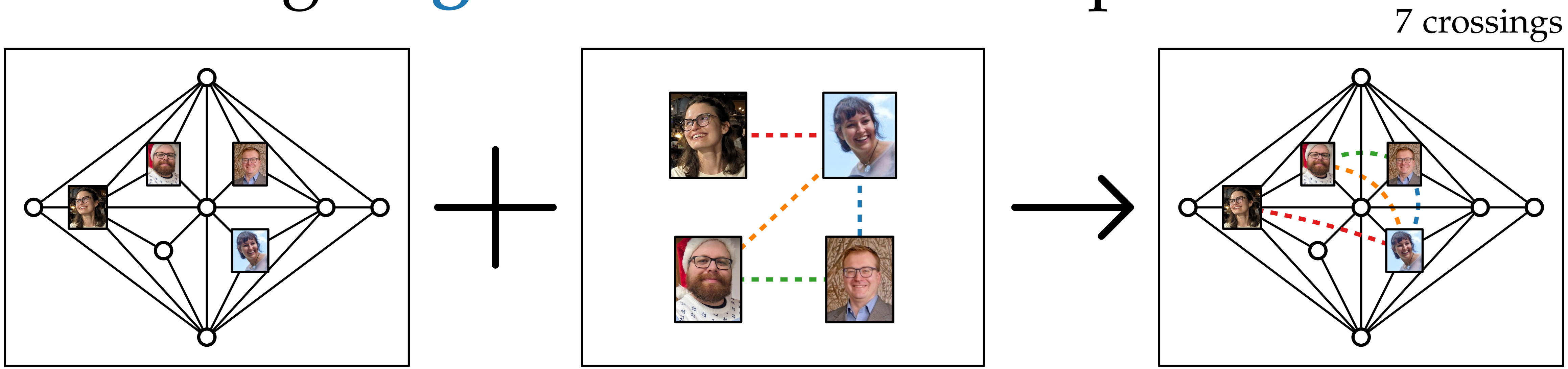
… even if G is biconnected.

Inserting **Edges** Into a Plane Graph



plane graph G | edges E' btw. vtcs in G | crossing-min. drawing of $G + E'$ that keeps the embedding of G

This problem is NP-hard. [Ziegler '01]
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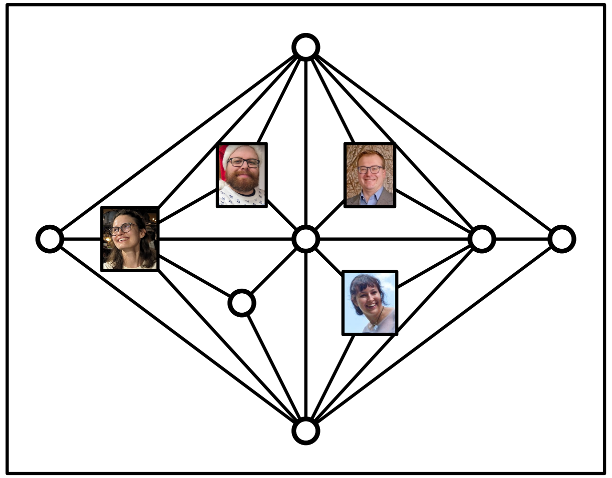
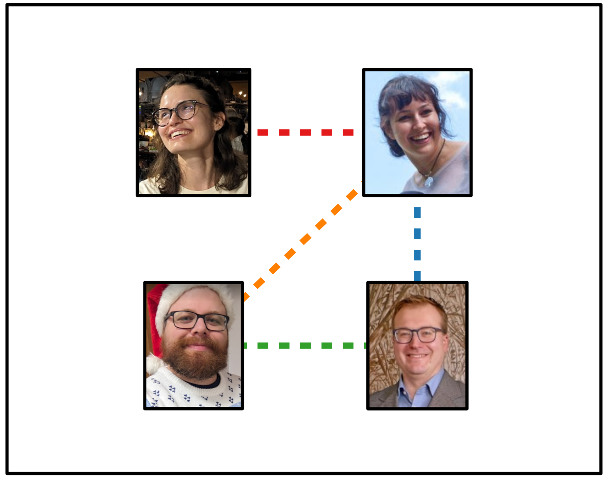
This problem is in FPT parameterized by #crossings. [Hamm & Hliněný '22]

Inserting **Edges** Into a Plane Graph

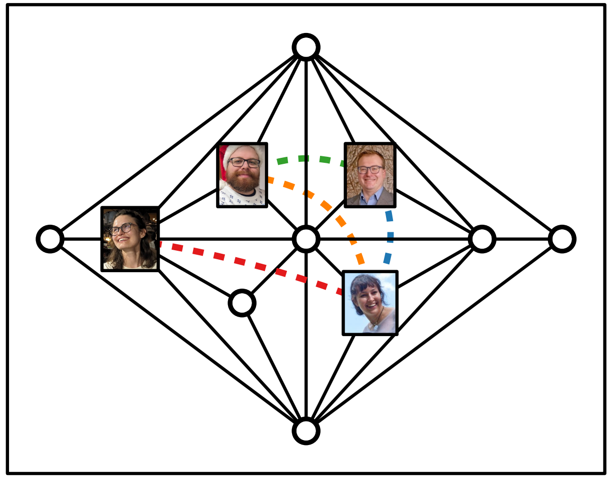


plane graph G



edges E' btw. vtcs in G

7 crossings



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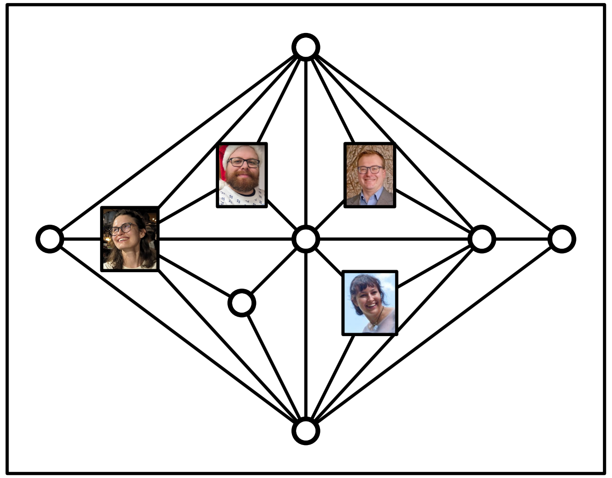
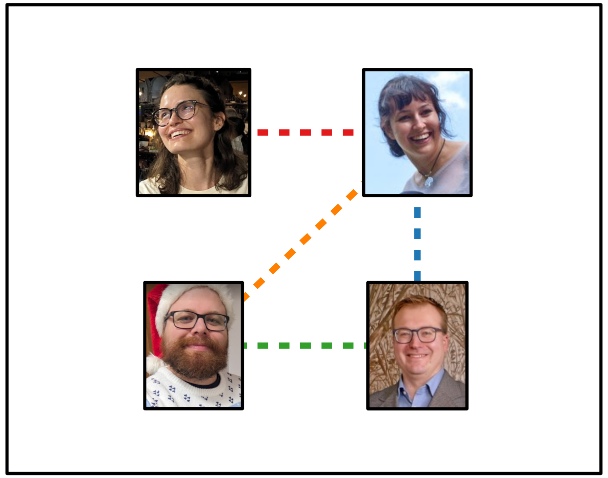
Inserting **Edges** Into a Plane Graph

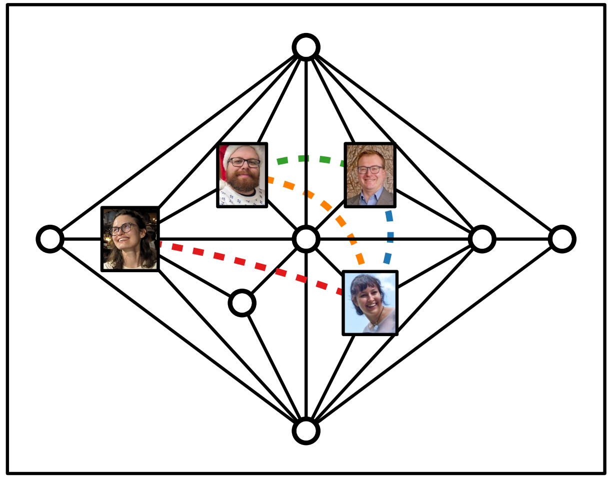


plane graph G



edges E' btw. vtcs in G

7 crossings



crossing-min. drawing of $G + E'$ that keeps the embedding of G

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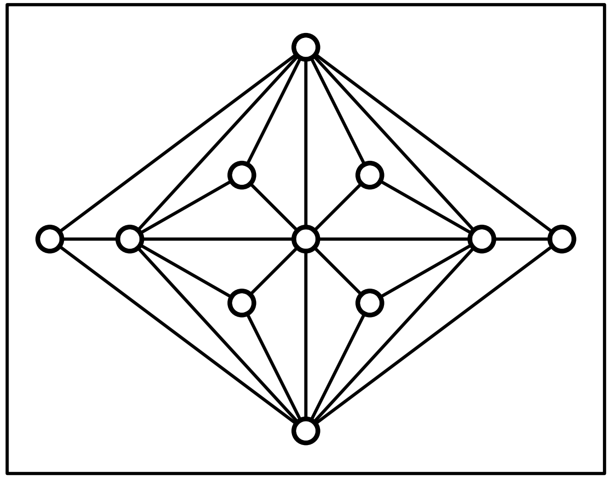
… even if G is non-planar (or drawn with crossings)

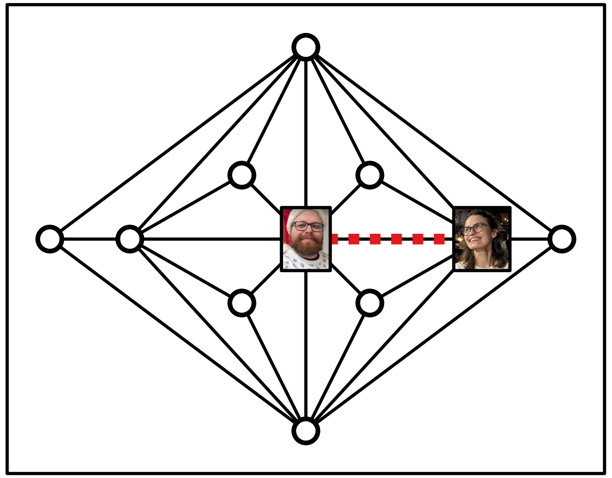
This problem is in FPT parameterized by $|E'|$ [Chimani & Hliněný '23]
if G is biconnected or all cutvertices have constant degree.

Partial Embedding – General Definition



graph G +
drawing style Φ
(e.g., straight-line planar)

Partial Embedding – General Definition



graph G +
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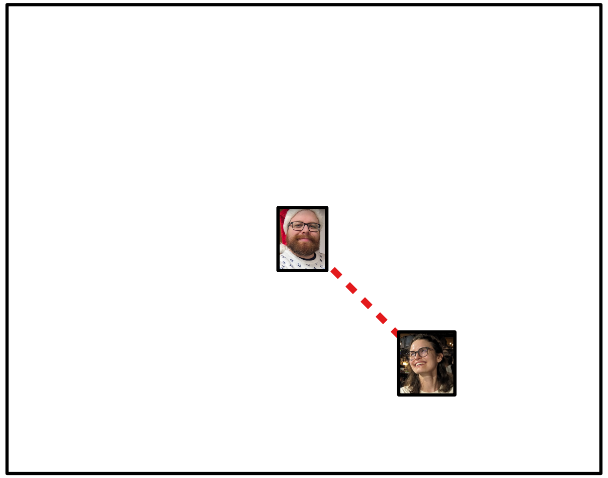
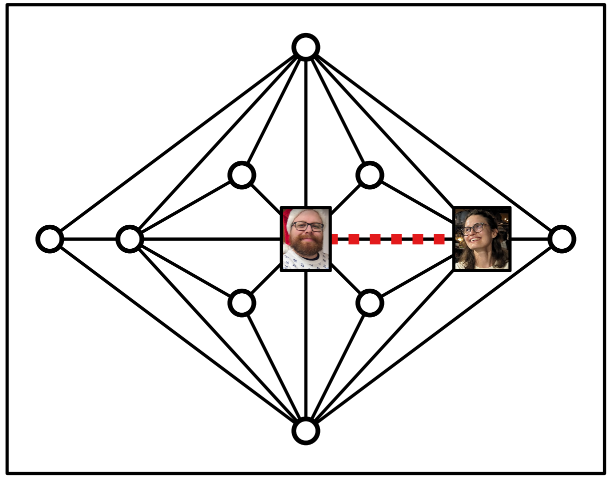
+



drawing with style Φ of
a subgraph $H \subseteq G$

Partial Embedding – General Definition



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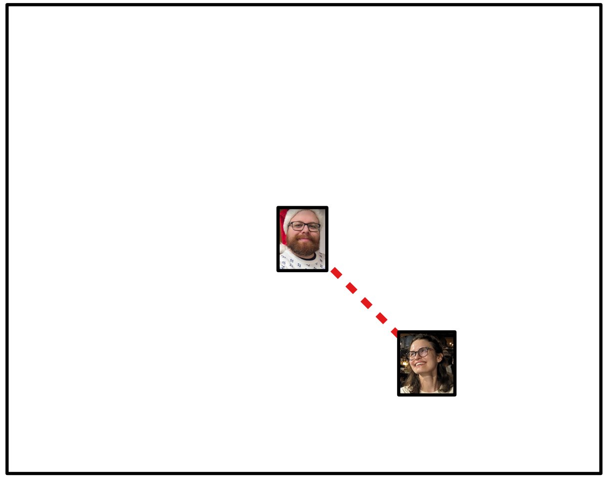

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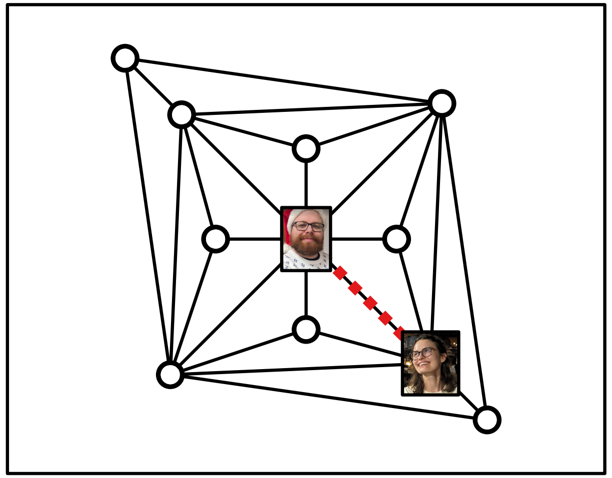
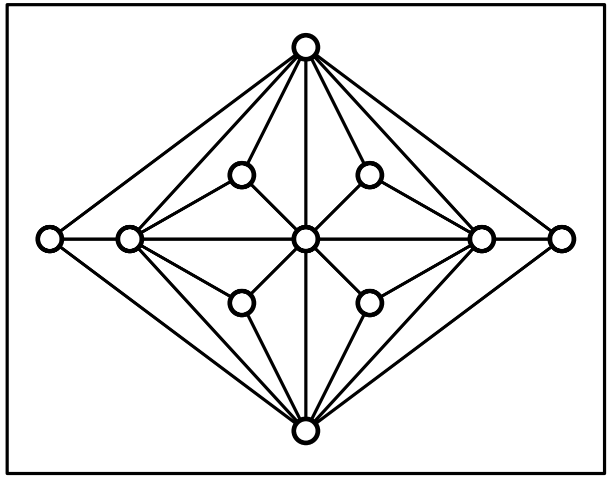


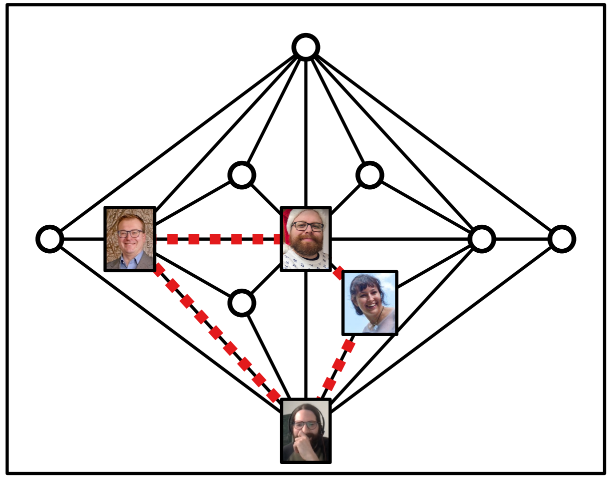
drawing with style Φ of G
s.t. H **keeps its drawing**

Partial Embedding of **Planar** Graphs



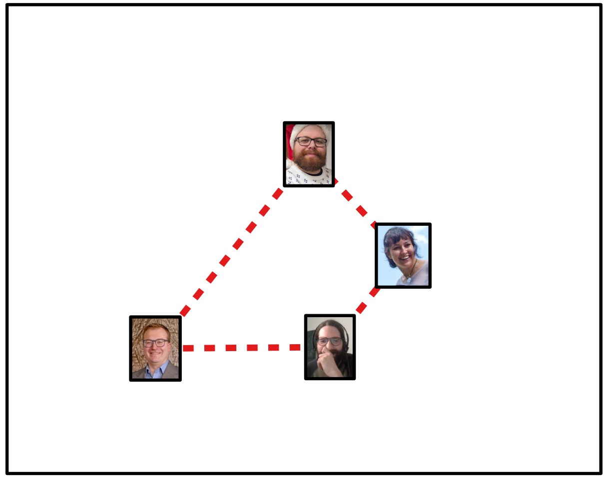
planar graph G

Partial Embedding of **Planar** Graphs



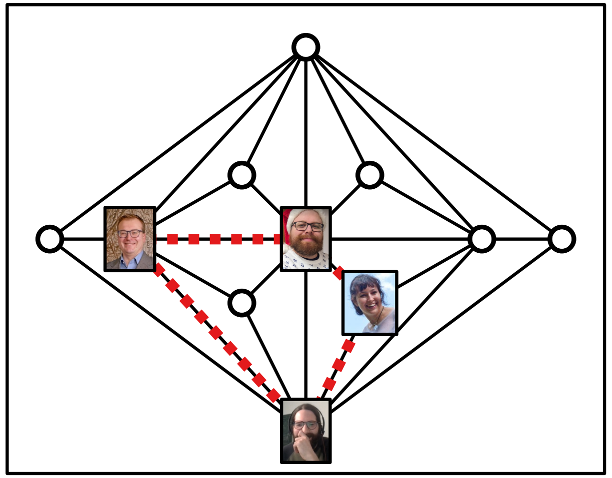
planar graph G

+

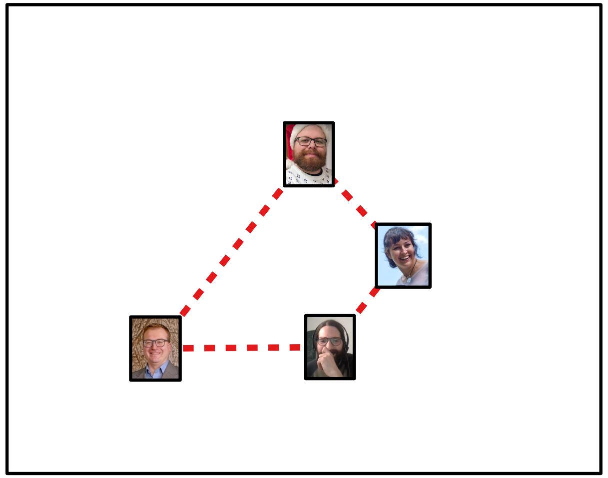


planar drawing of a subgraph $H \subseteq G$

Partial Embedding of **Planar** Graphs



planar graph G



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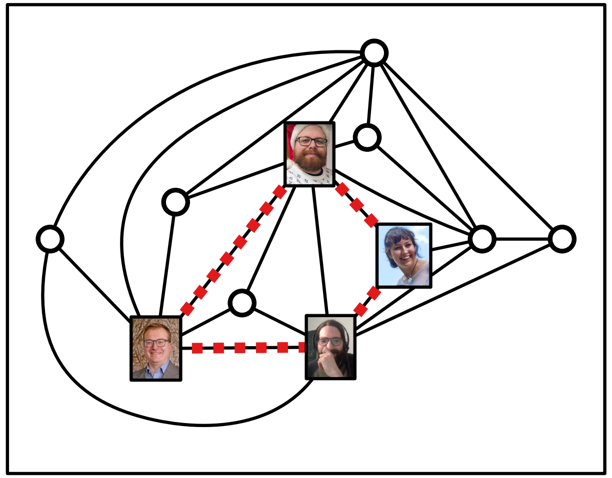
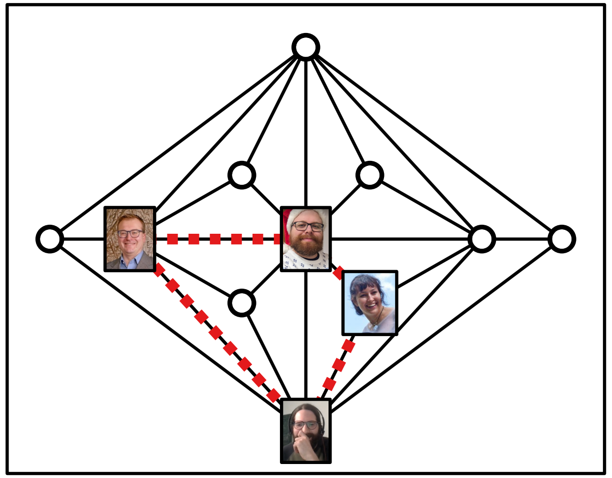
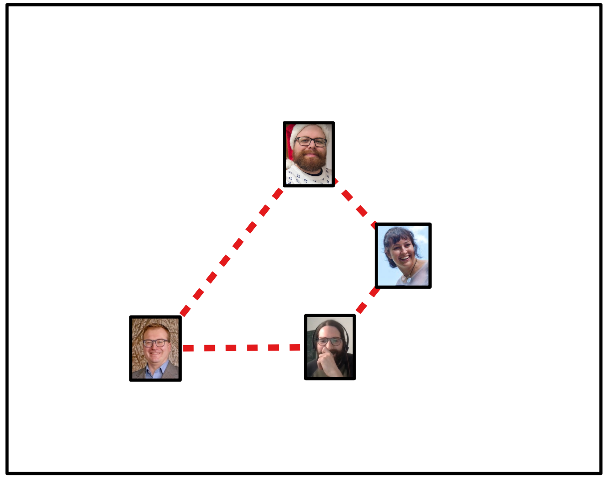


planar drawing of G s.t. H keeps its drawing

Partial Embedding of **Planar** Graphs



planar graph G



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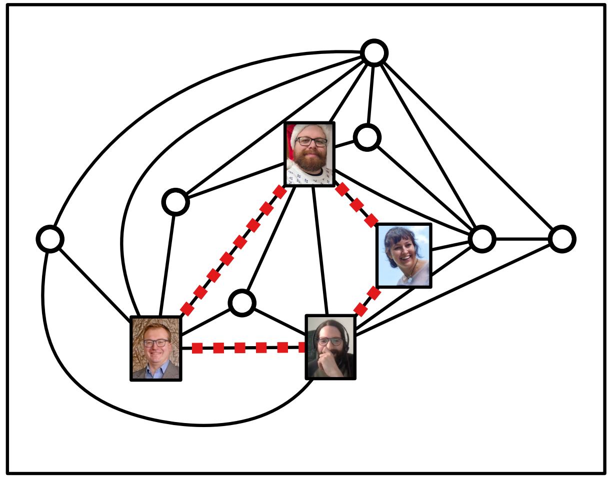
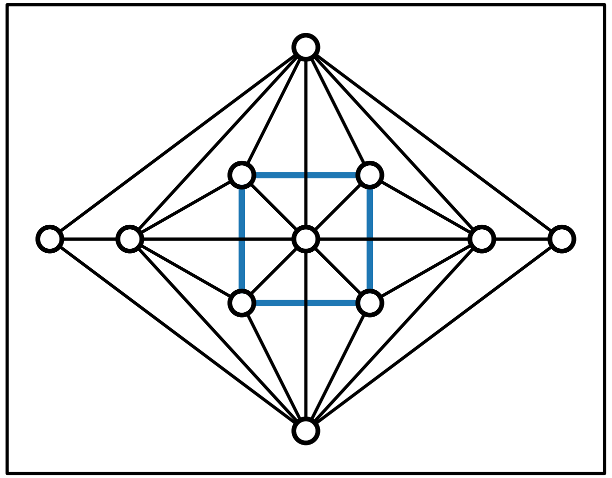


planar drawing of G s.t. H keeps its drawing

This problem can be solved in $\mathcal{O}(n)$ time.

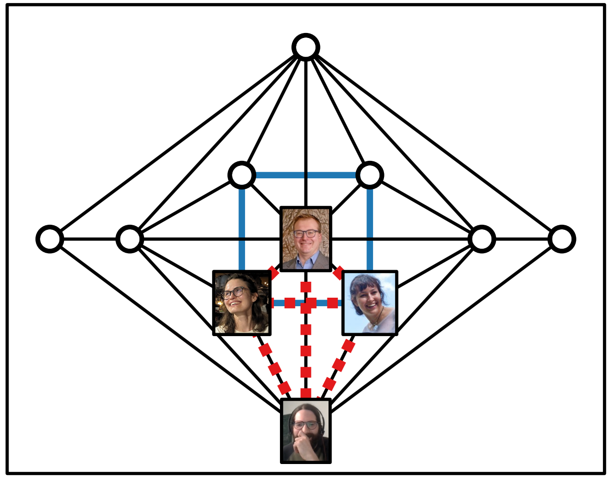
[Angelini, Di Battista, Frati, Jelínek, Kratochvíl, Patrignani, Rutter '10]

Partial Embedding of **1-Planar** Graphs



1-planar graph G
(can be drawn s.t. every edge is crossed at most once)

Partial Embedding of **1-Planar** Graphs



1-planar graph G
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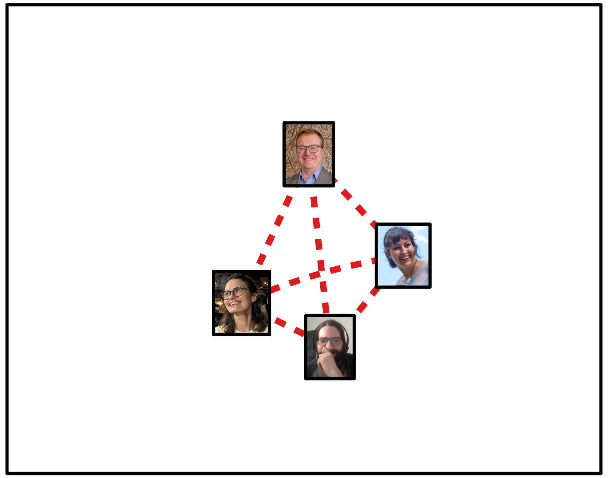
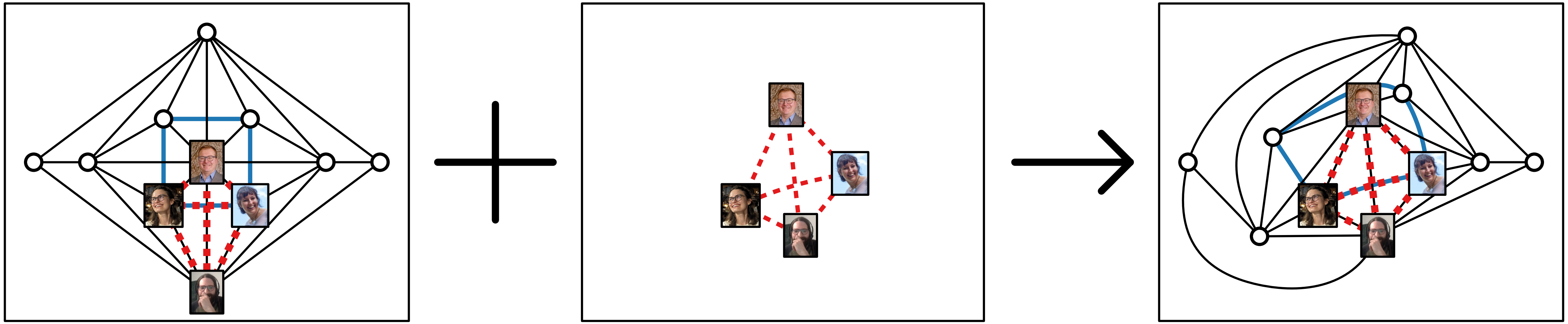
+



1-planar drawing of a subgraph $H \subseteq G$

Partial Embedding of **1-Planar** Graphs

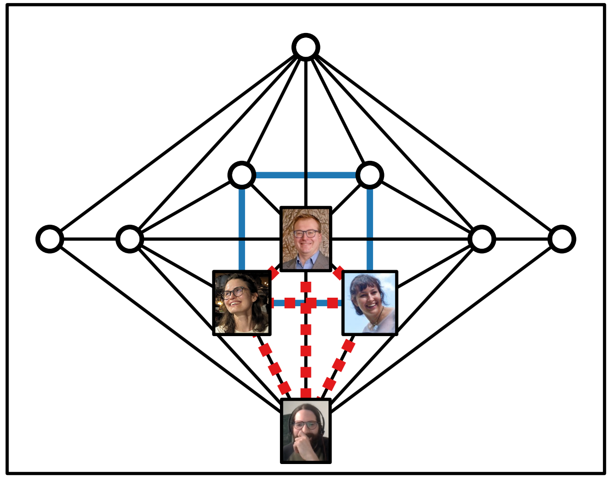


1-planar graph G (can be drawn s.t. every edge is crossed at most once)

1-planar drawing of a subgraph $H \subseteq G$

1-planar drawing of G s.t. H keeps its drawing

Partial Embedding of 1-Planar Graphs



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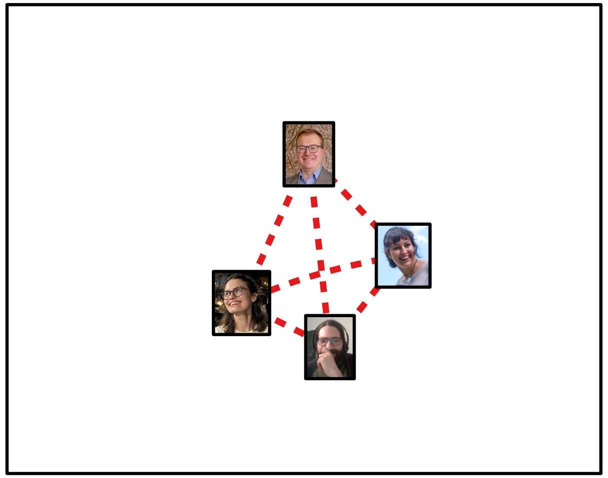
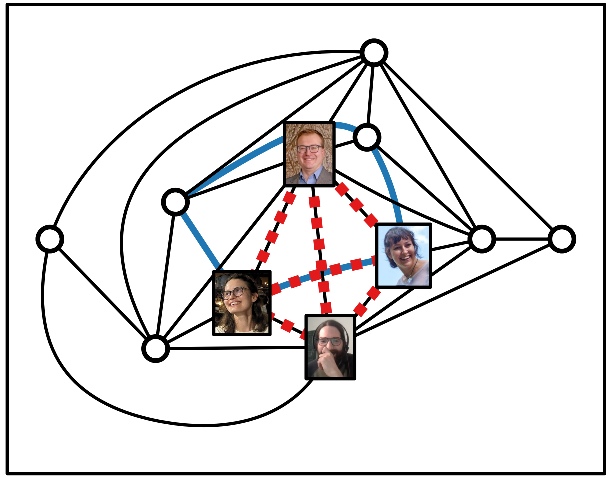


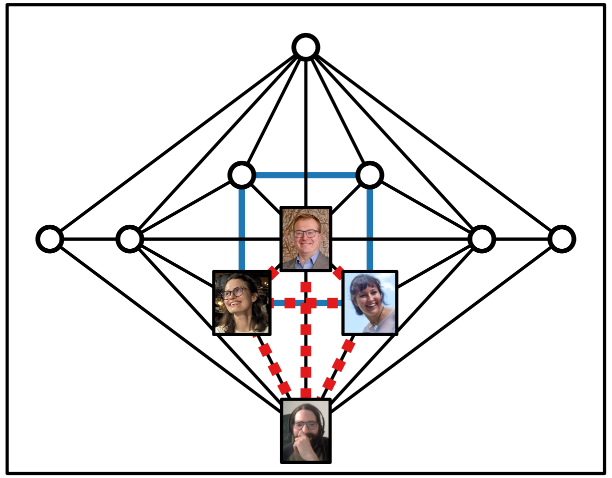
1-planar drawing of a subgraph $H \subseteq G$



1-planar drawing of G s.t. H keeps its drawing

This problem is NP-hard even if $H = \emptyset$. [Grigoriev & Bodlaender '07]

Partial Embedding of 1-Planar Graphs



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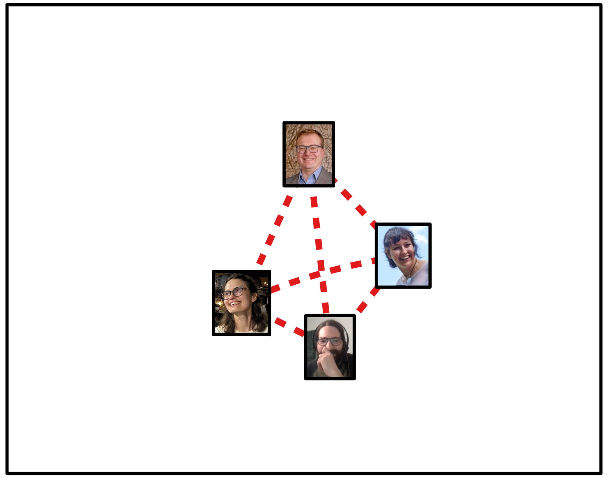

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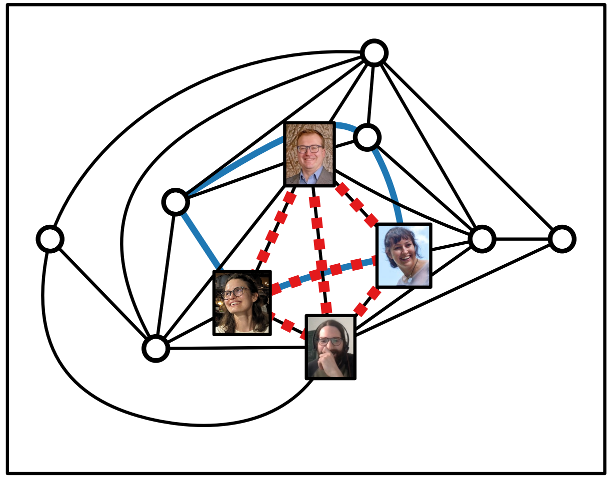
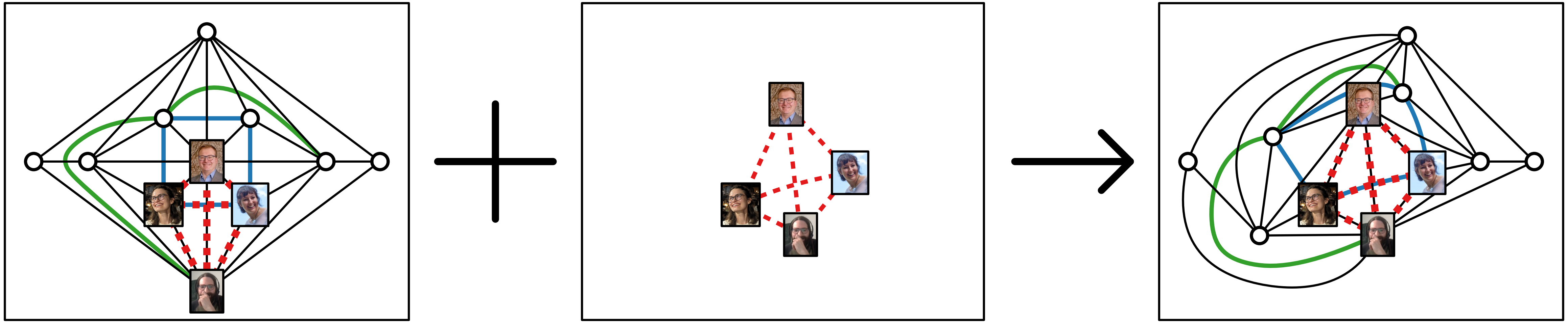


1-planar drawing of G s.t. H keeps its drawing

This problem is NP-hard even if $H = \emptyset$. [Grigoriev & Bodlaender '07]

This problem is in FPT parameterized by the vertex+edge deletion distance between G and H. [Eiben, Ganian, Hamm, Klute & Nöllenburg '20]

Partial Embedding of *k*-**Planar** Graphs

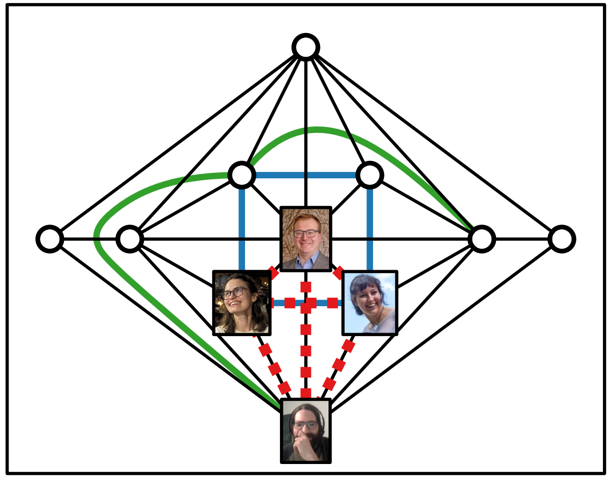


k-**planar** graph G (can be drawn s.t. every edge is crossed at most k times)

k-**planar** drawing of a subgraph $H \subseteq G$

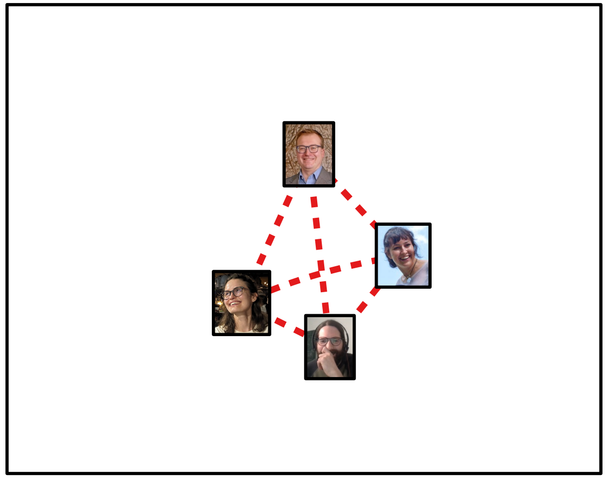
k-**planar** drawing of G s.t. H keeps its drawing

Partial Embedding of *k*-**Planar** Graphs



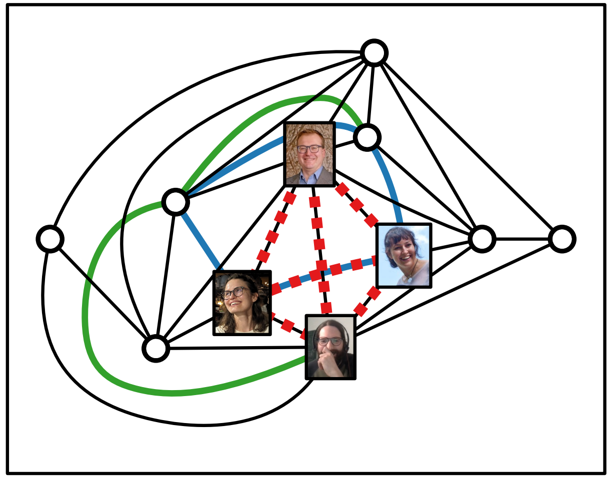
k-**planar** graph G
(can be drawn s.t. every edge is crossed at most k times)

\+



k-**planar** drawing of a subgraph $H \subseteq G$

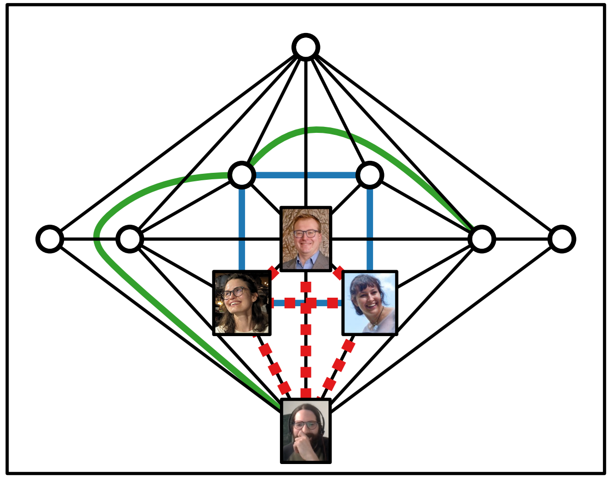
→



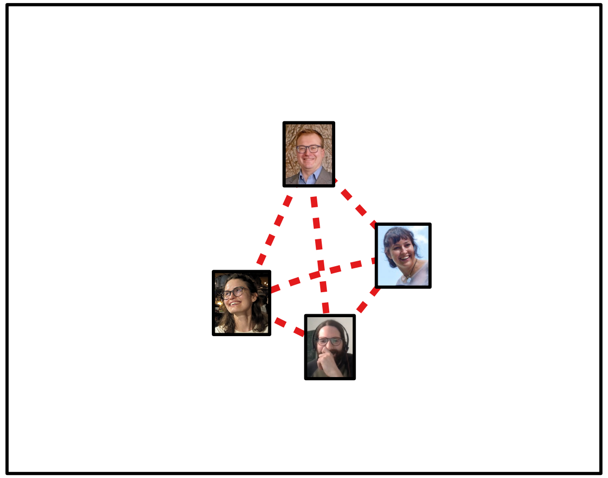
k-**planar** drawing of G s.t. H keeps its drawing

This problem is NP-hard for any constant k even if $H = \emptyset$. [Urschel & Wellens '21]

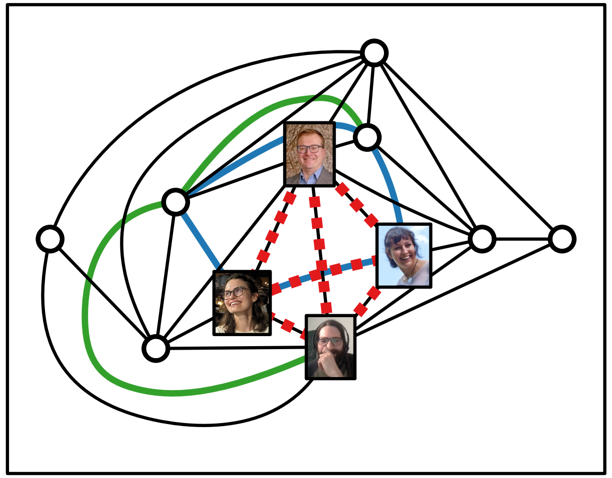
Partial Embedding of *k*-**Planar** Graphs



k-**planar** graph G (can be drawn s.t. every edge is crossed at most k times)



k-**planar** drawing of a subgraph $H \subseteq G$

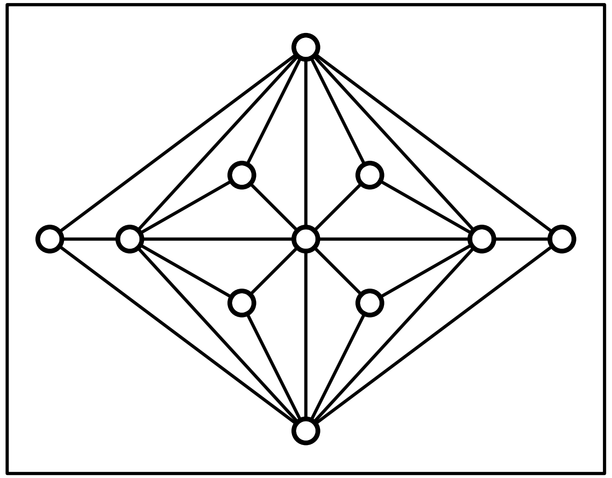


k-**planar** drawing of G s.t. H keeps its drawing

This problem is NP-hard for any constant k even if $H = \emptyset$. [Urschel & Wellens '21]

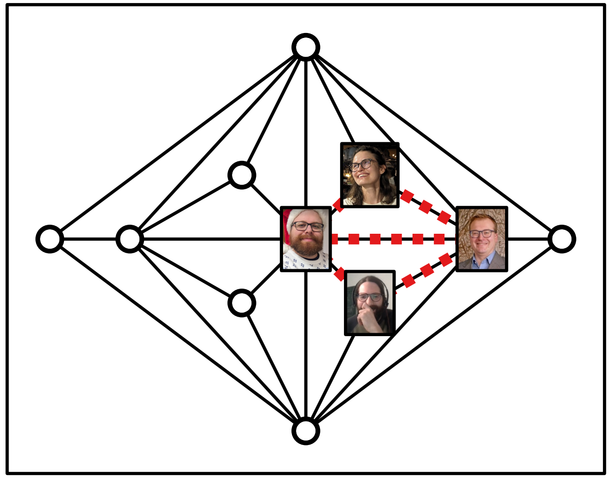
This problem is in FPT parameterized by k + #edges in $G - H$. [Ganian, Hamm, Klute, Parada & Vogtenhuber '21]

Generalization of Partial Embedding



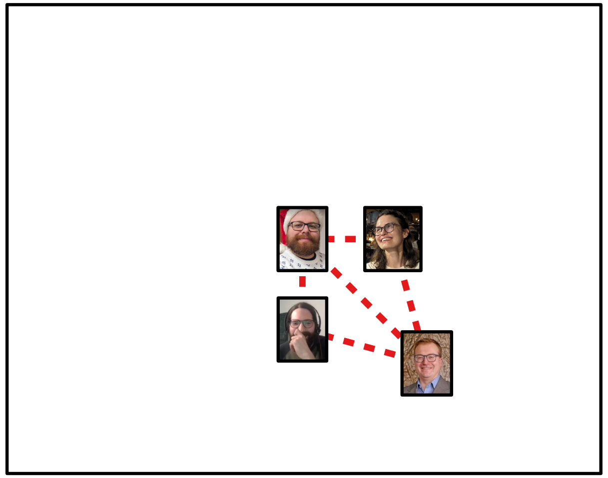
Graph G +
drawing style Φ
(e.g., planar)

Generalization of Partial Embedding



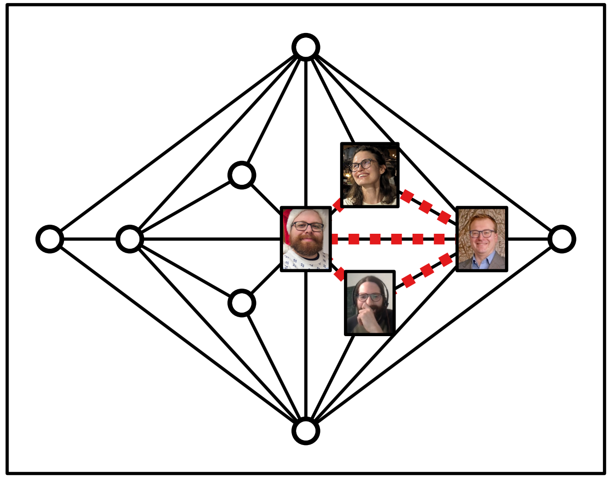
Graph G +
drawing style Φ
(e.g., planar)

+

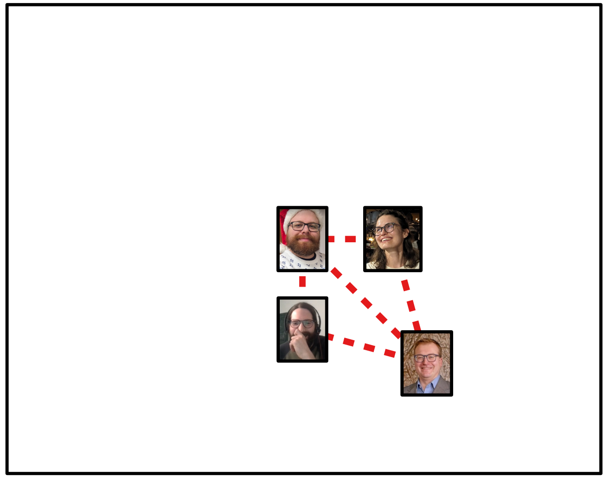


Drawing with style
$\Phi' \subseteq \Phi$
(e.g., straight-line planar)
of a subgraph $H \subseteq G$

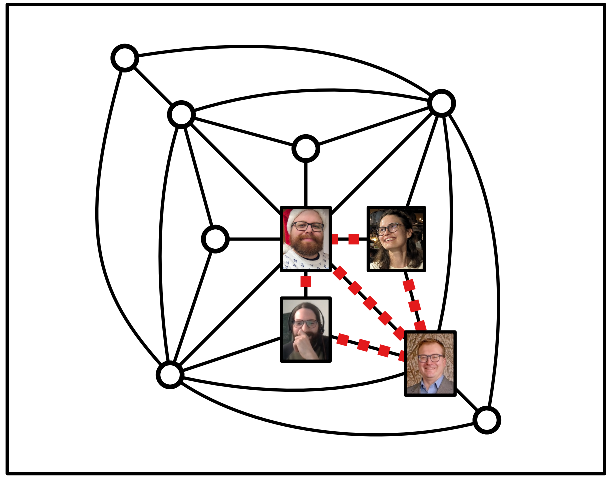
Generalization of Partial Embedding



Graph G +
drawing style Φ
(e.g., planar)

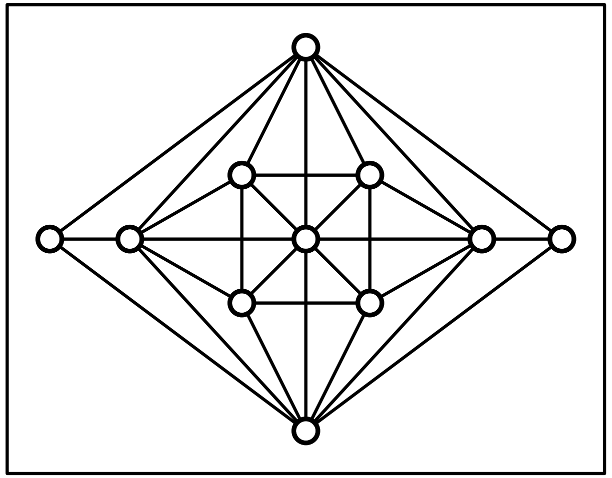


Drawing with style
$\Phi' \subseteq \Phi$
(e.g., straight-line planar)
of a subgraph $H \subseteq G$



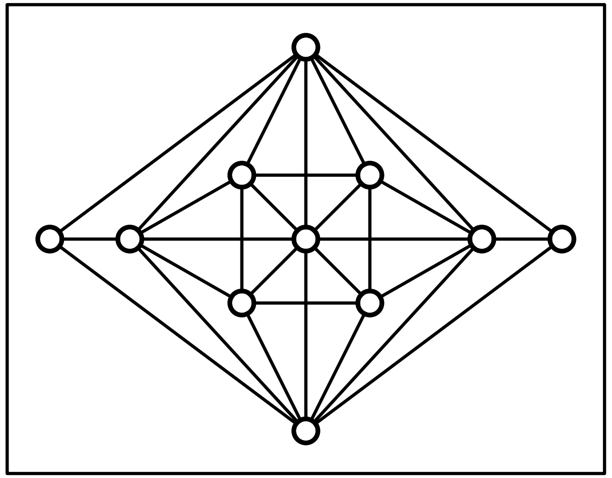
Drawing with style Φ of G
s.t. H **keeps its drawing**

1-Plane Insertion Into a Plane Graph



1-planar graph G

1-Plane Insertion Into a Plane Graph



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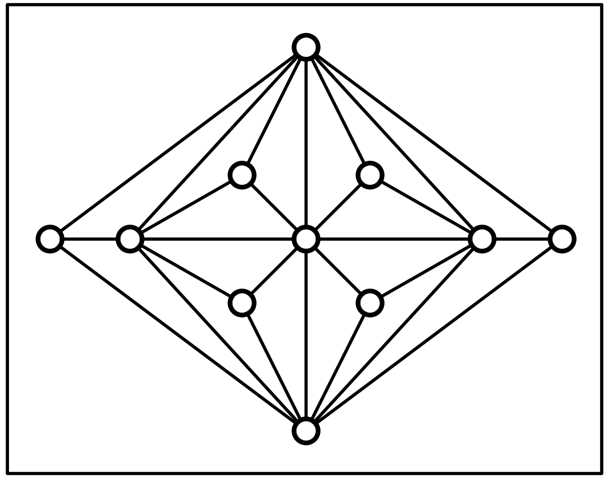
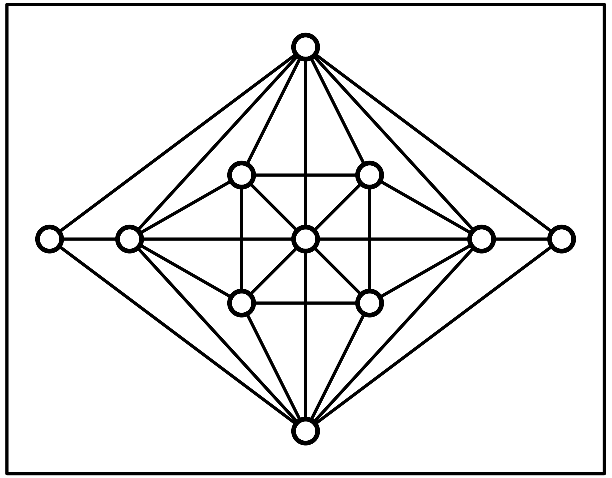


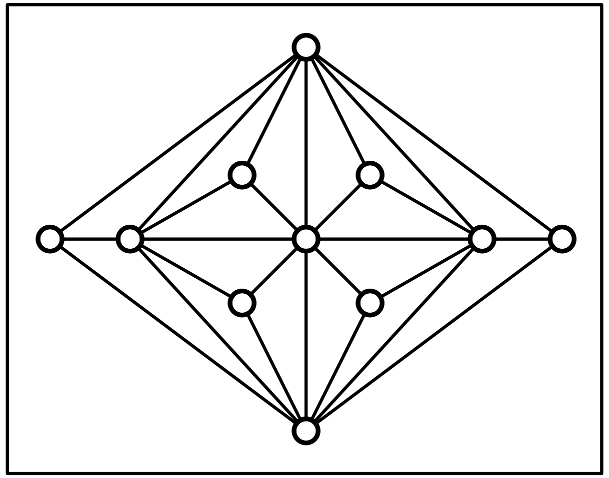
planar drawing of a **spanning** subgraph $H \subseteq G$

1-Plane Insertion Into a Plane Graph



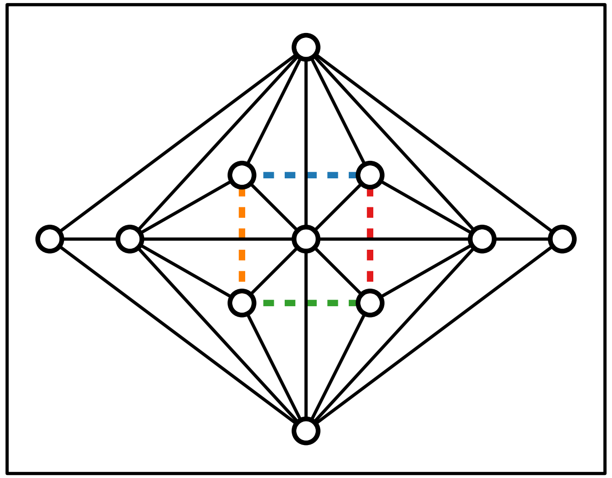
1-planar graph G

+



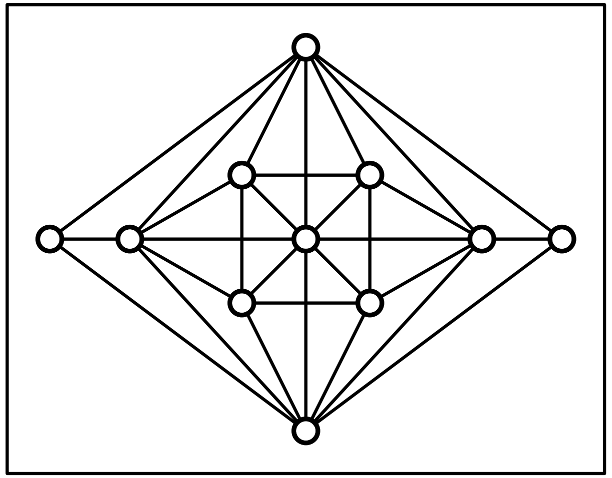
planar drawing of a **spanning** subgraph $H \subseteq G$

→

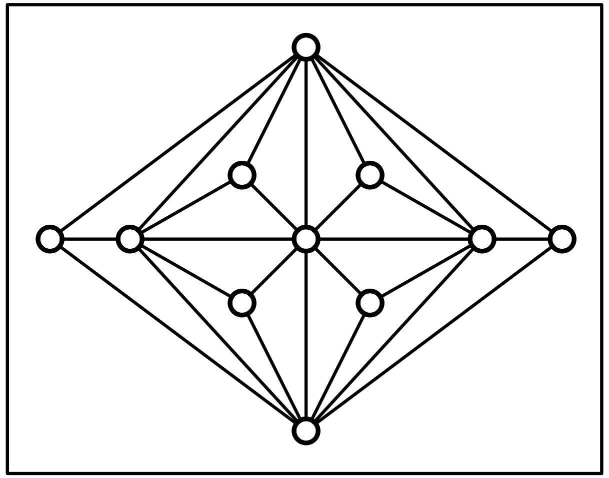


1-planar drawing of G that keeps the drawing of H

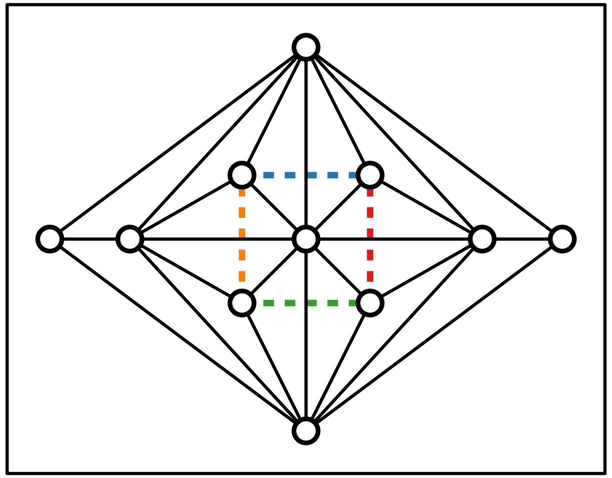
1-Plane Insertion Into a Plane Graph



1-planar graph G



planar drawing of a spanning subgraph $H \subseteq G$



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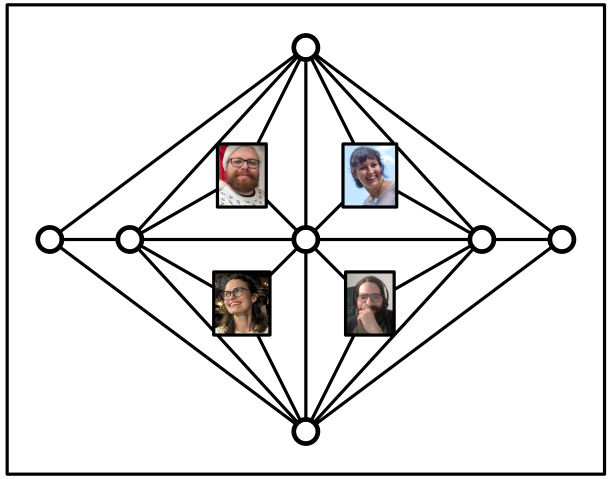
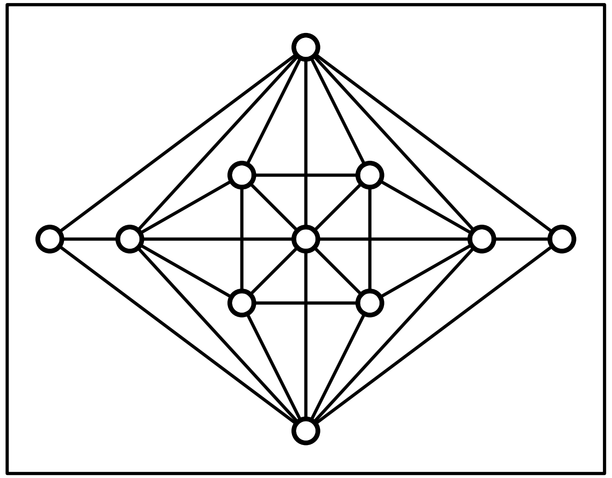
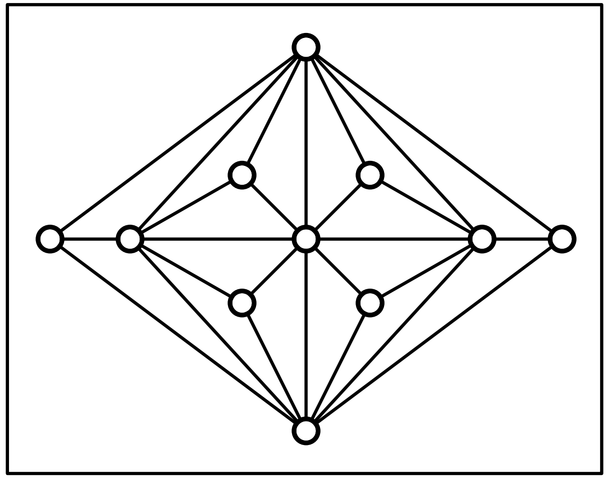


plane graph G

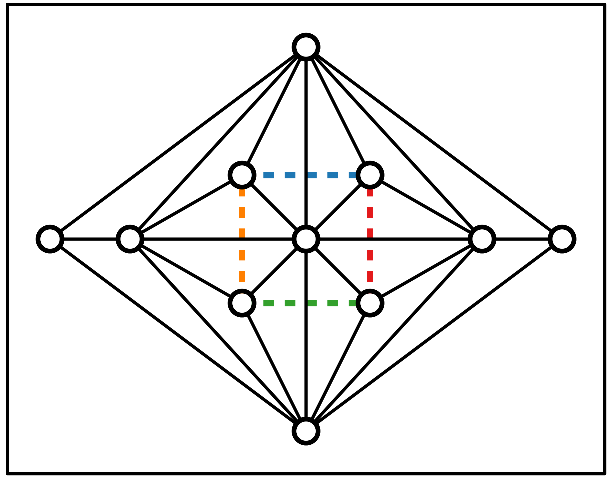
1-Plane Insertion Into a Plane Graph



1-planar graph G



planar drawing of a **spanning** subgraph $H \subseteq G$



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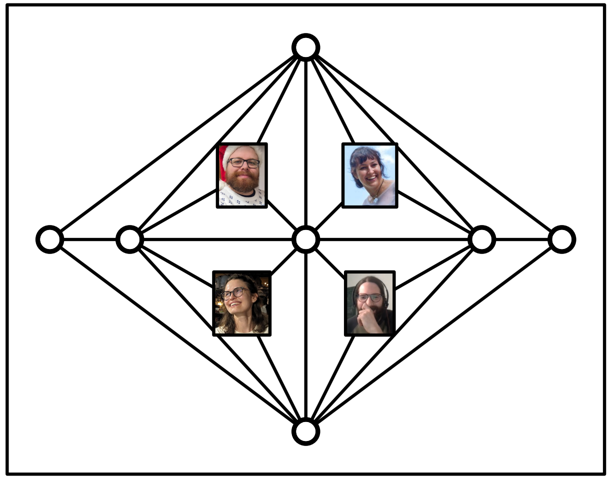
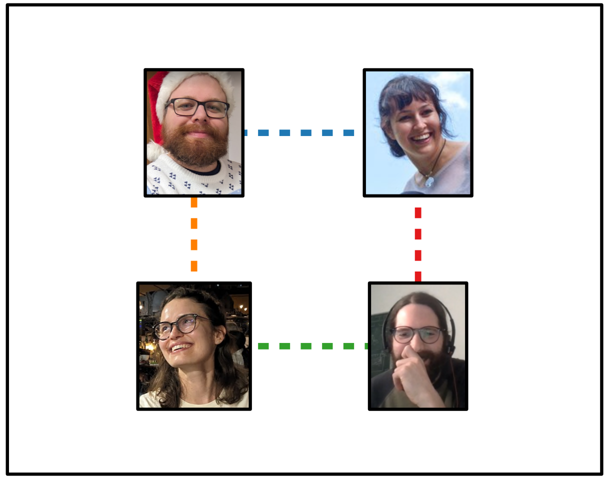
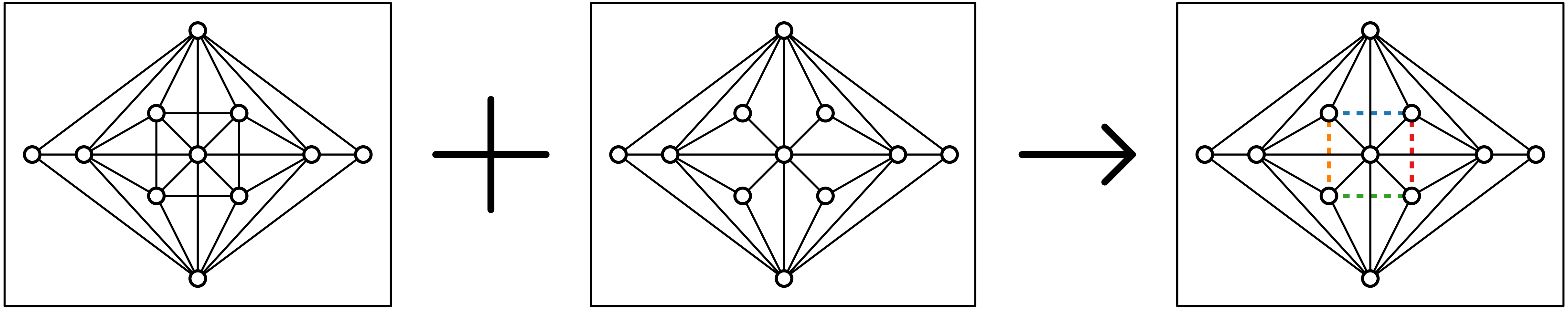


plane graph G



edges E' btw. vtcs in G

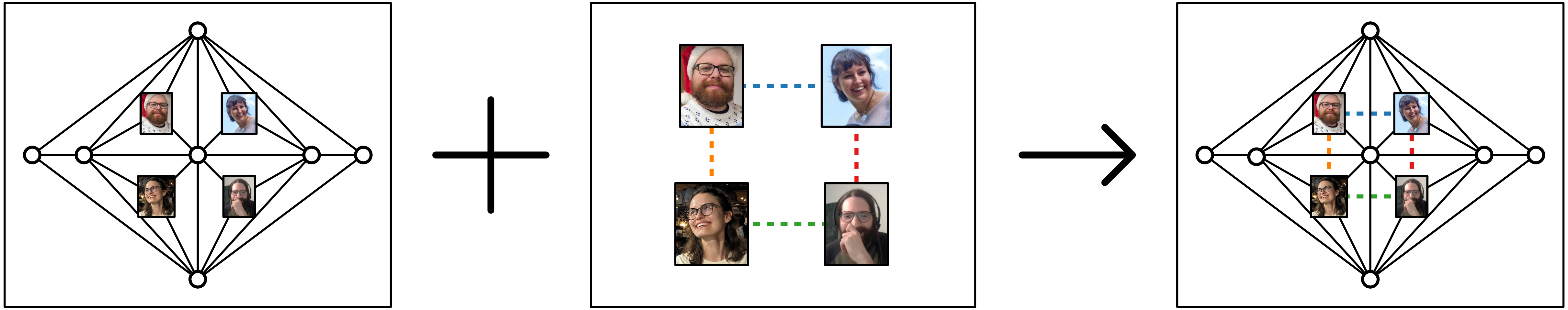
1-Plane Insertion Into a Plane Graph



1-planar graph G

planar drawing of a **spanning** subgraph $H \subseteq G$

1-planar drawing of G that keeps the drawing of H

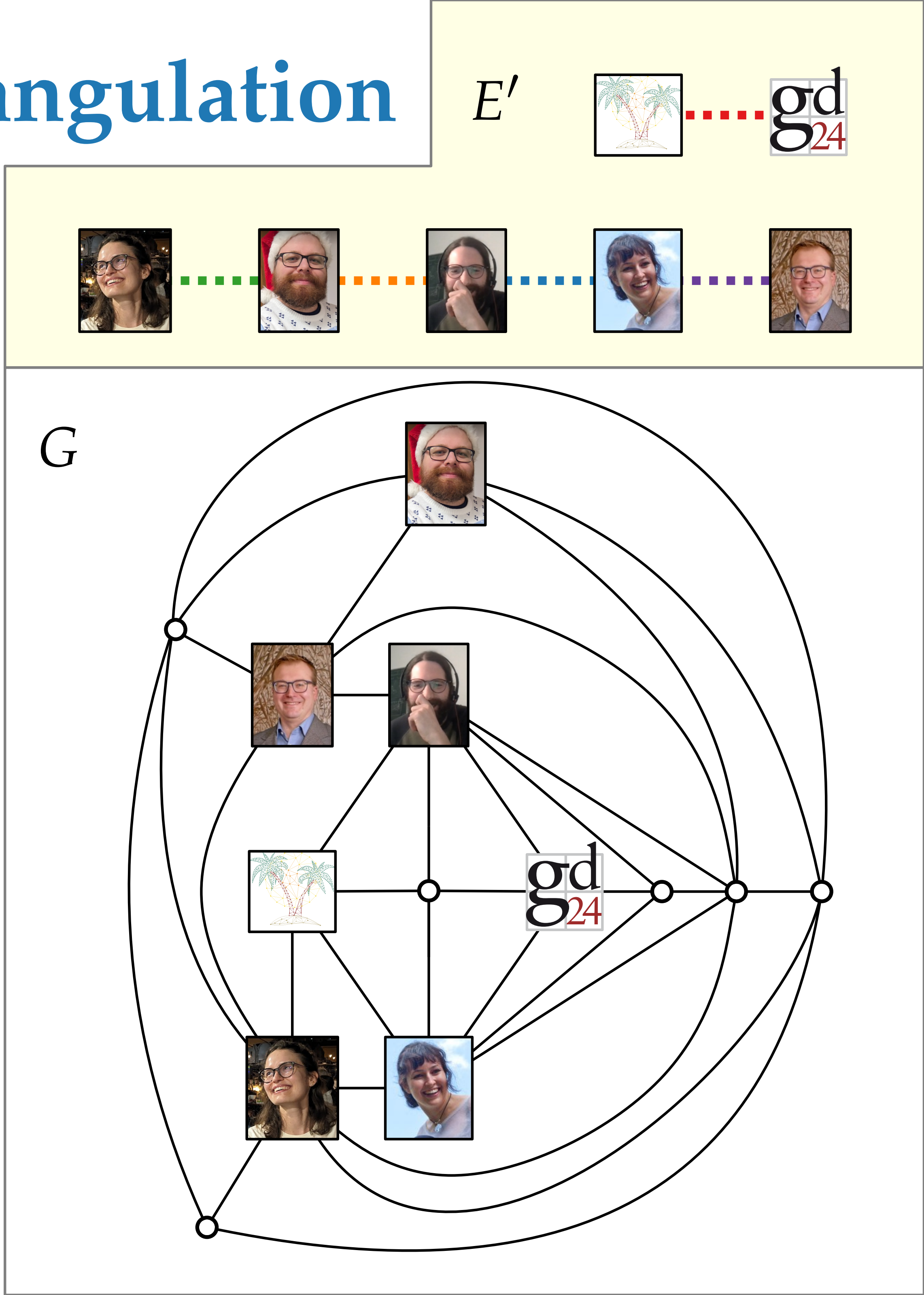


plane graph G

edges E' btw. vtcs in G

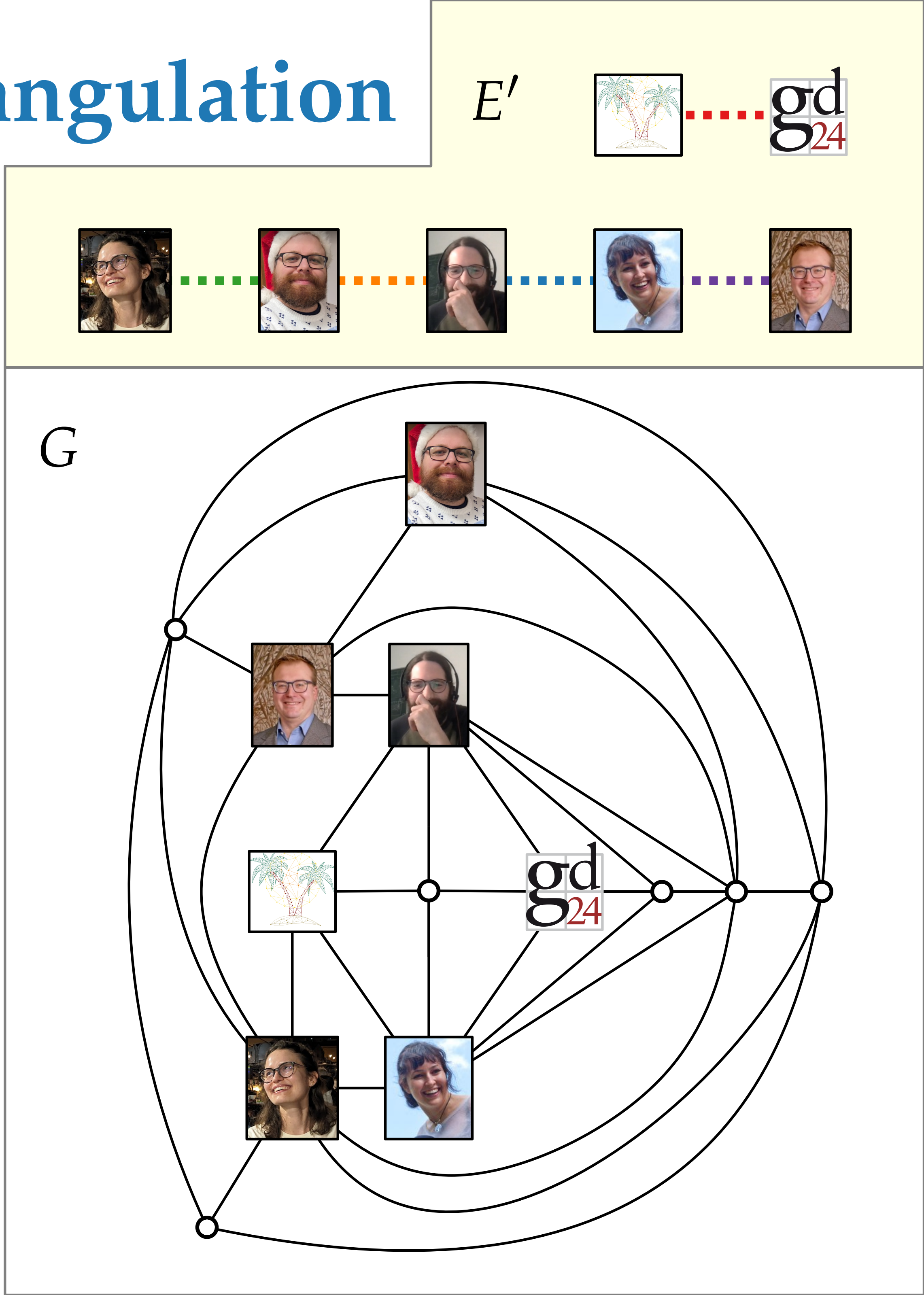
1-planar drawing of $G + E'$ that keeps the embedding of G

1-Plane Insertion Into a Plane **Triangulation**



1-Plane Insertion Into a Plane **Triangulation**

1. For each edge, find all possibilities to route it



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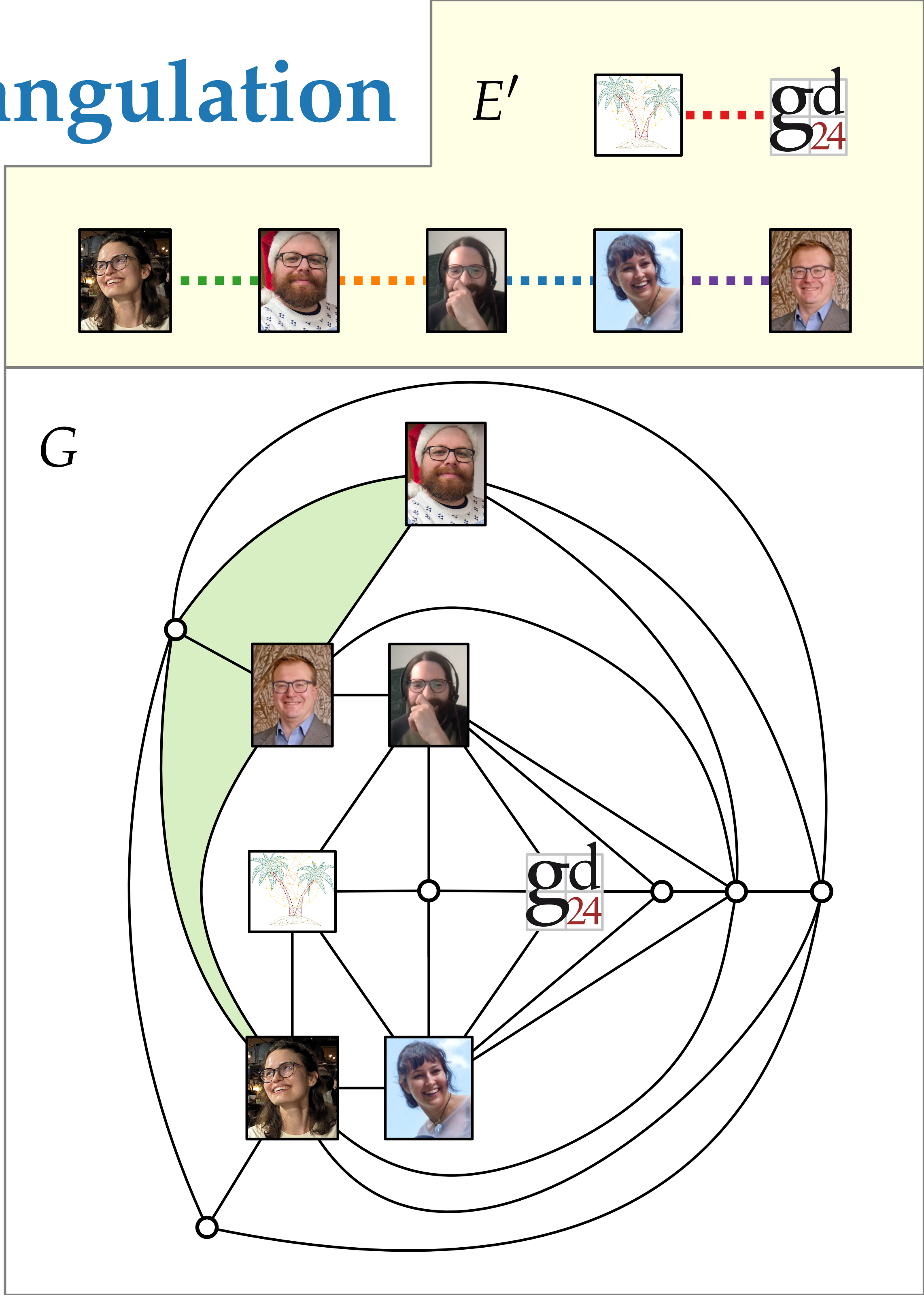

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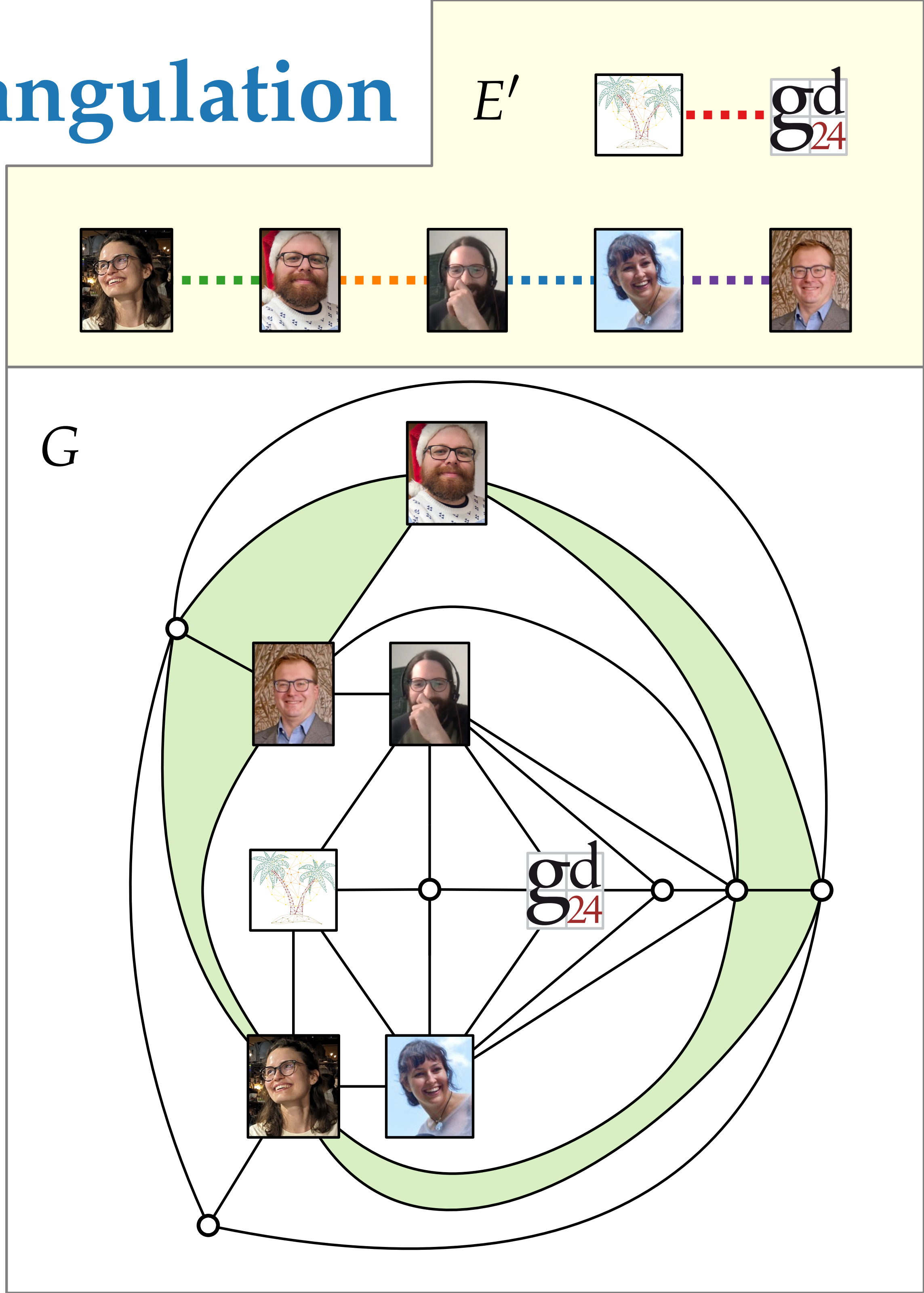

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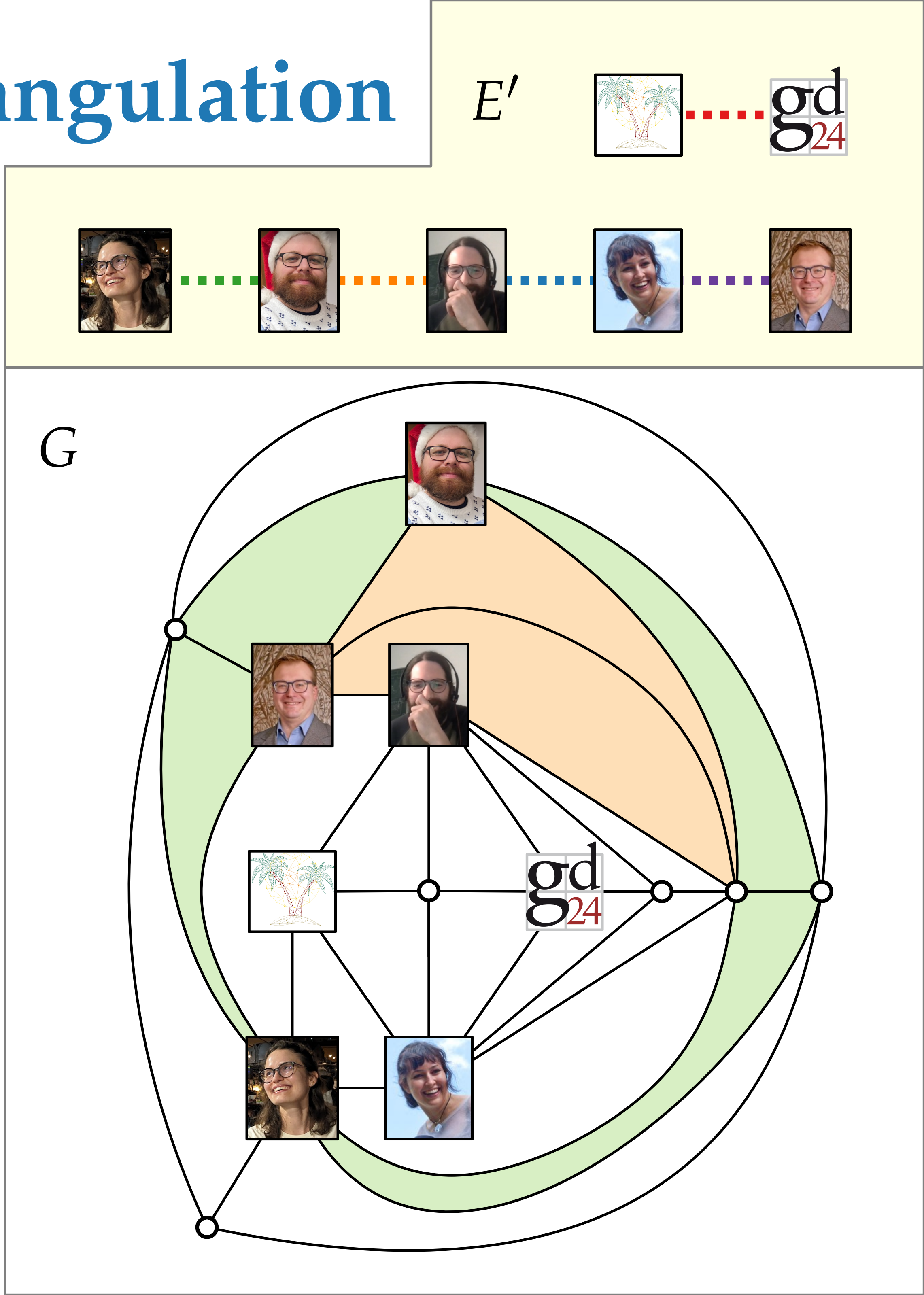

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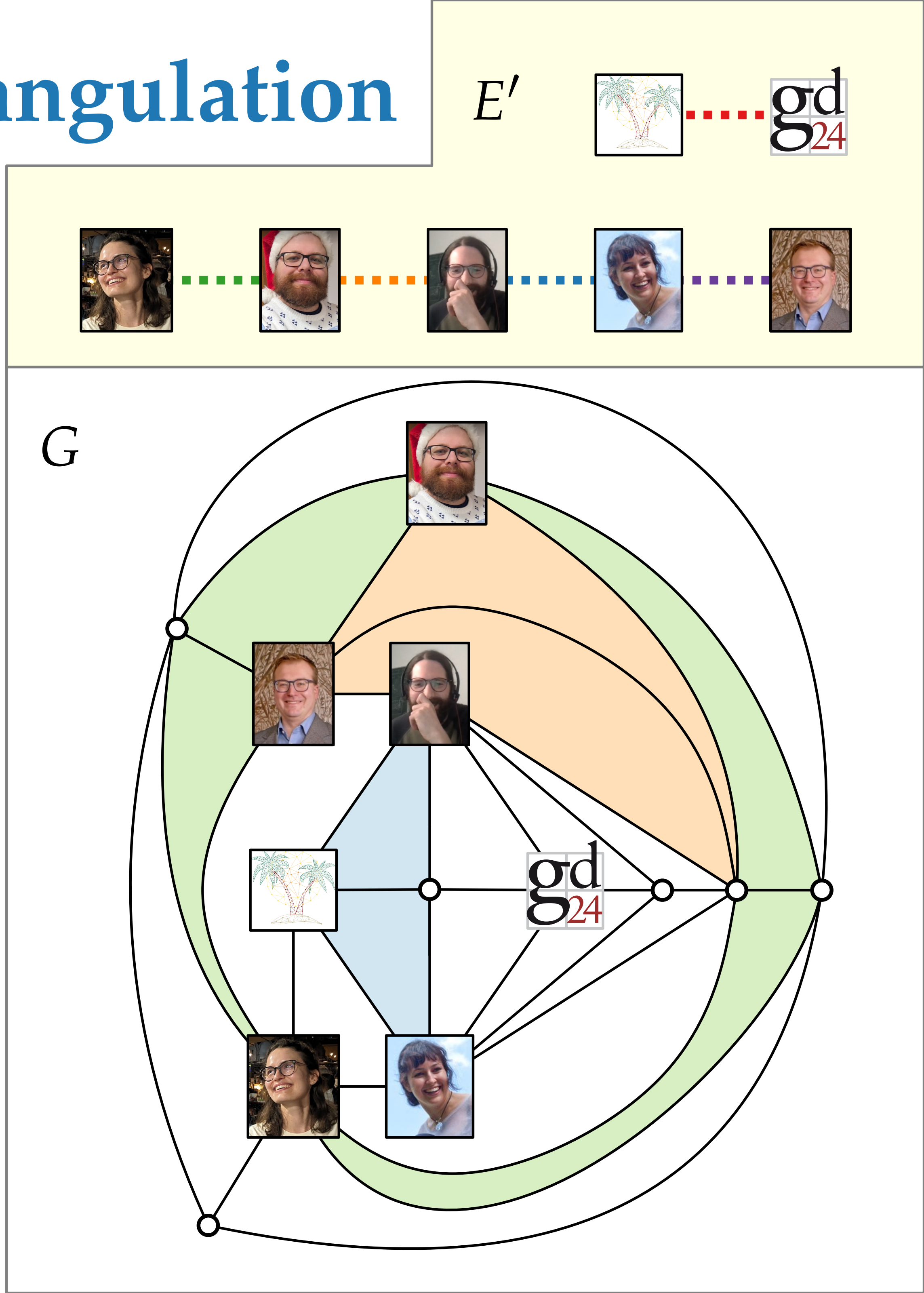

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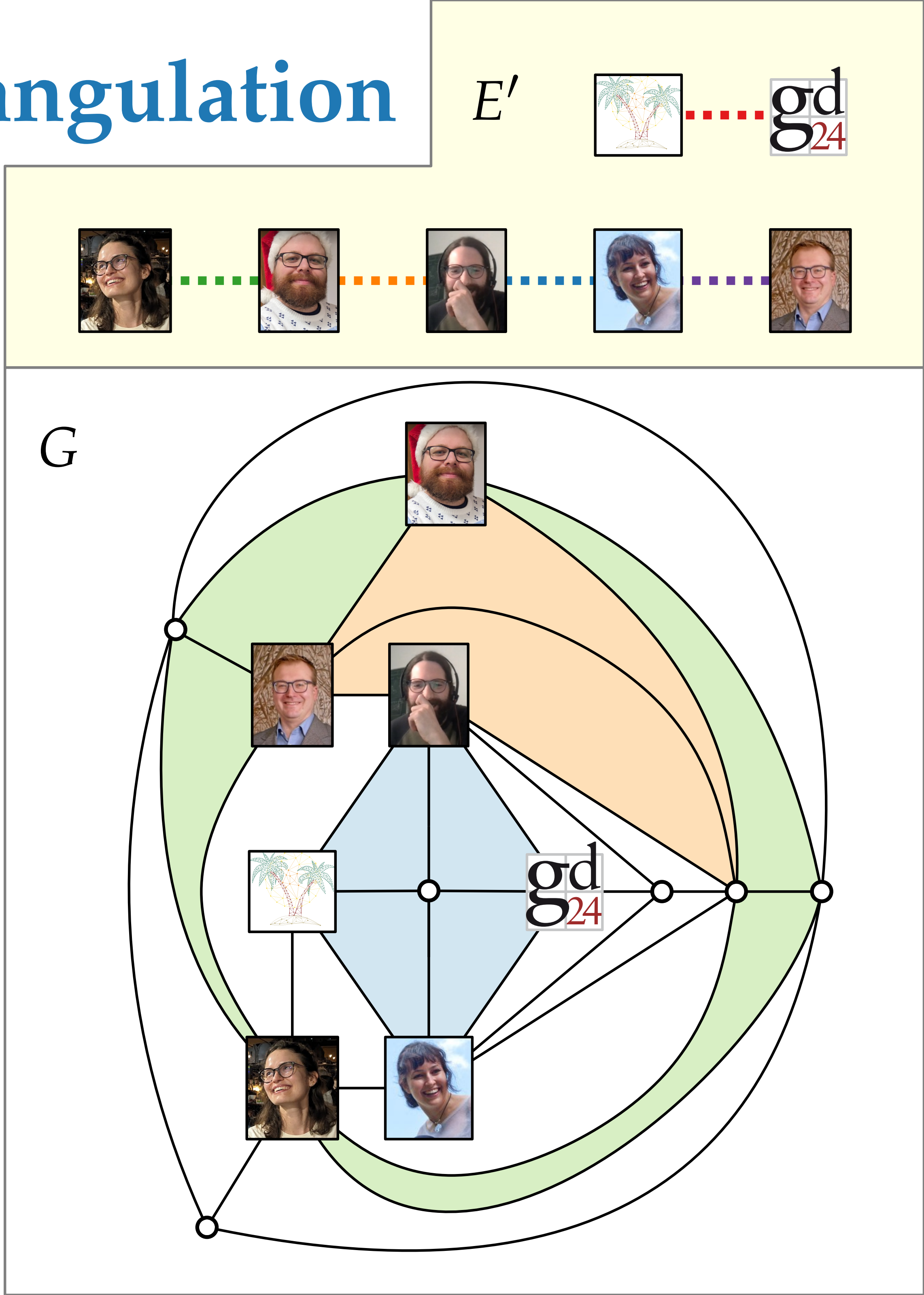

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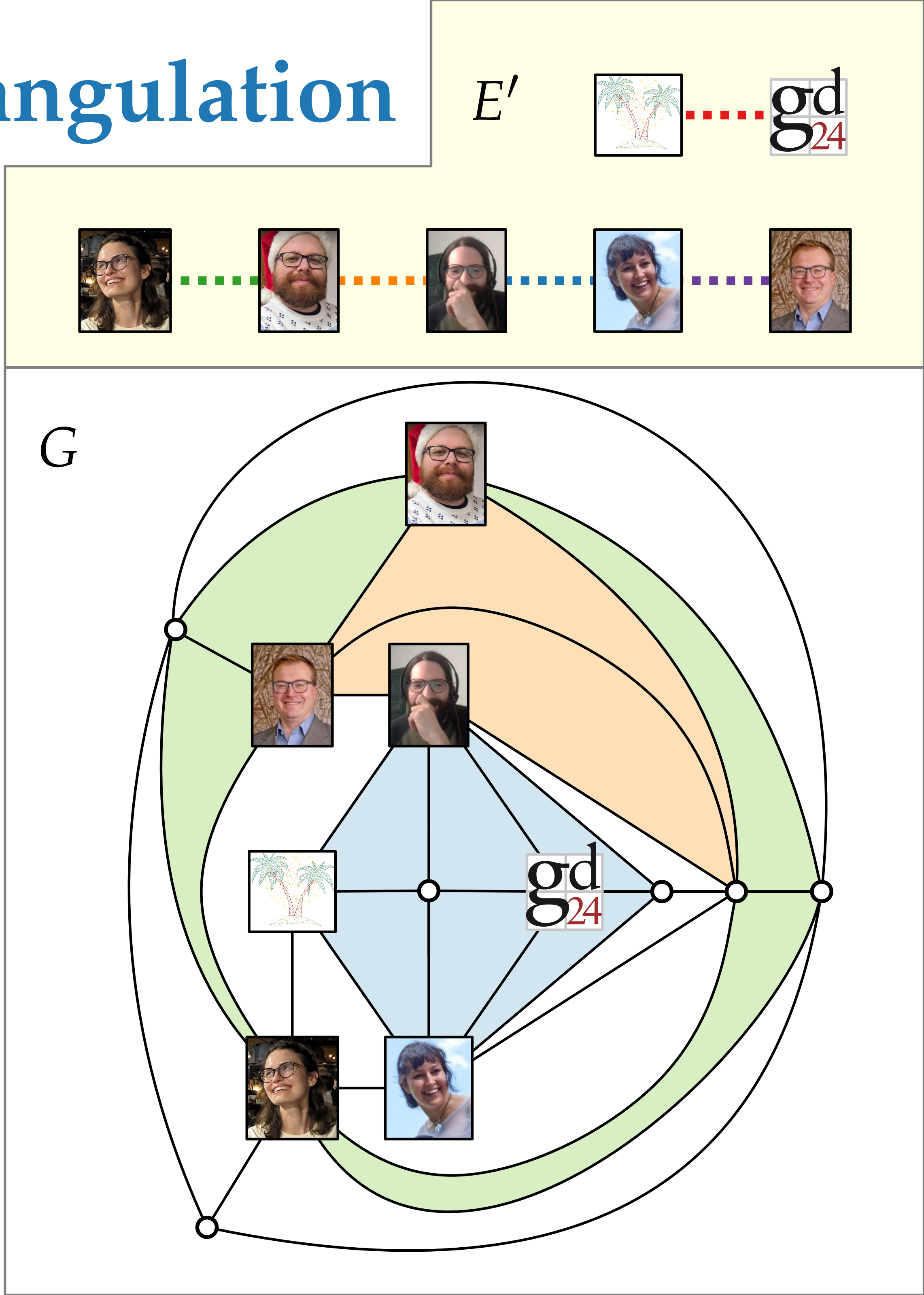

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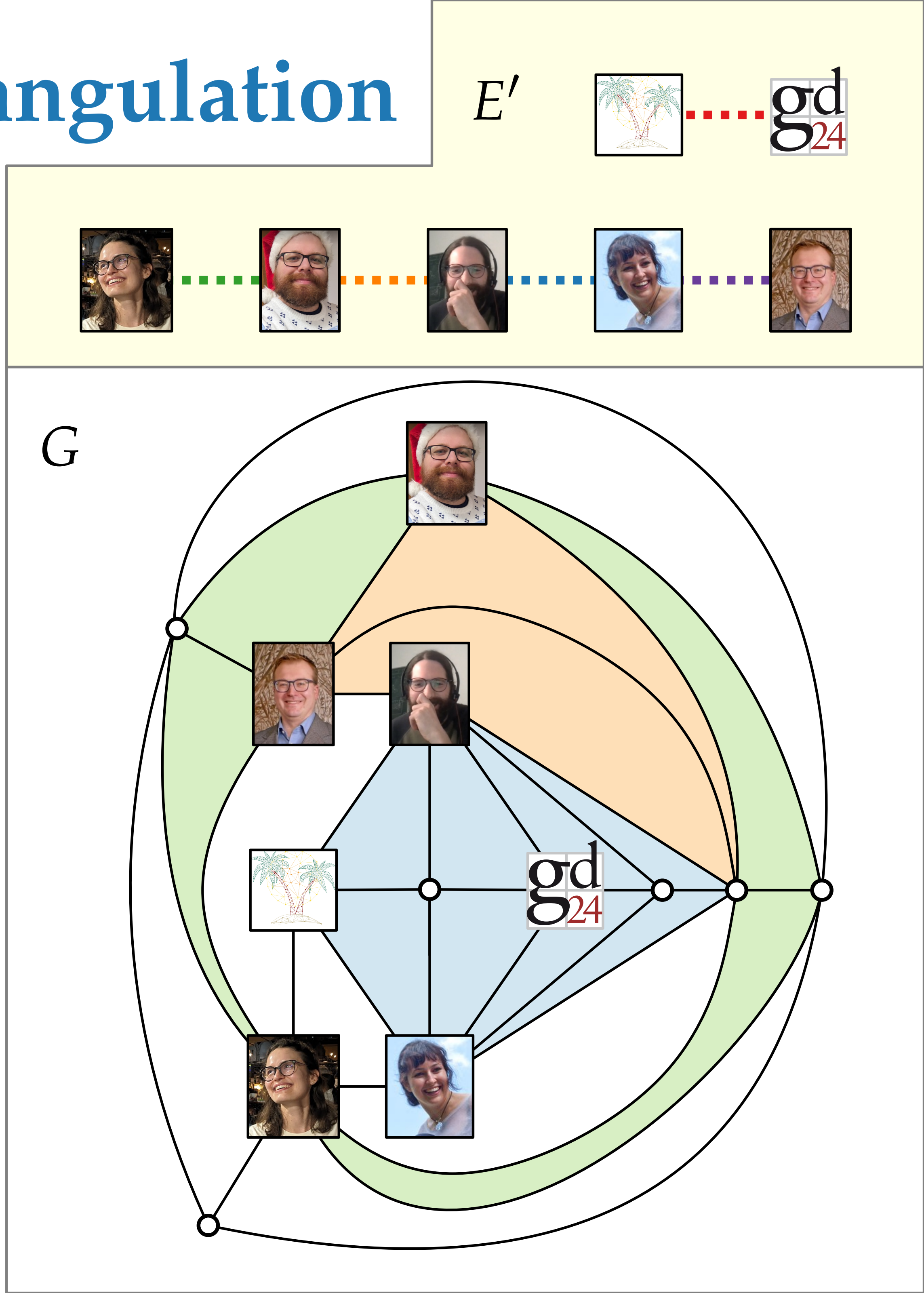

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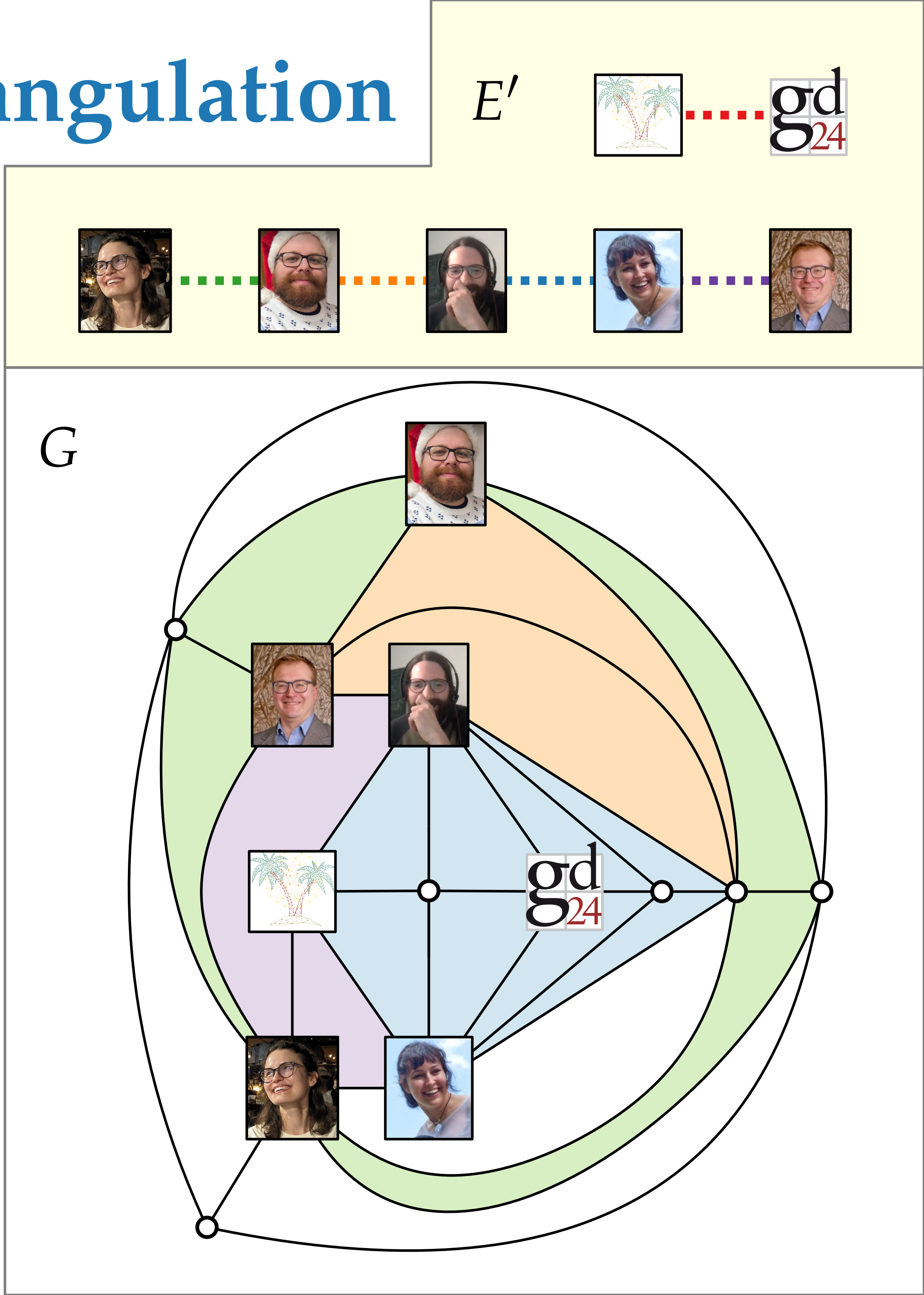

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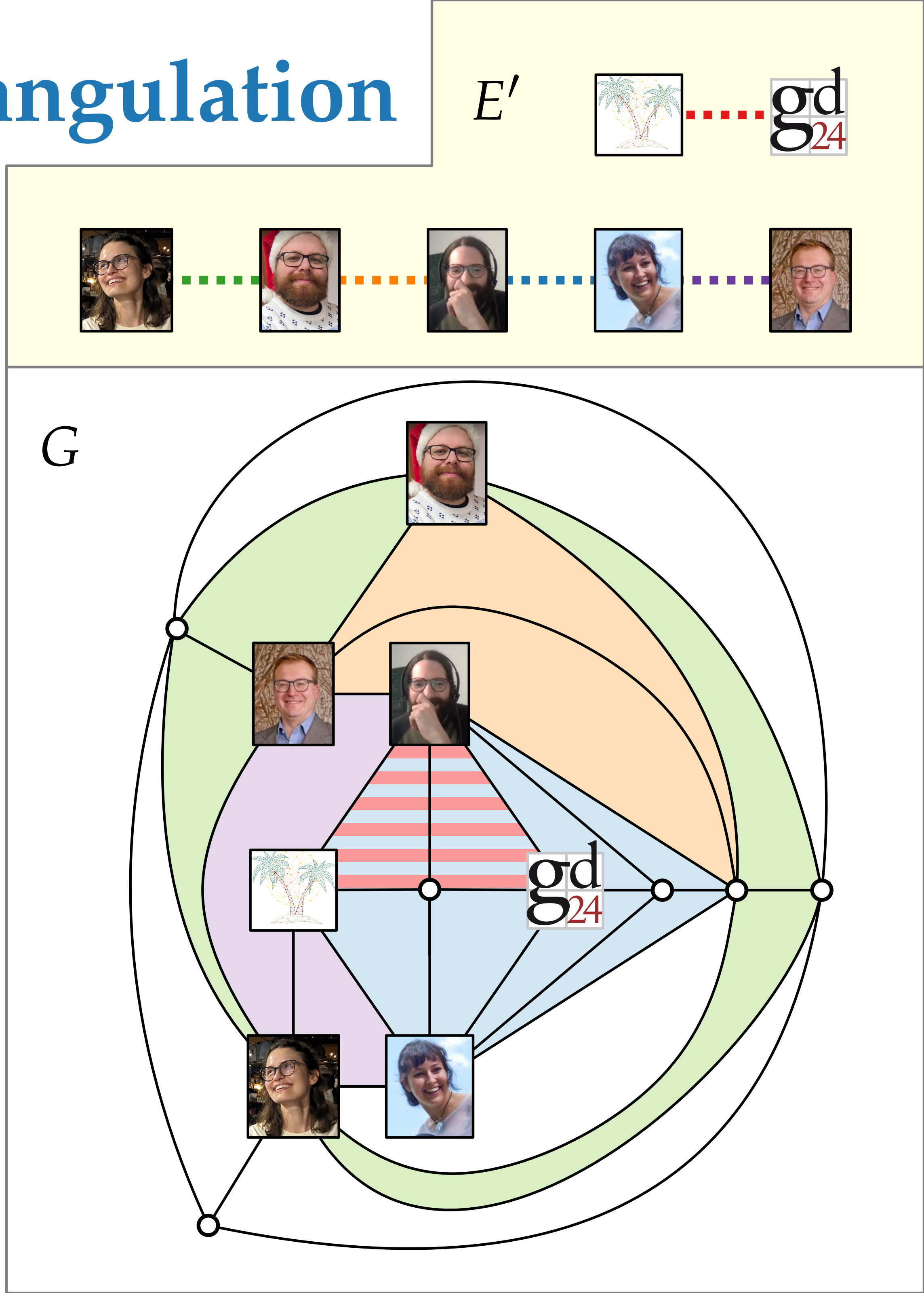

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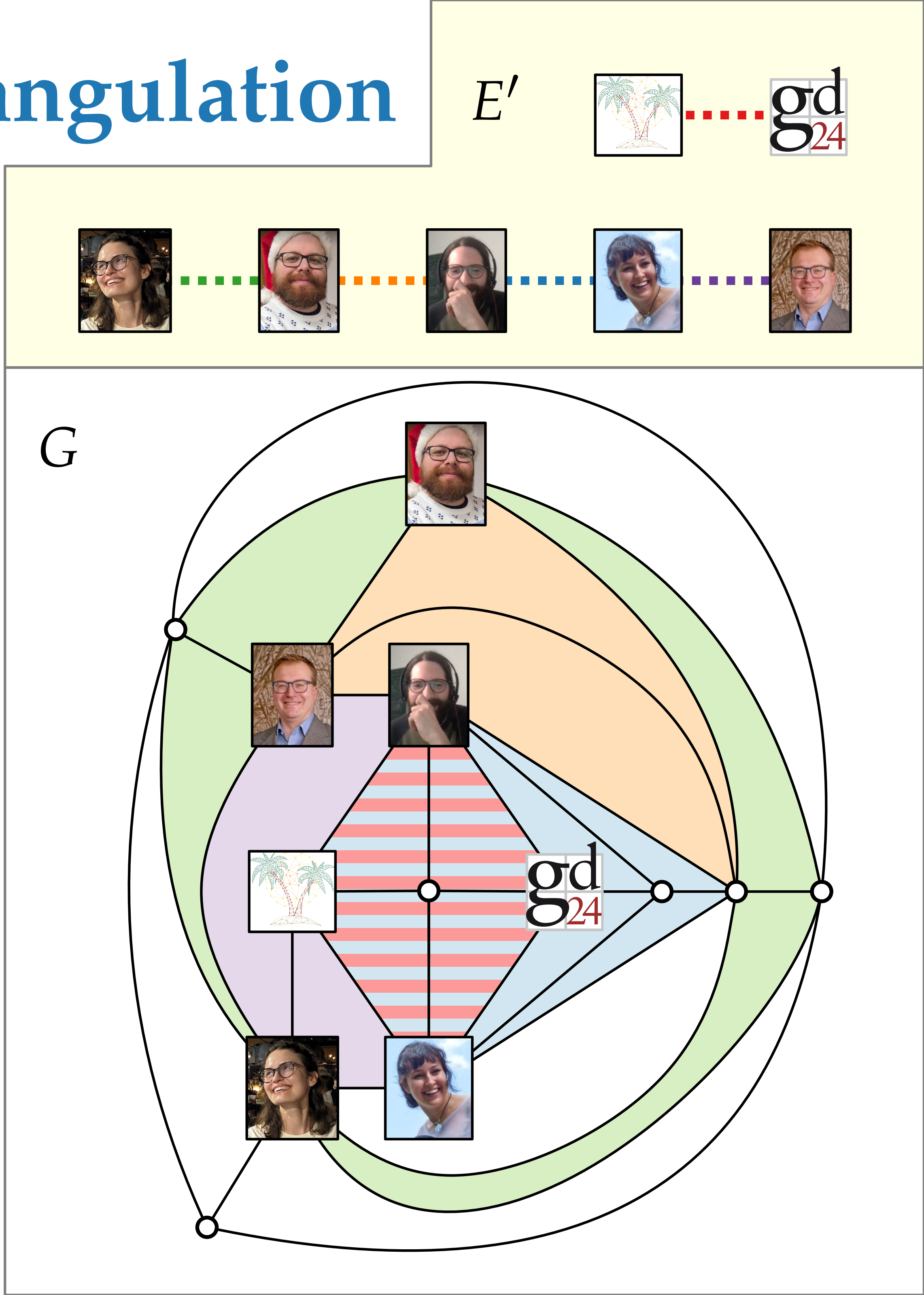

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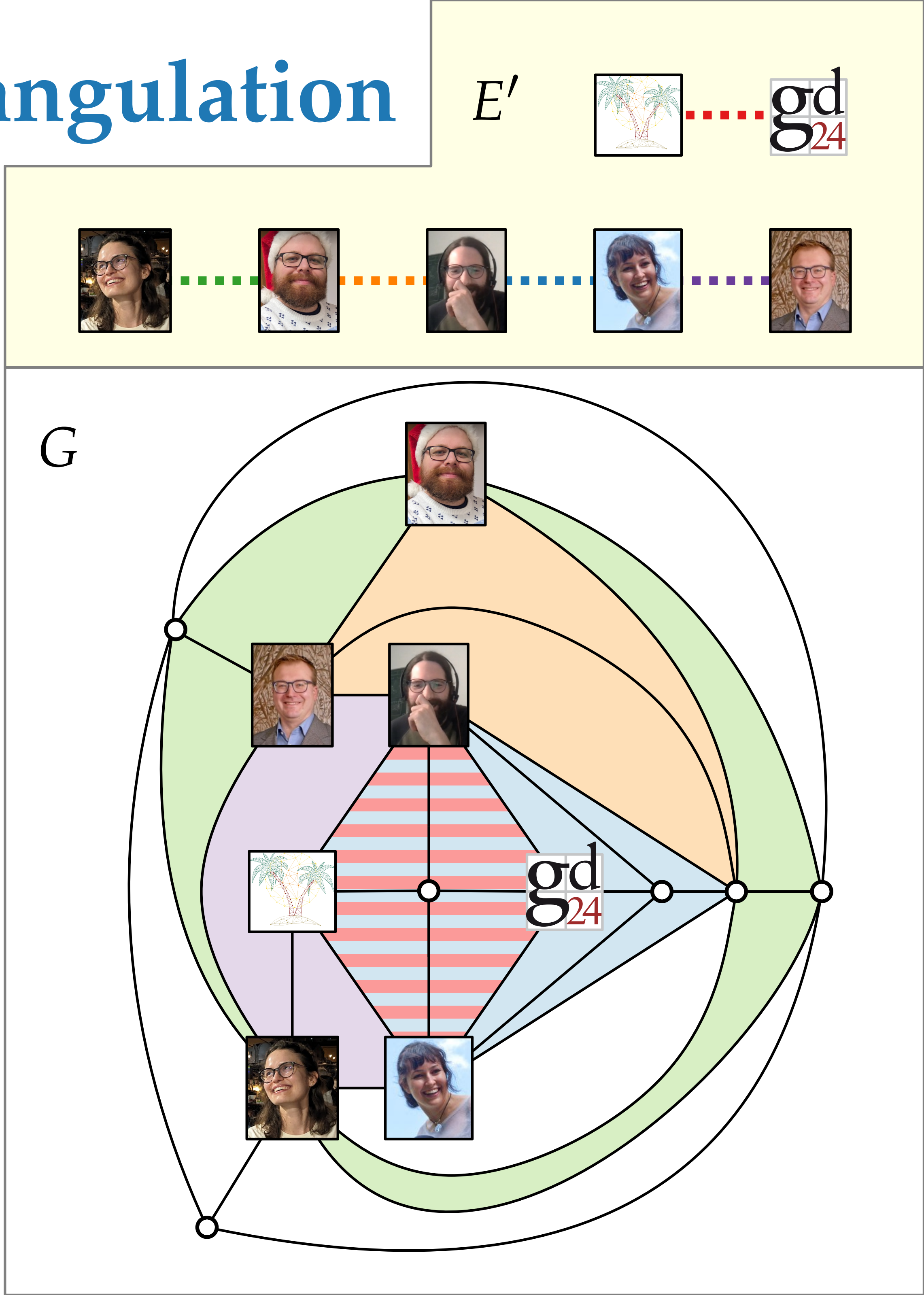
1. For each edge, find all possibilities to route it
2. Edge with 0 options ➡ no-instance



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1. For each edge, find all possibilities to route it
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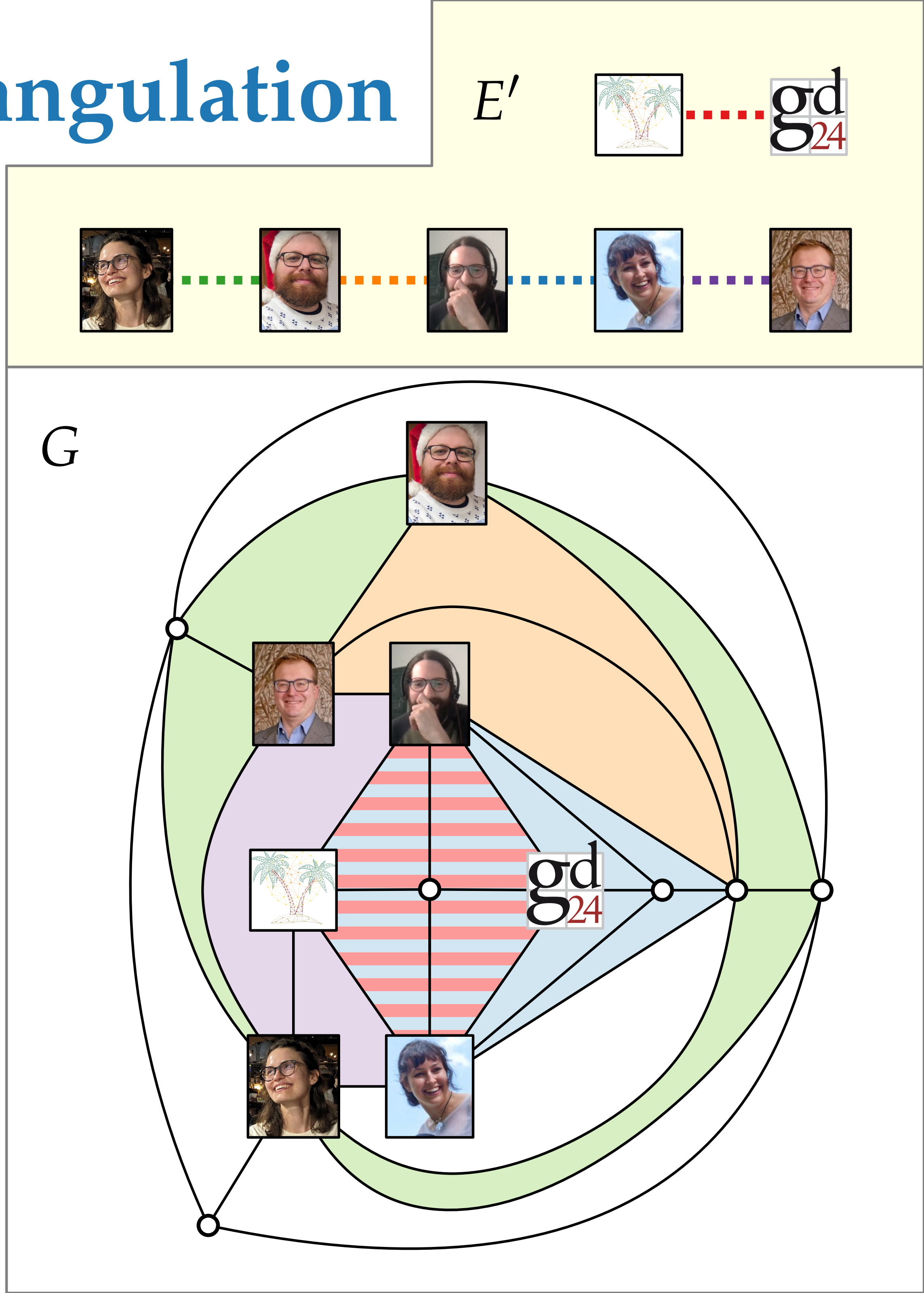

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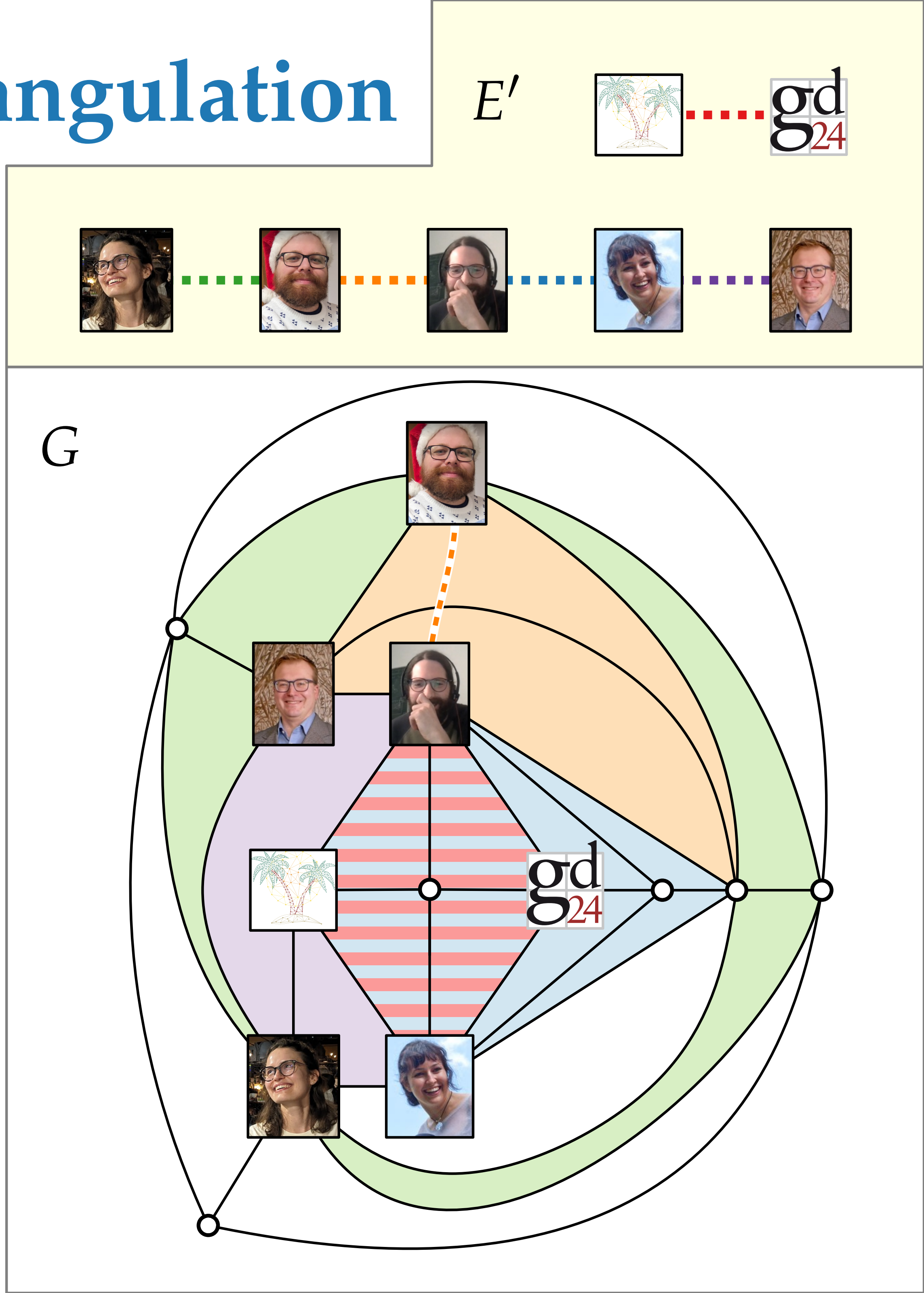

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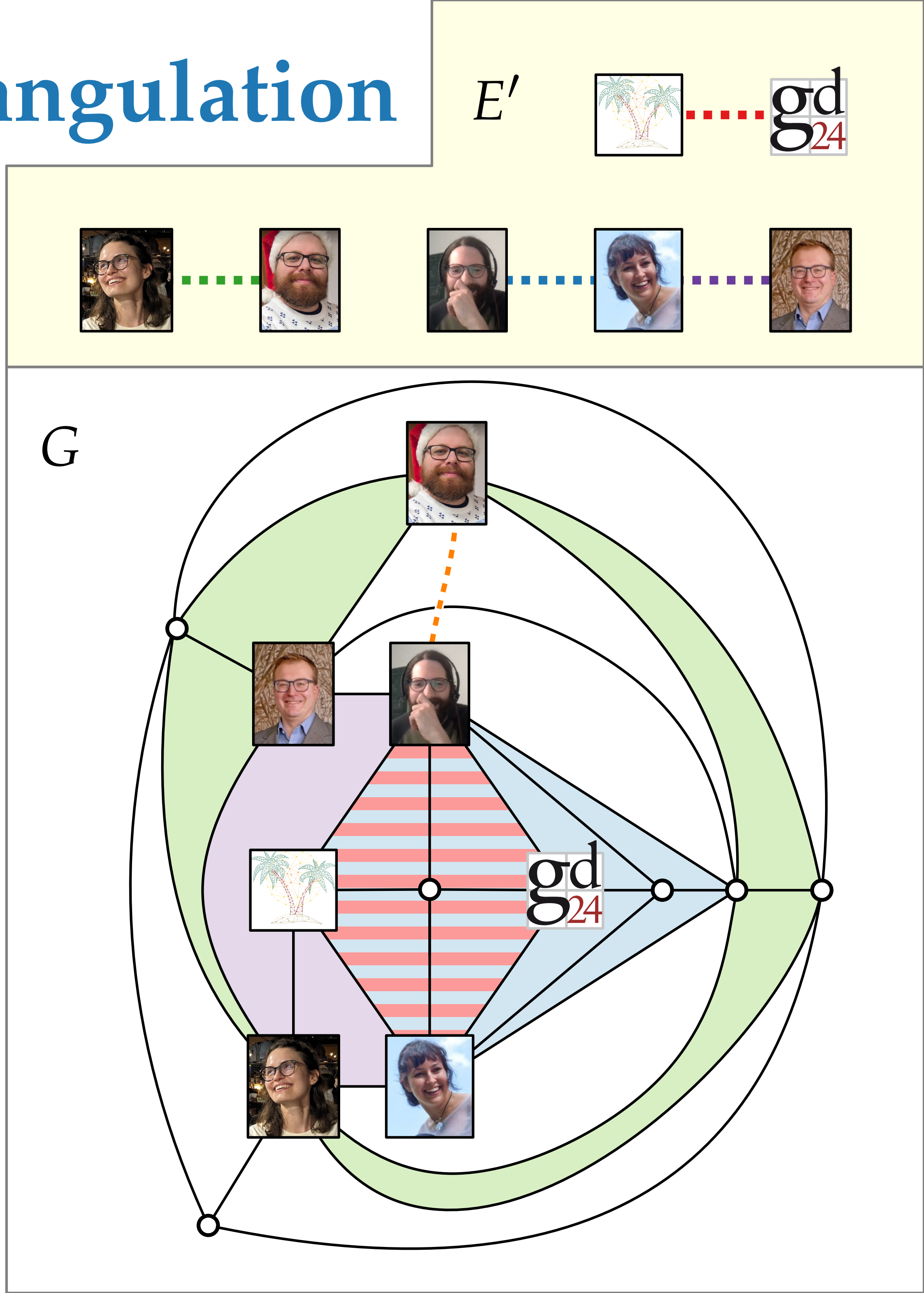

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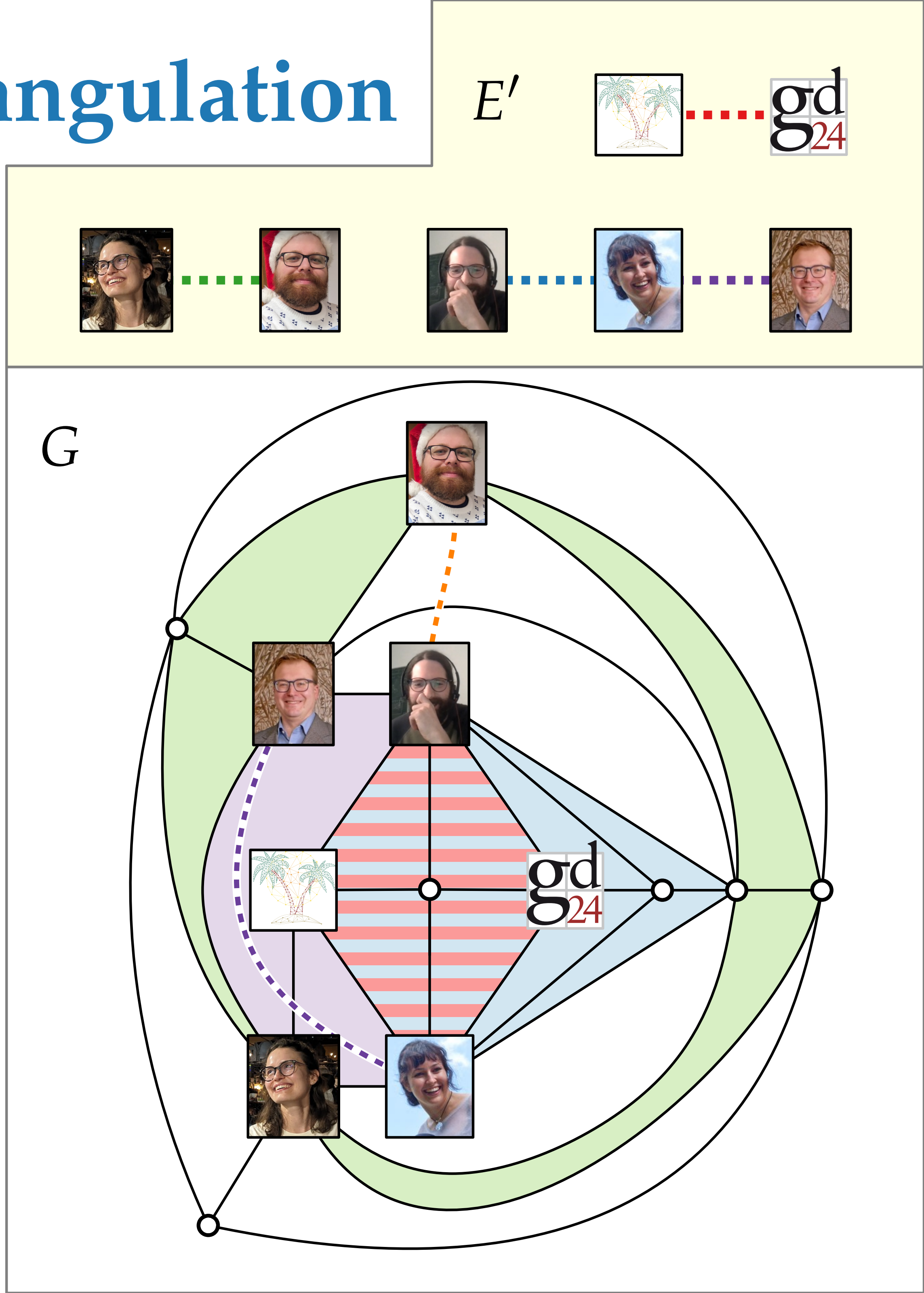

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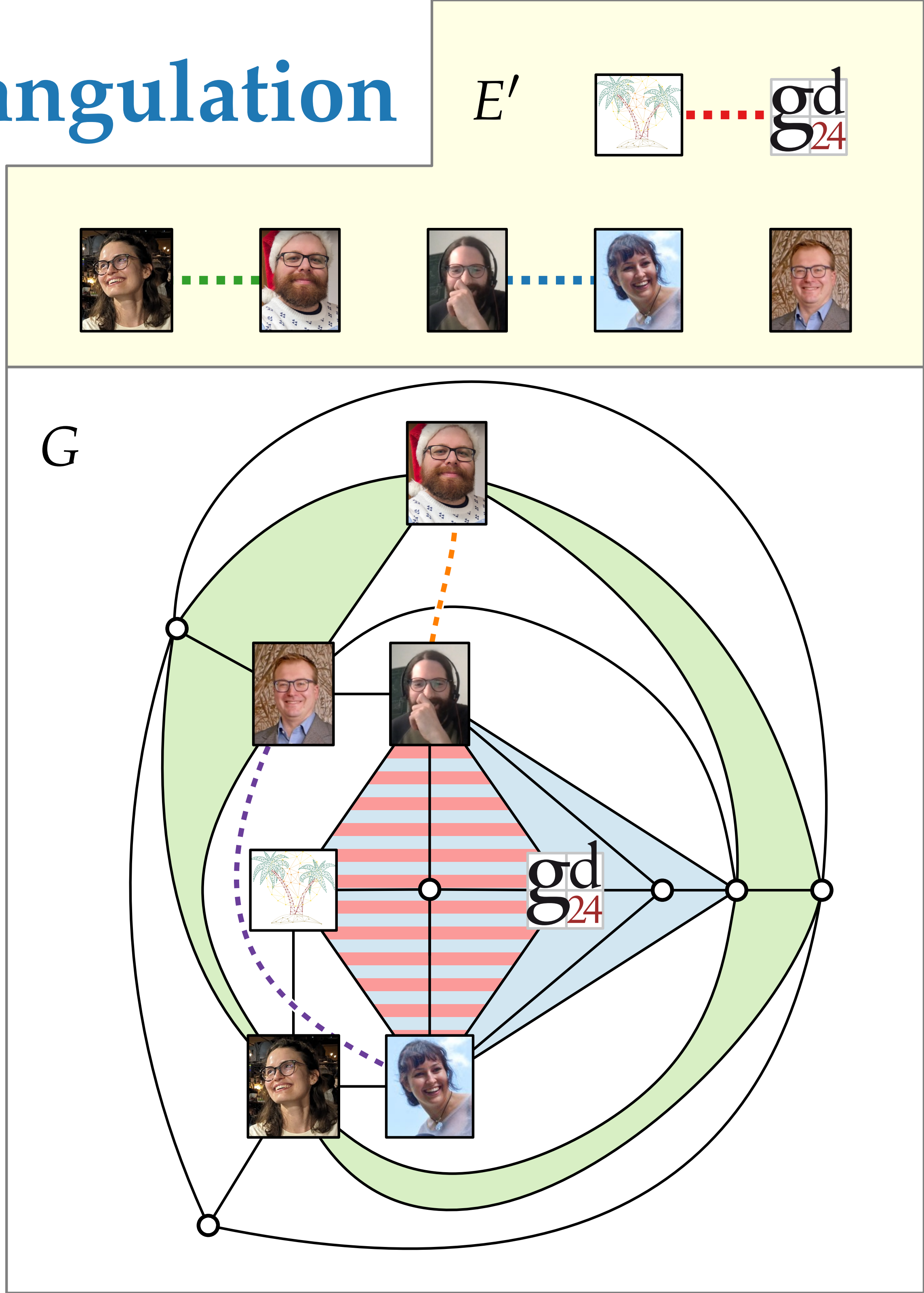

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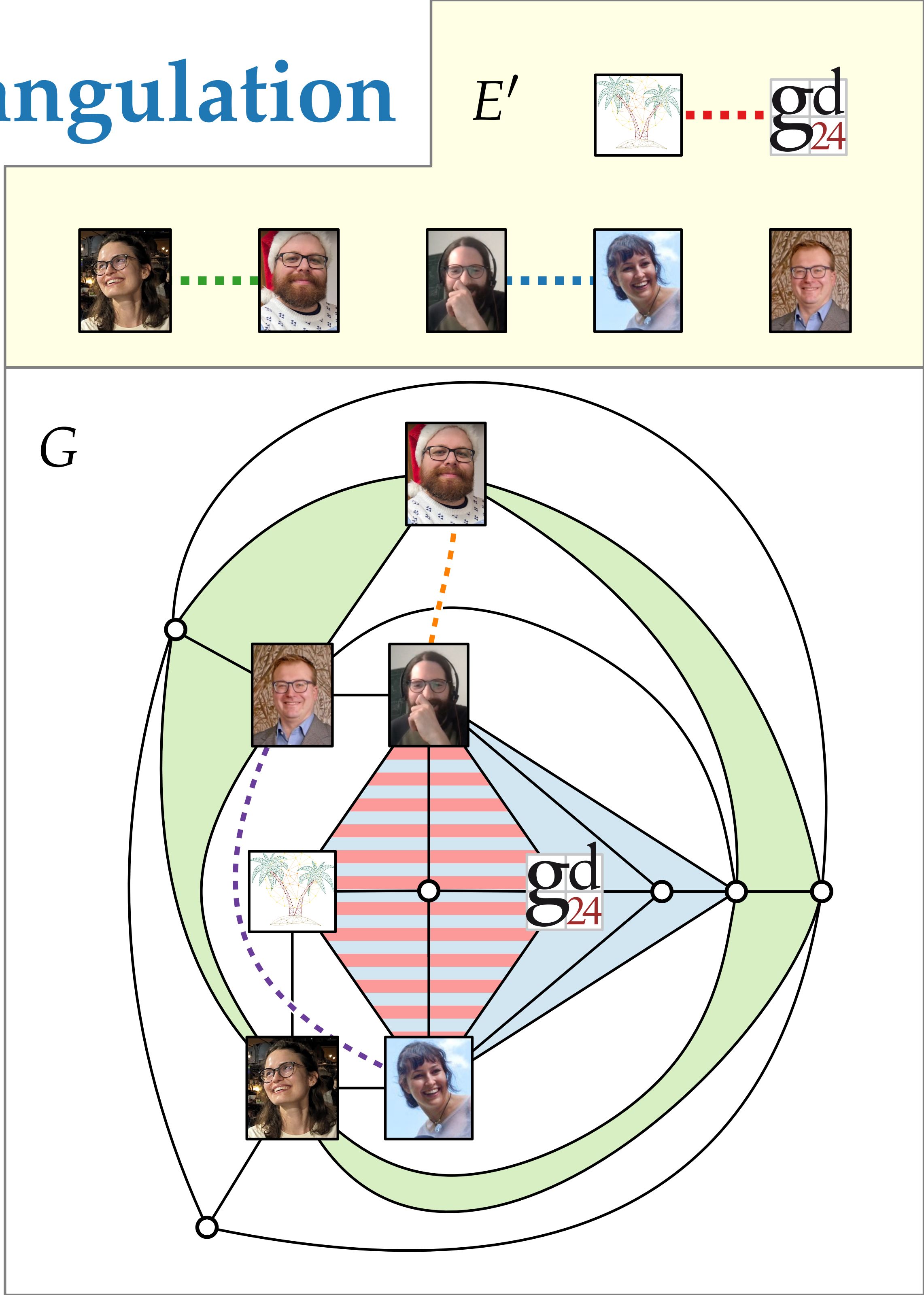

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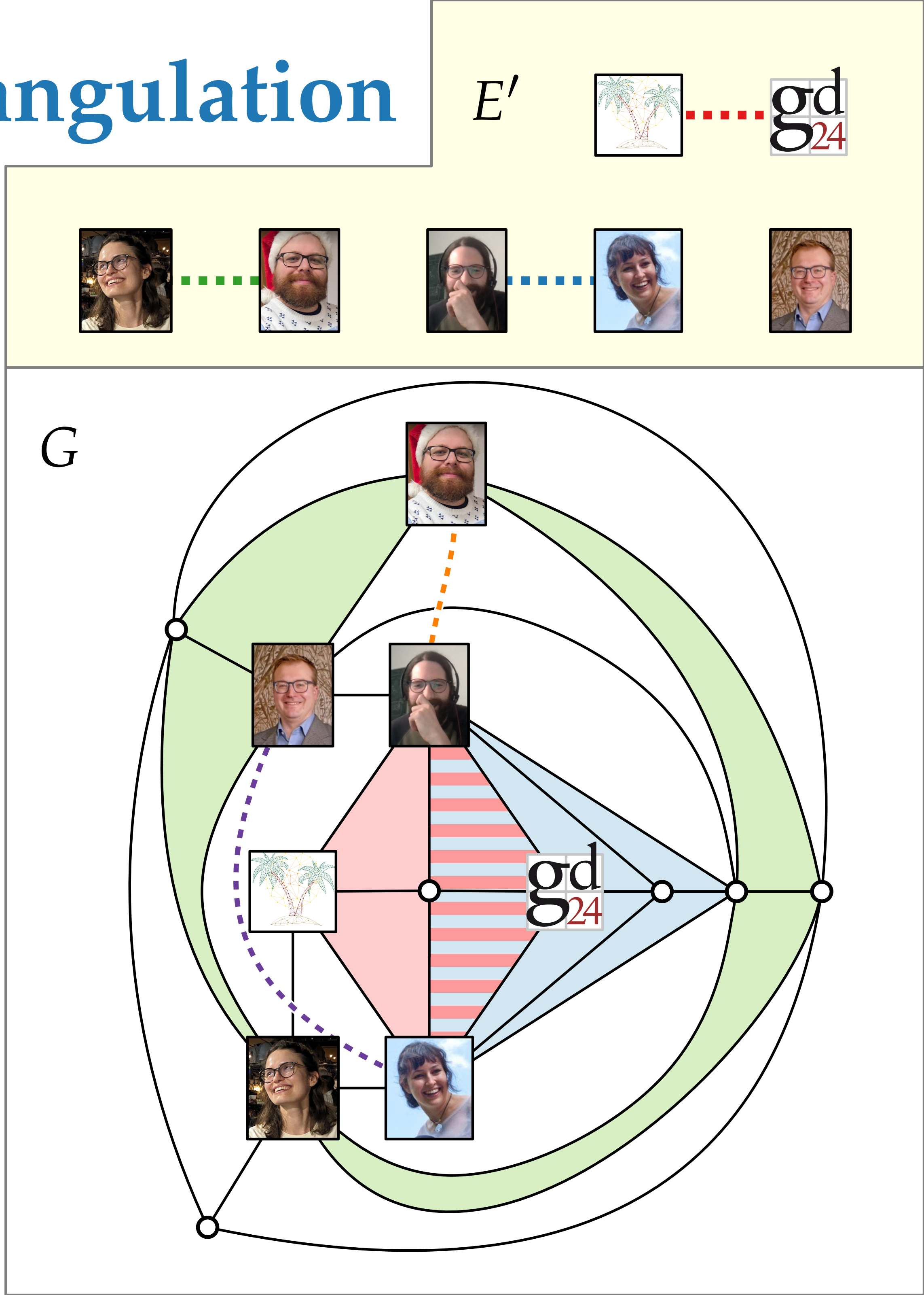

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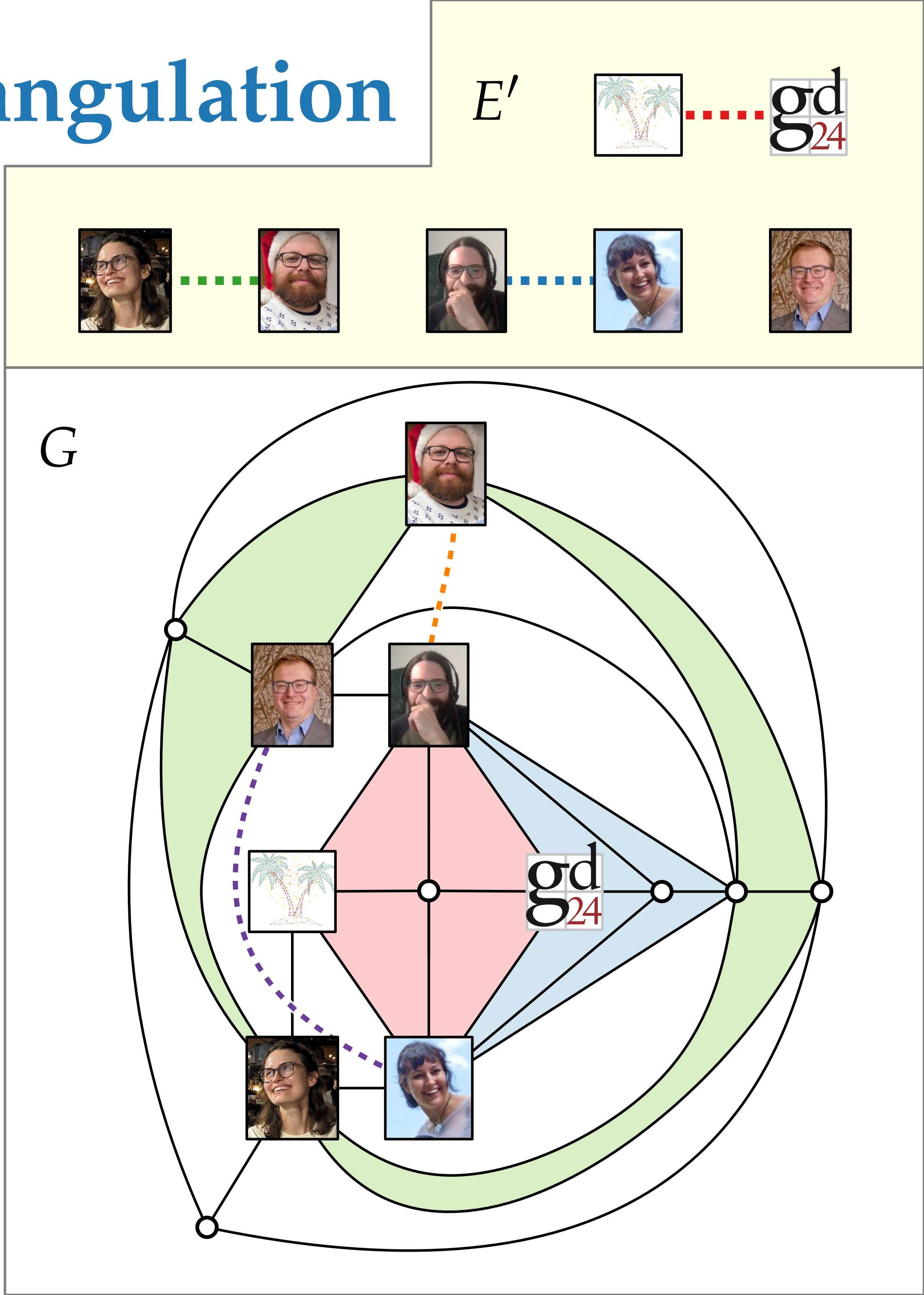

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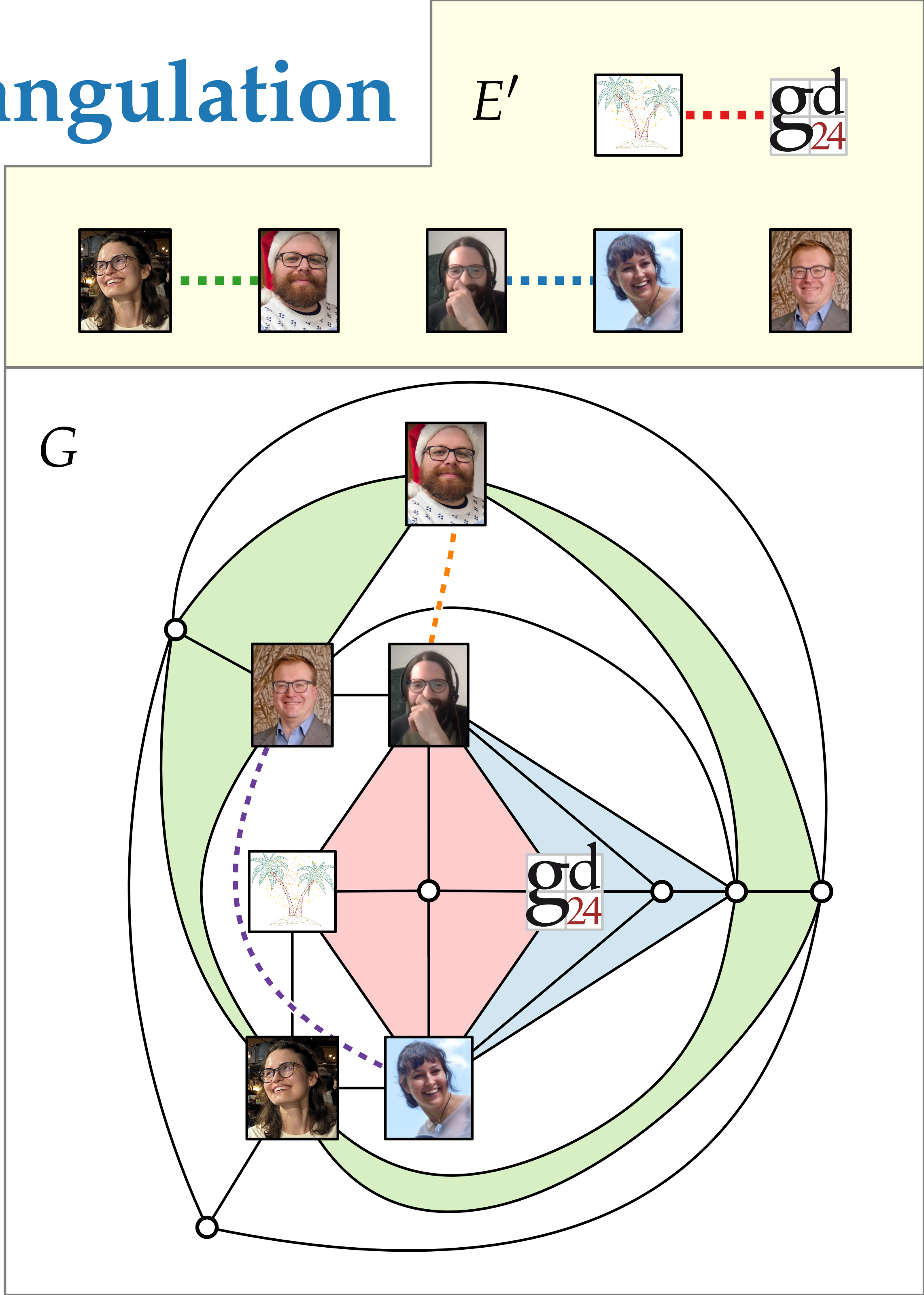
5. All edges have 2 options
   ➡ solve 2-SAT on conflict graph



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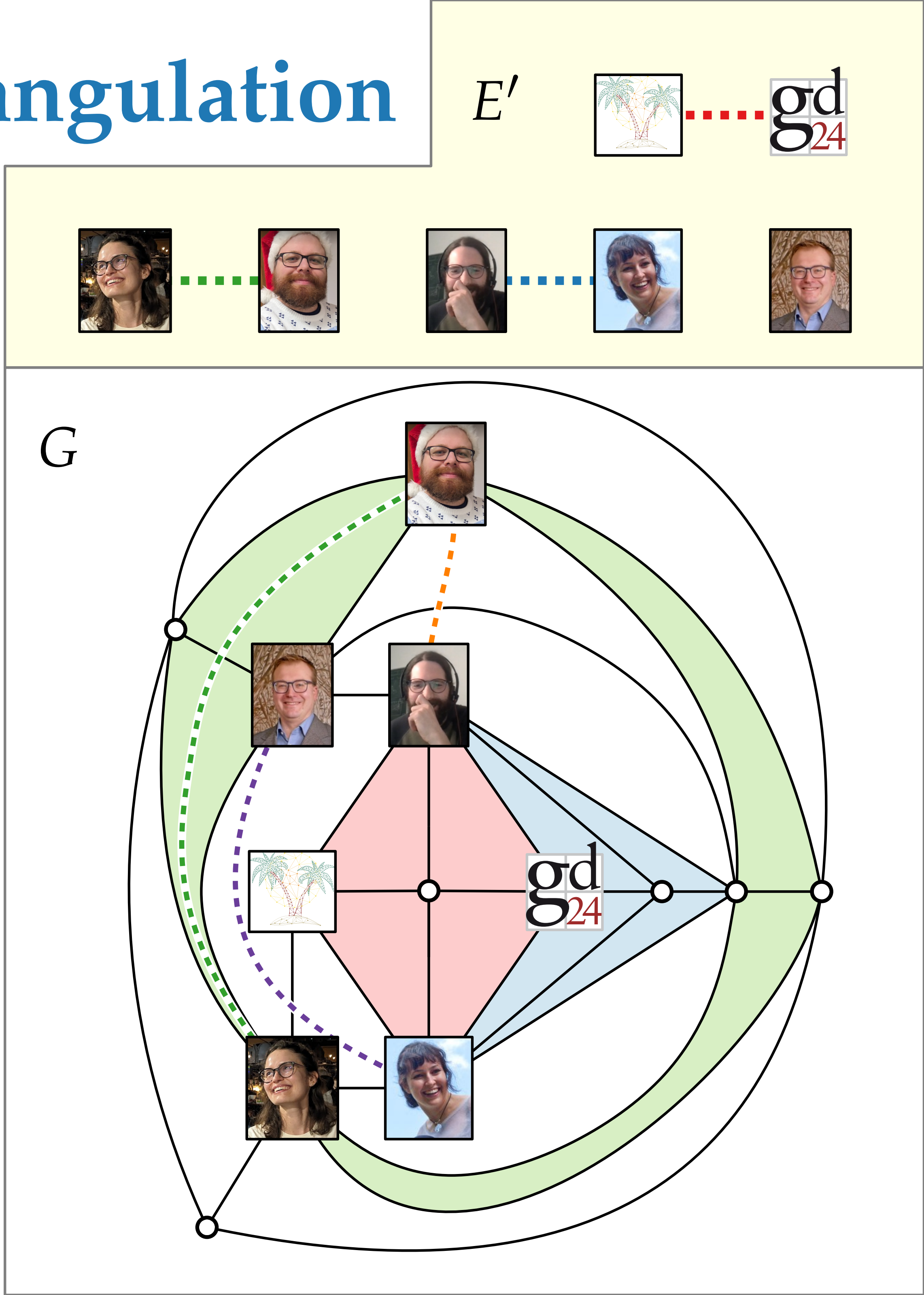

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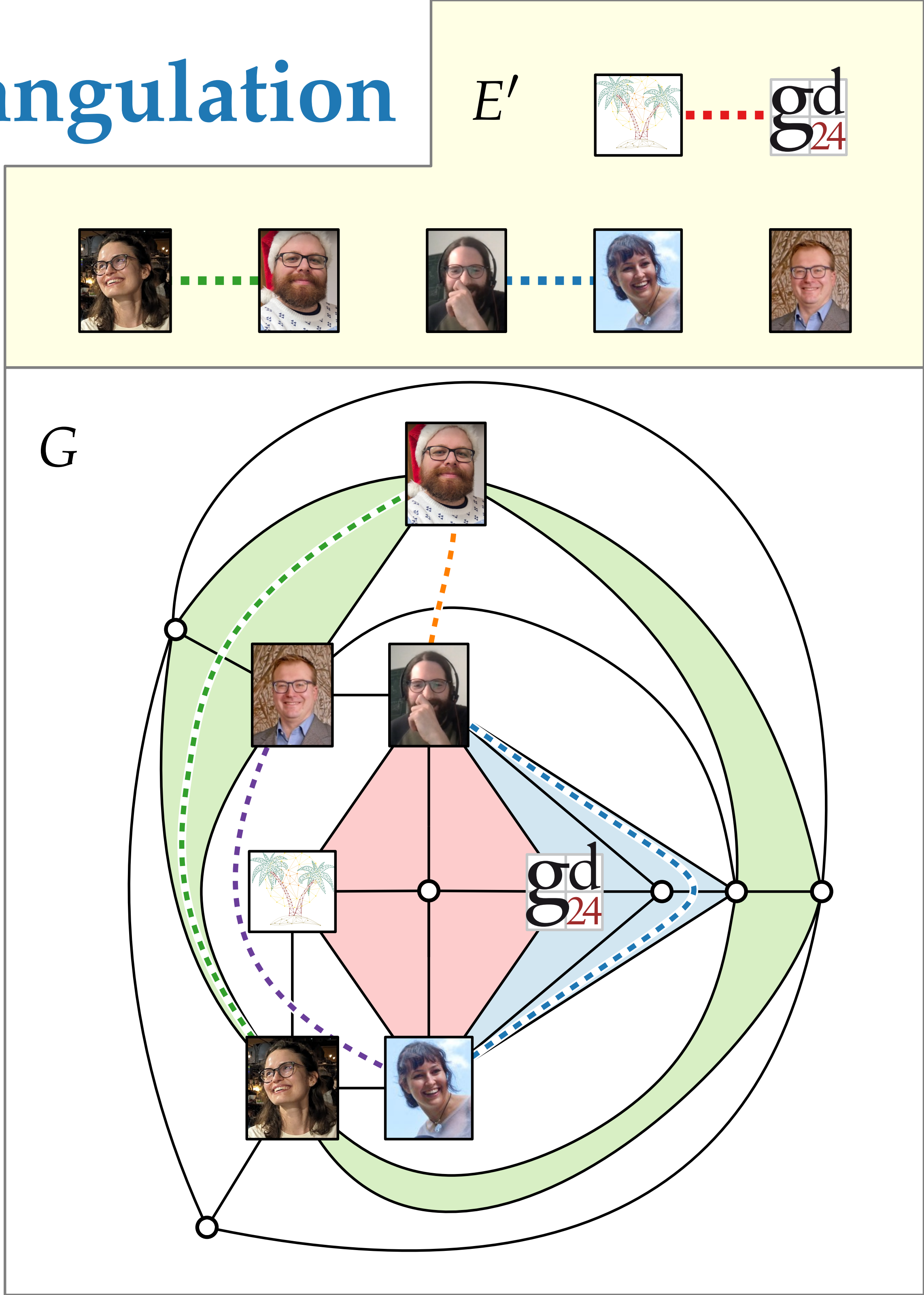

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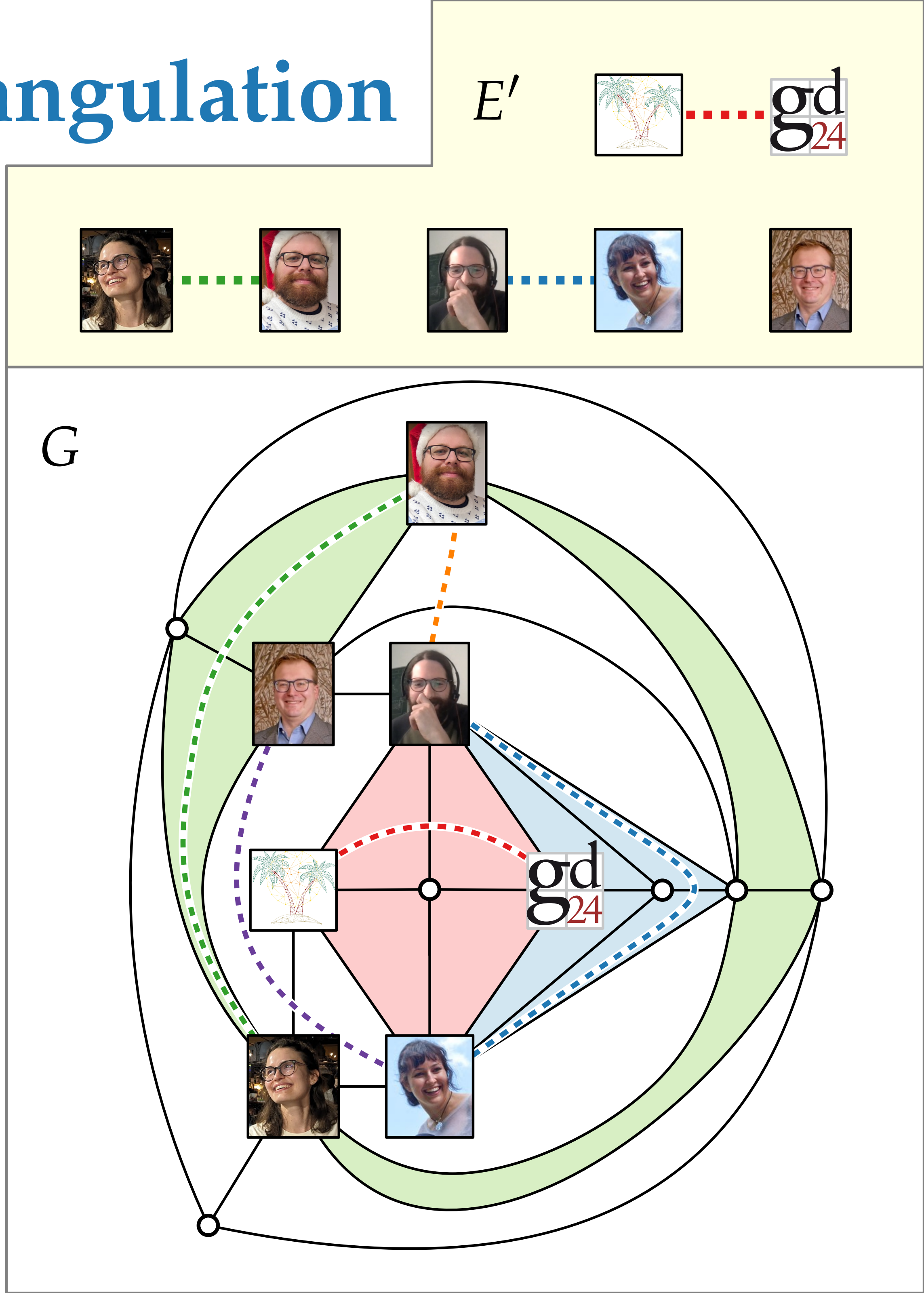

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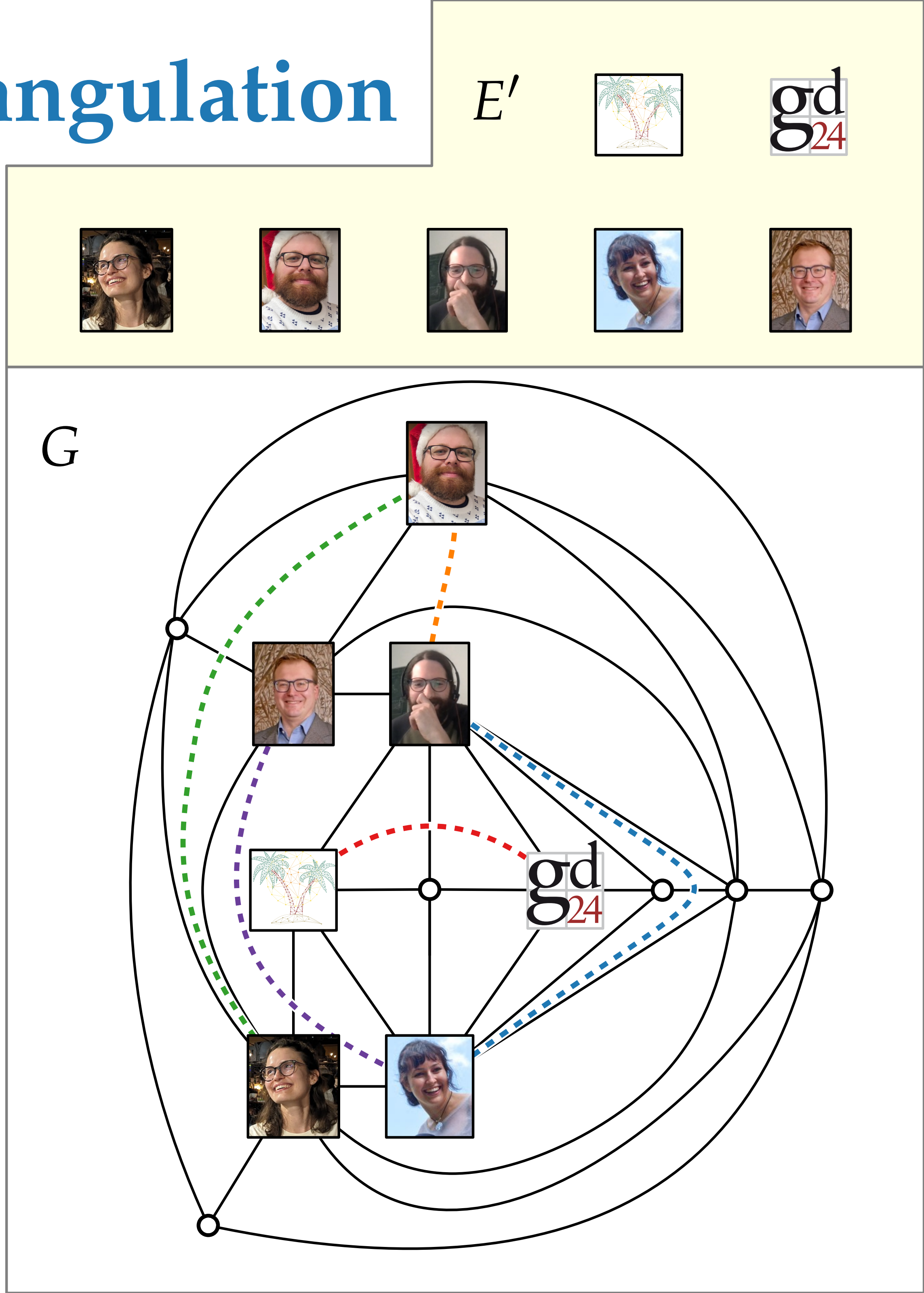

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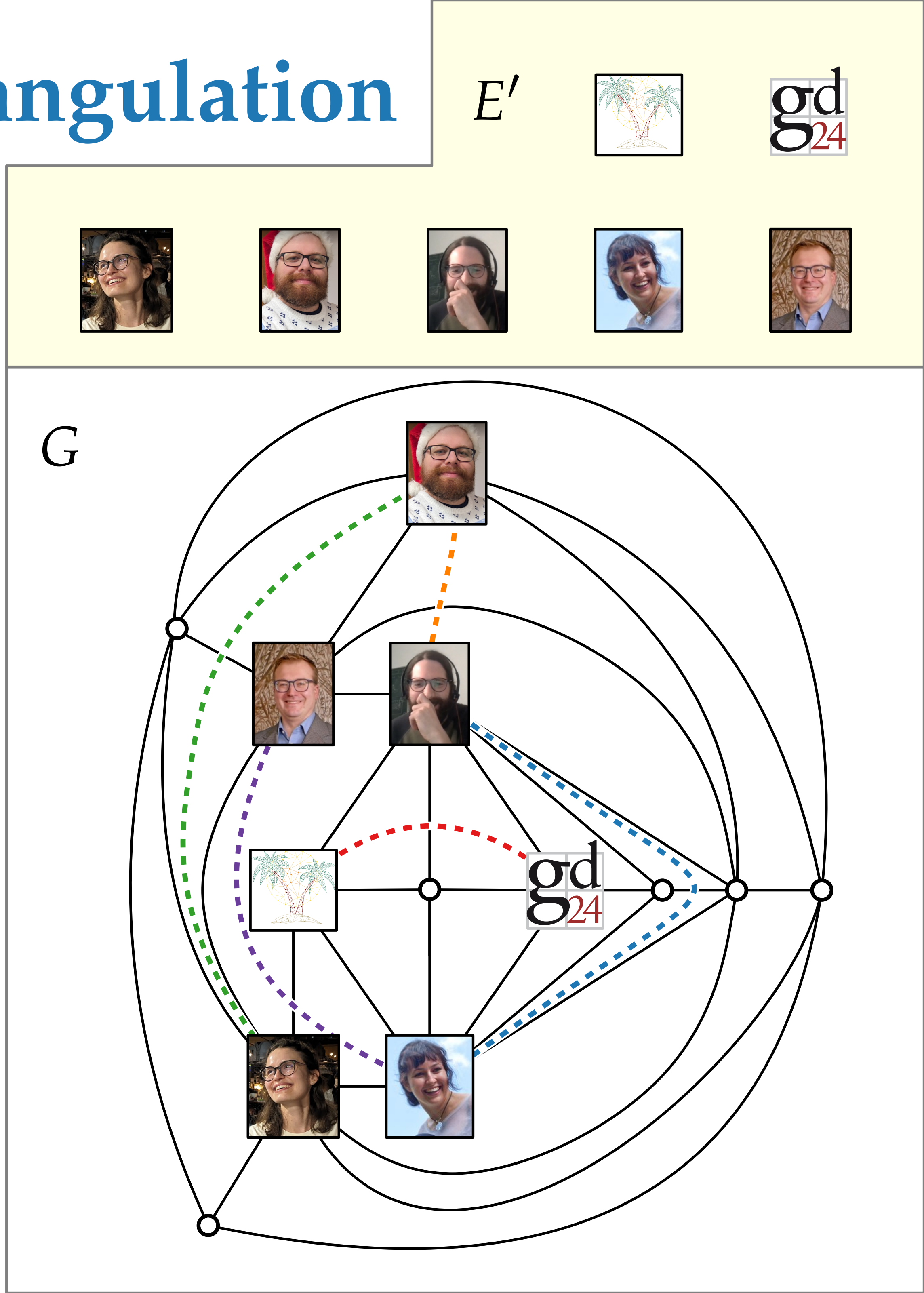

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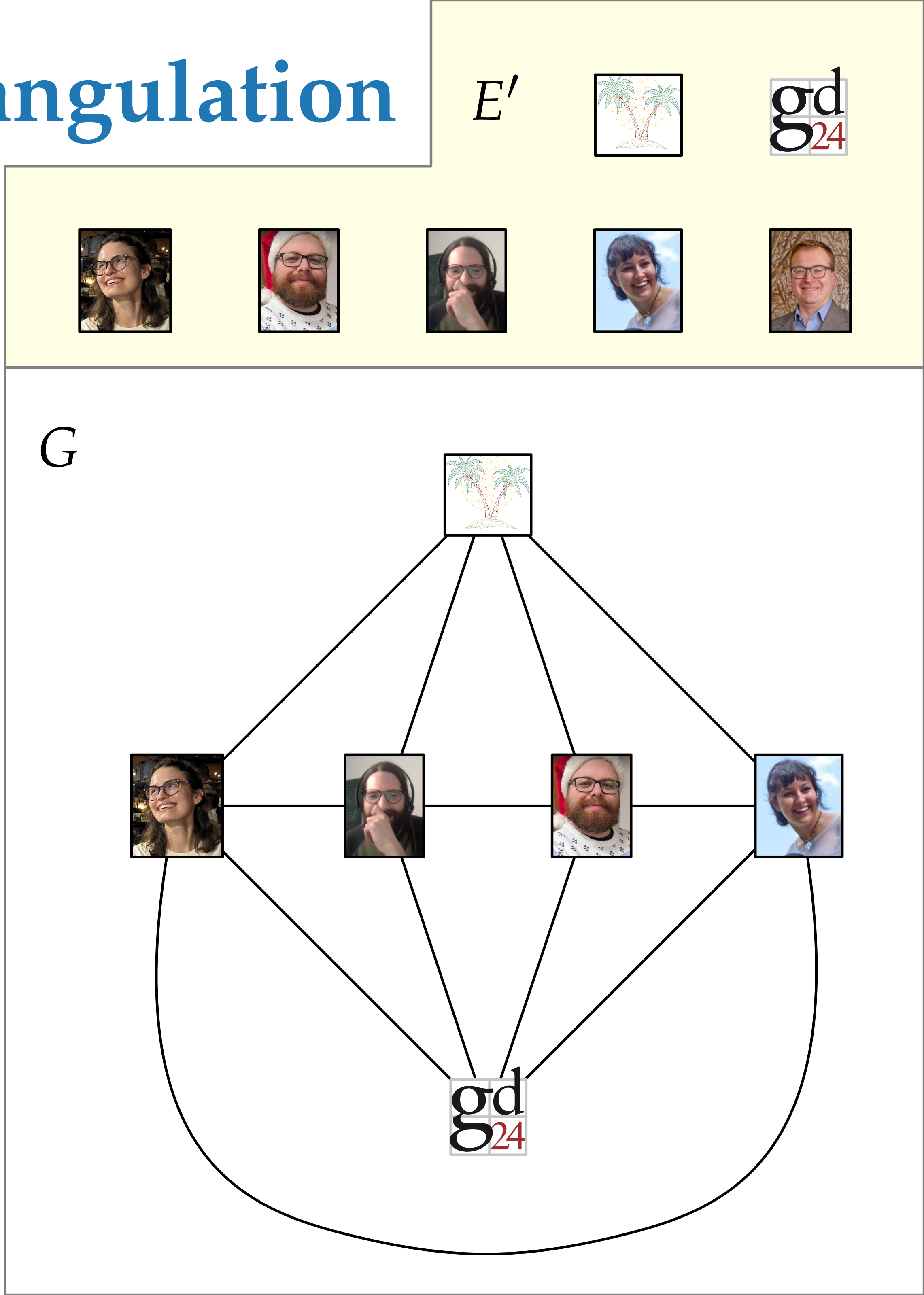

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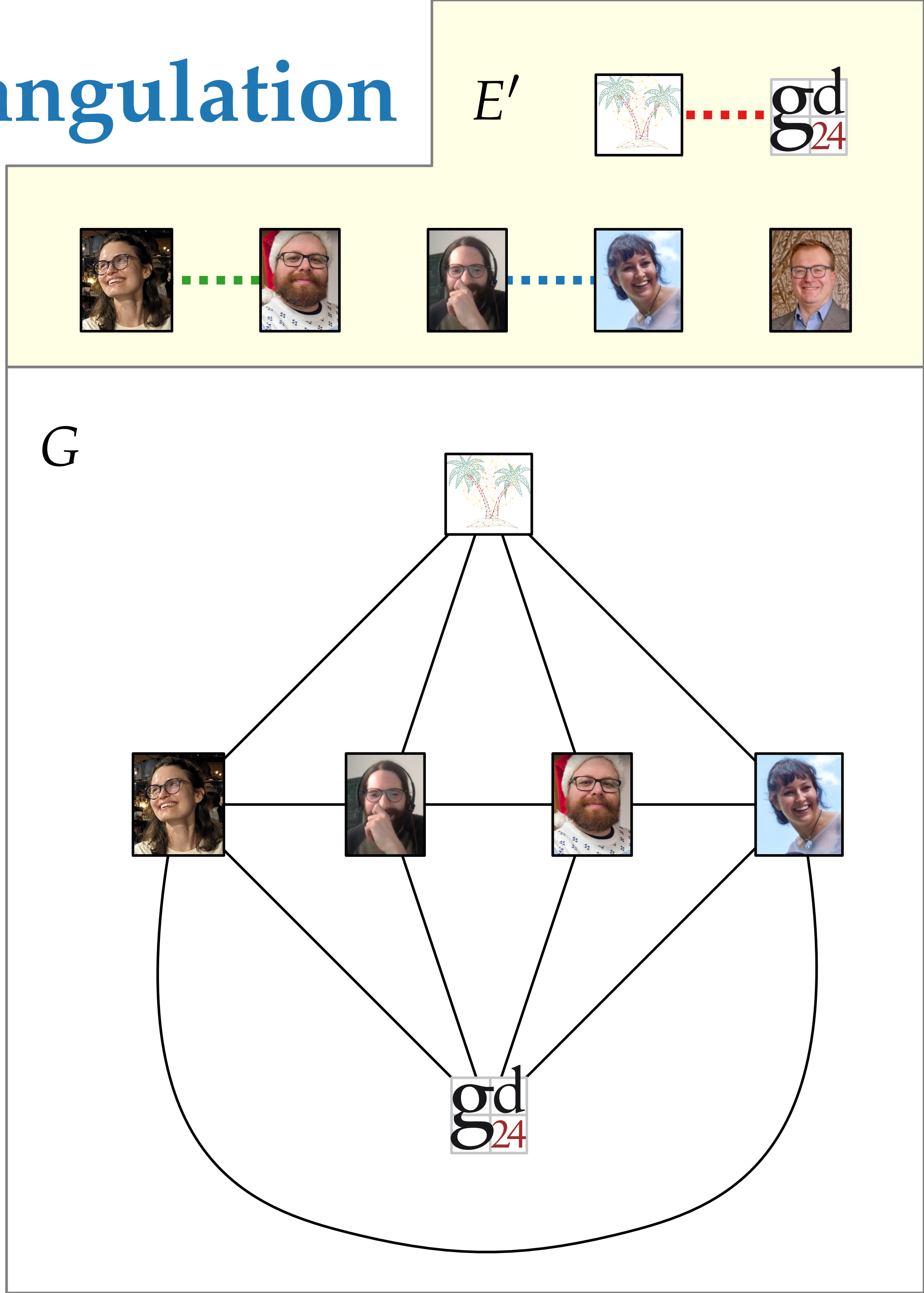

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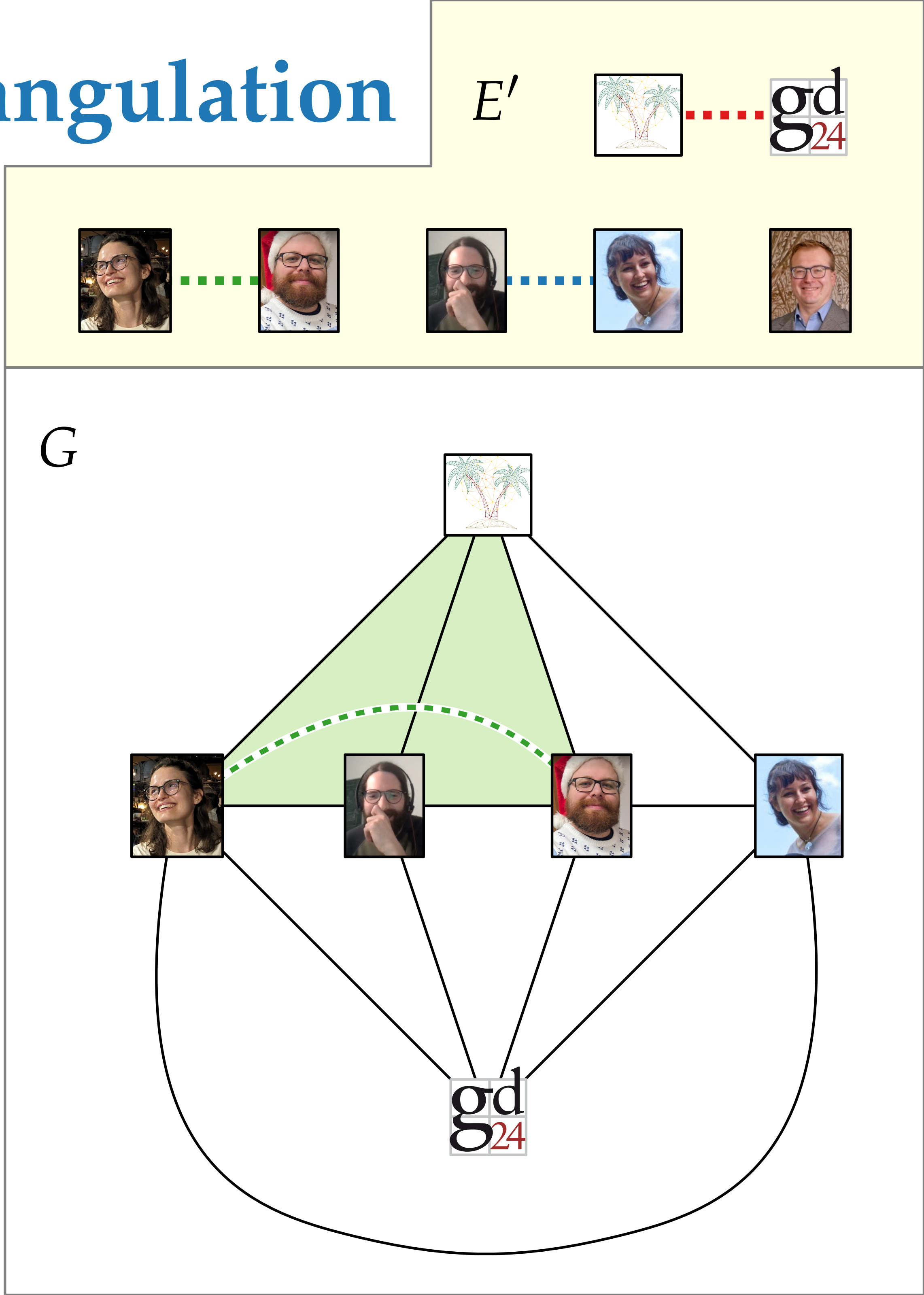

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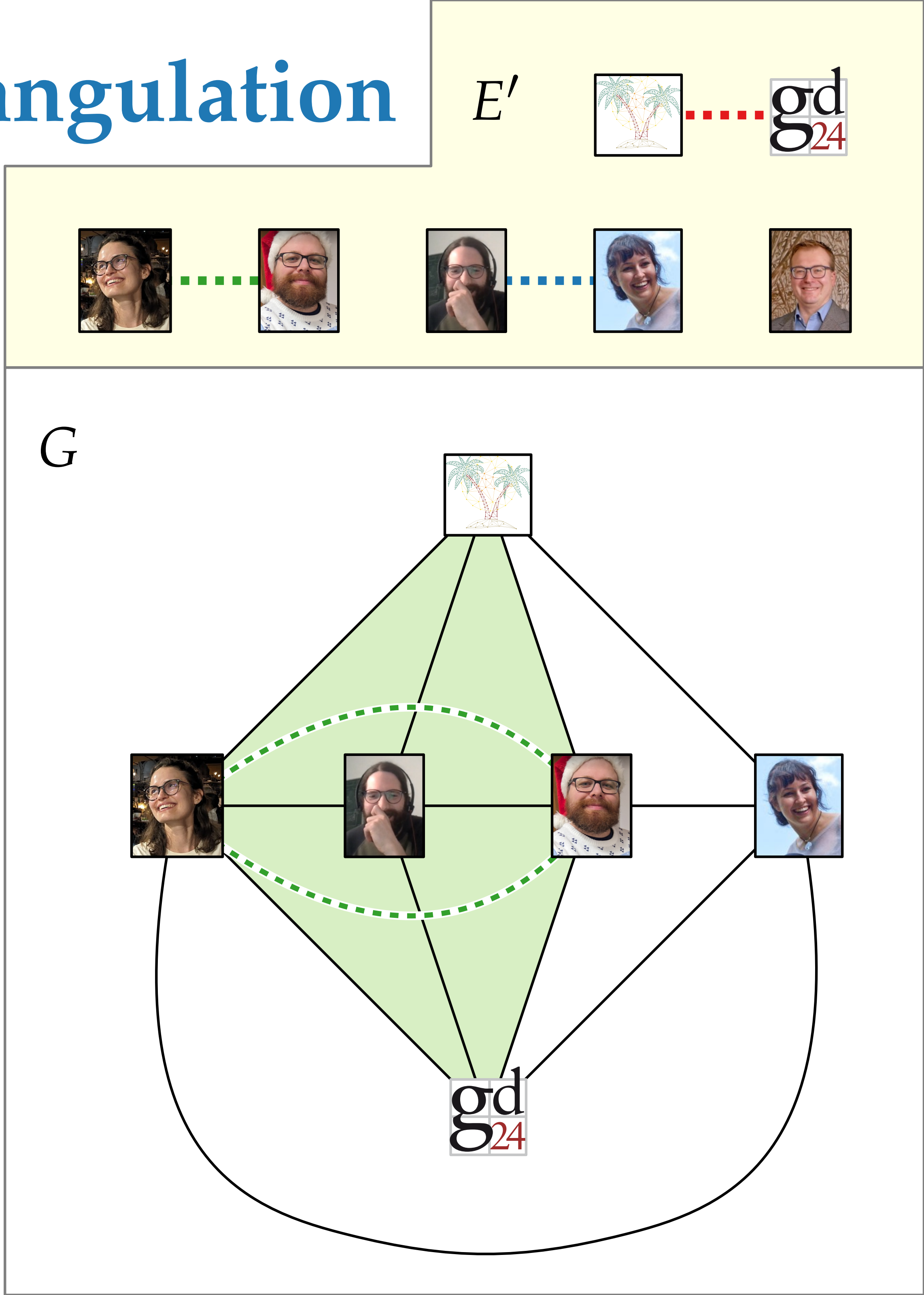

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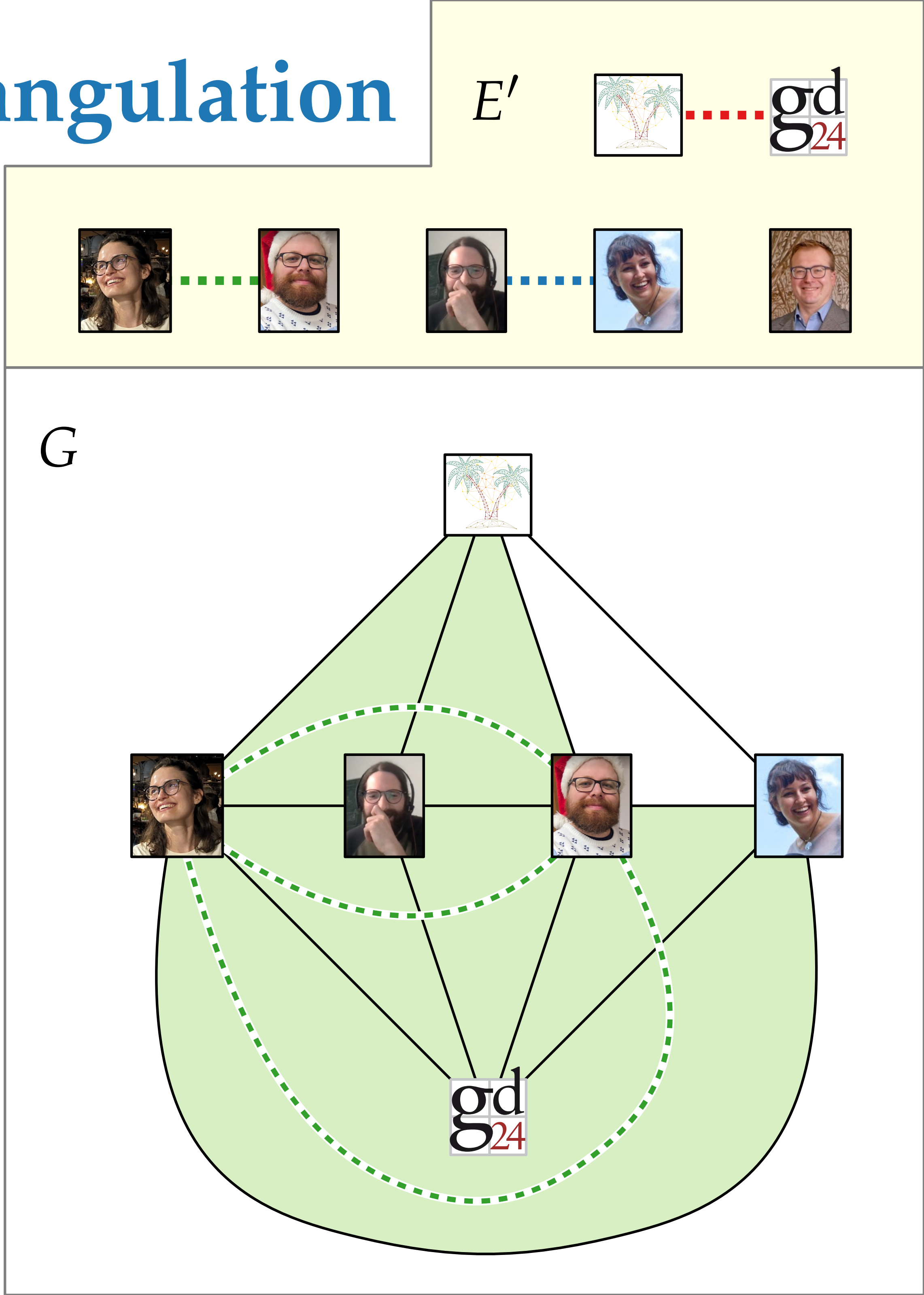

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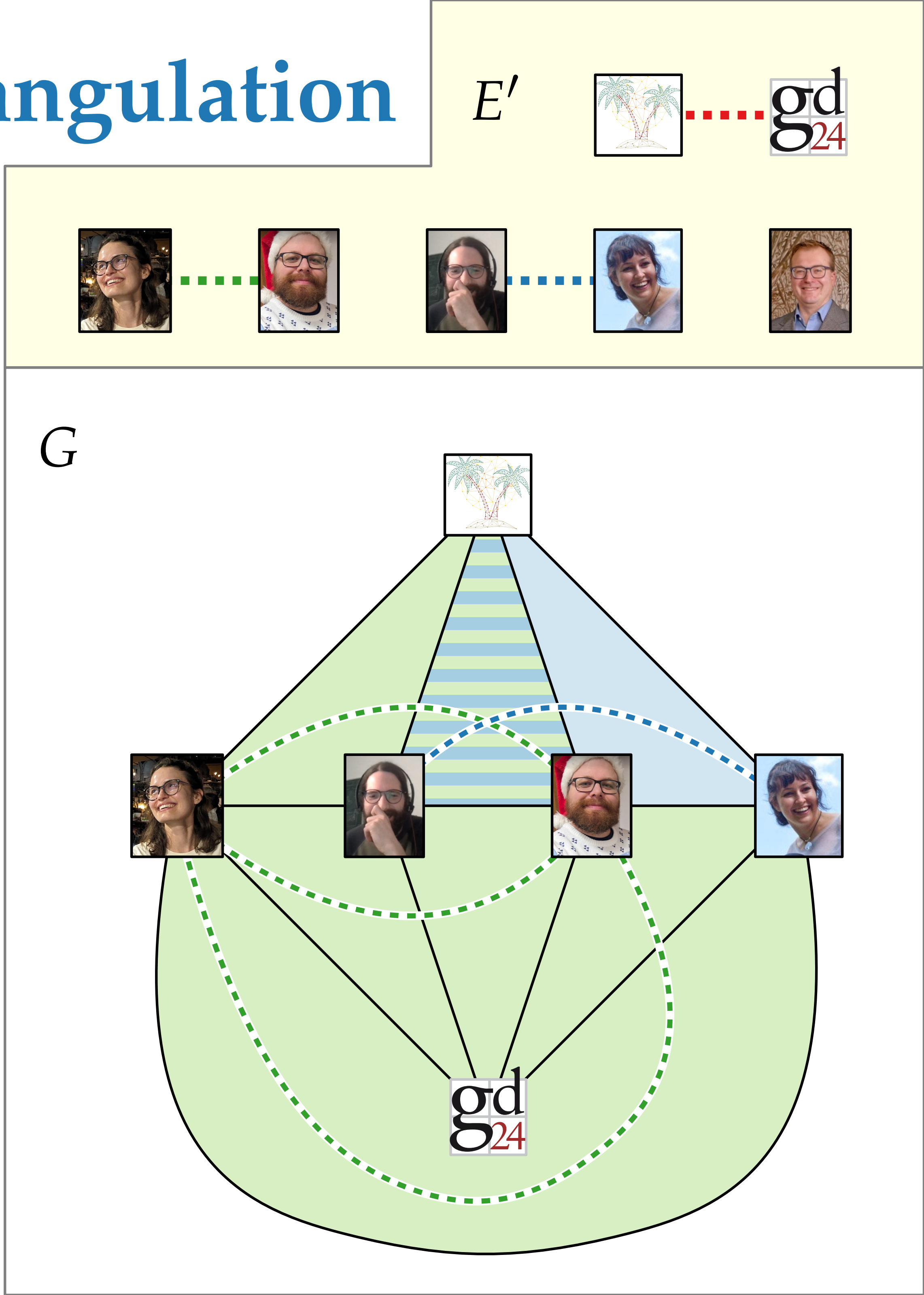

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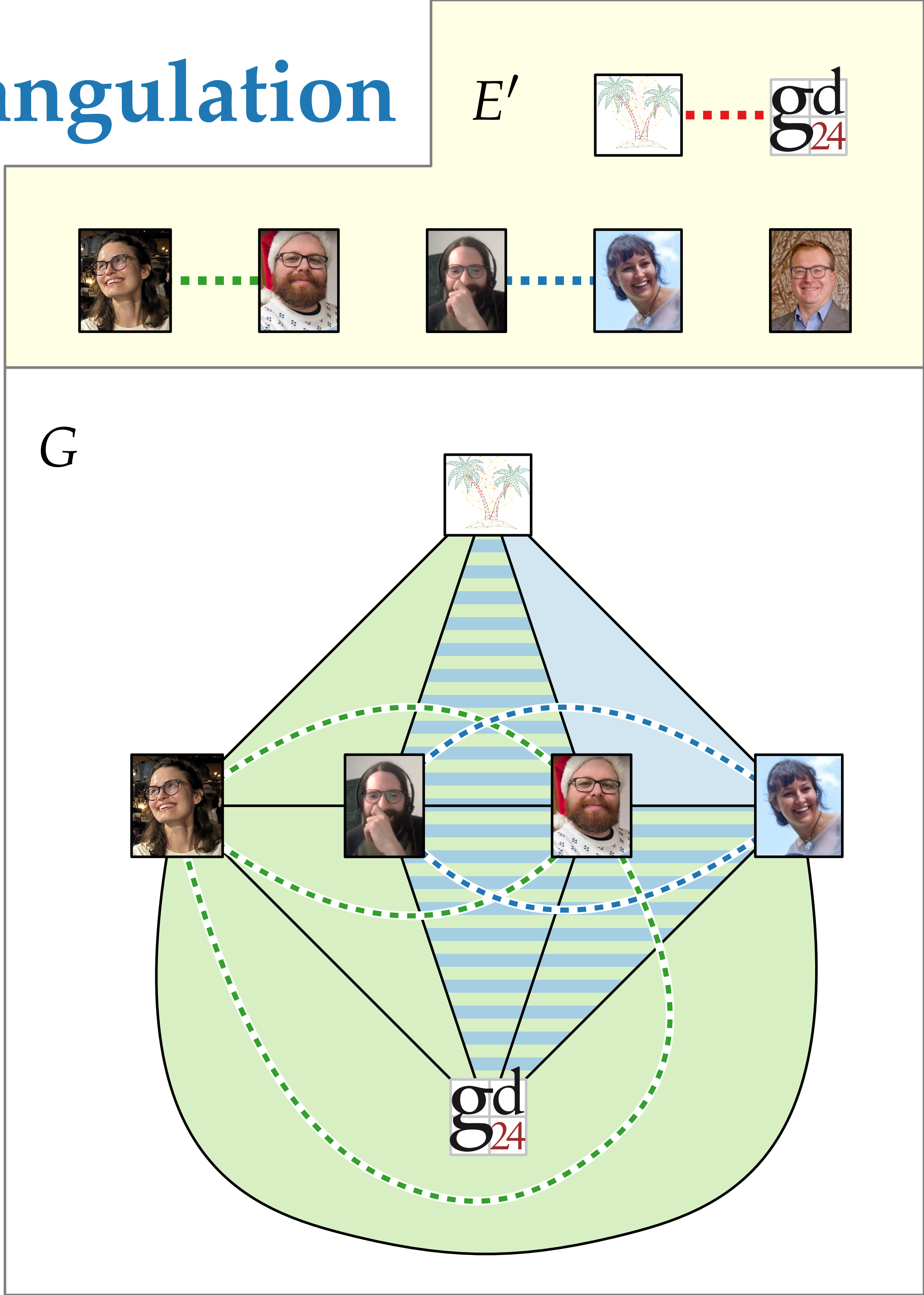

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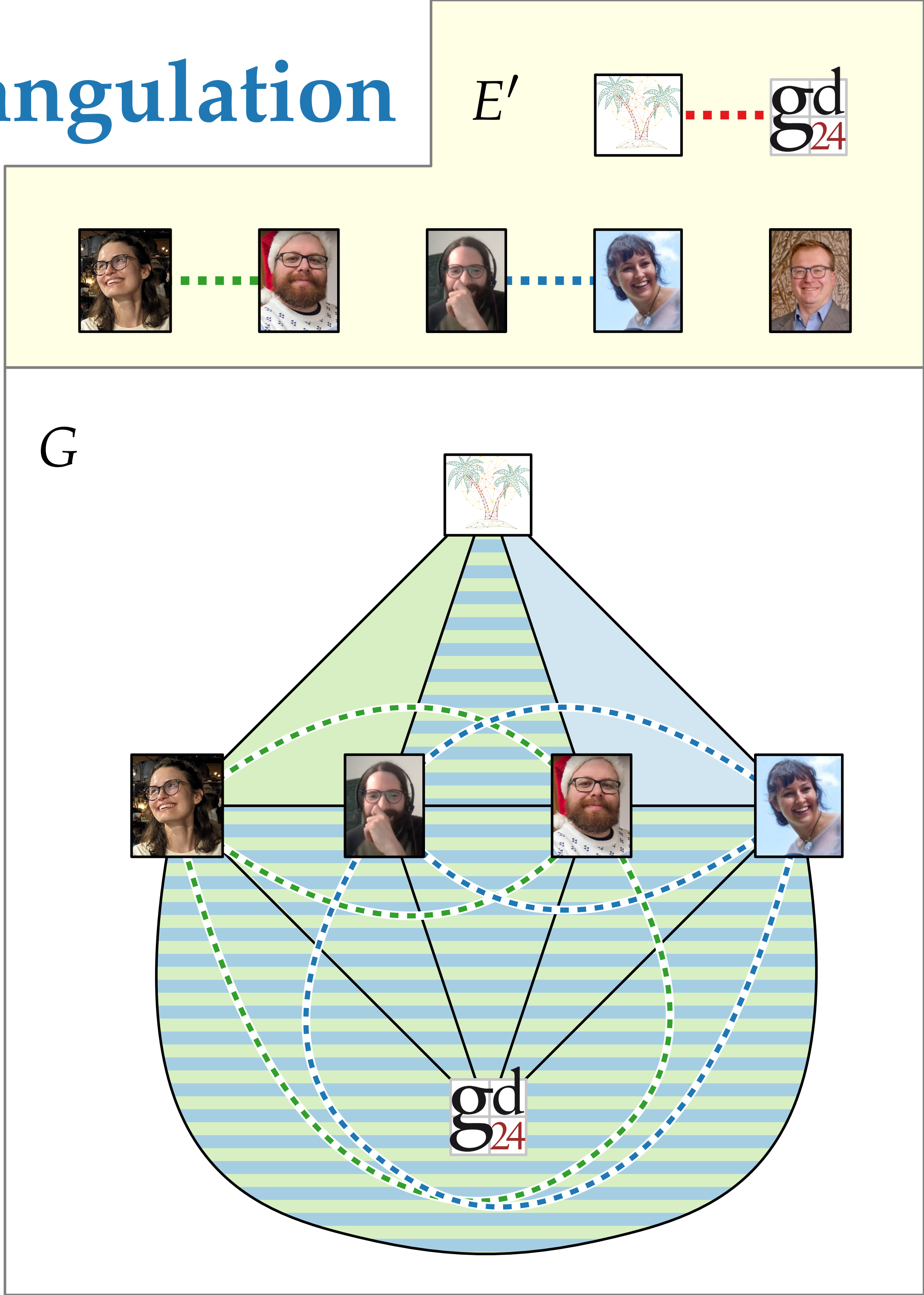

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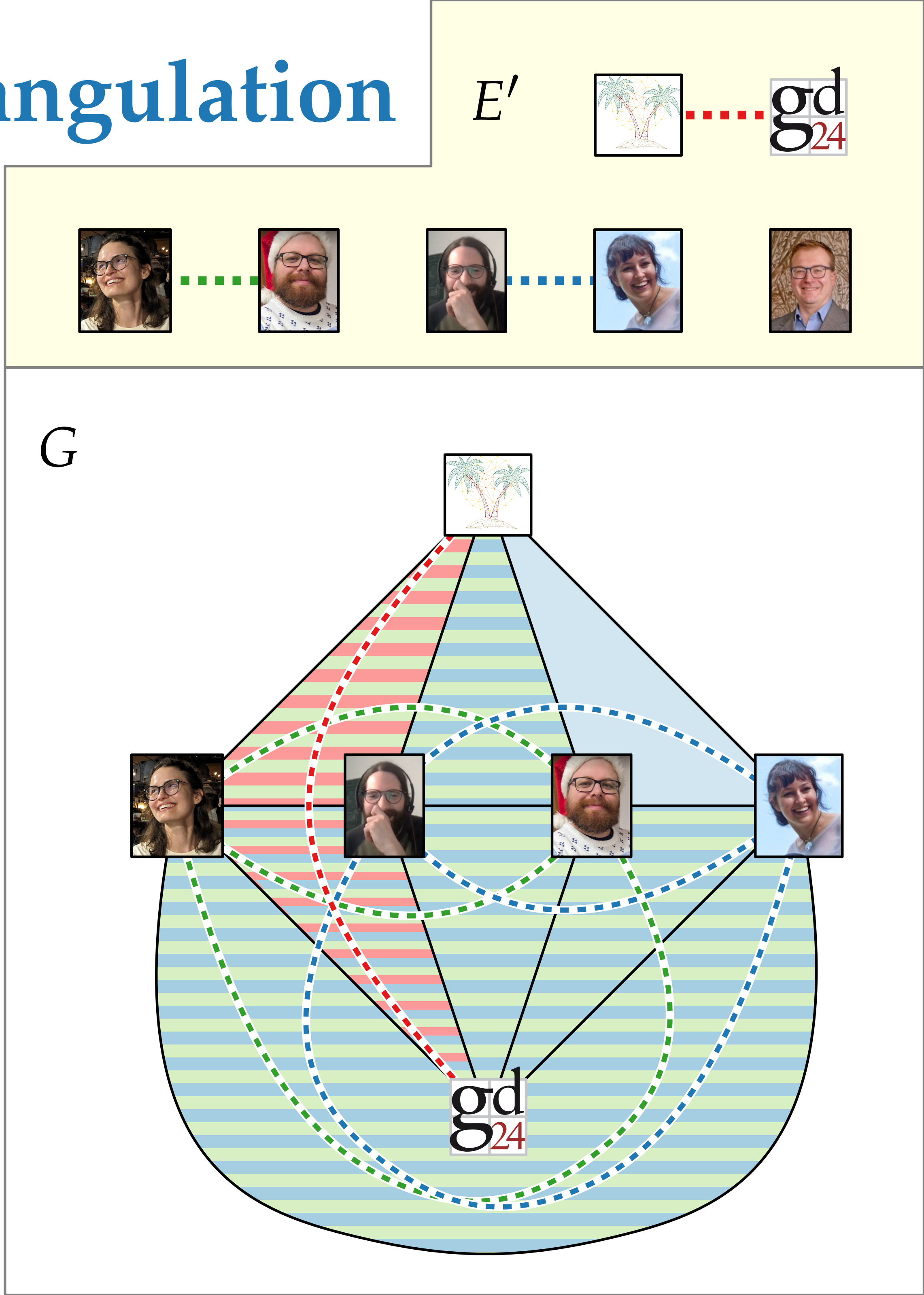

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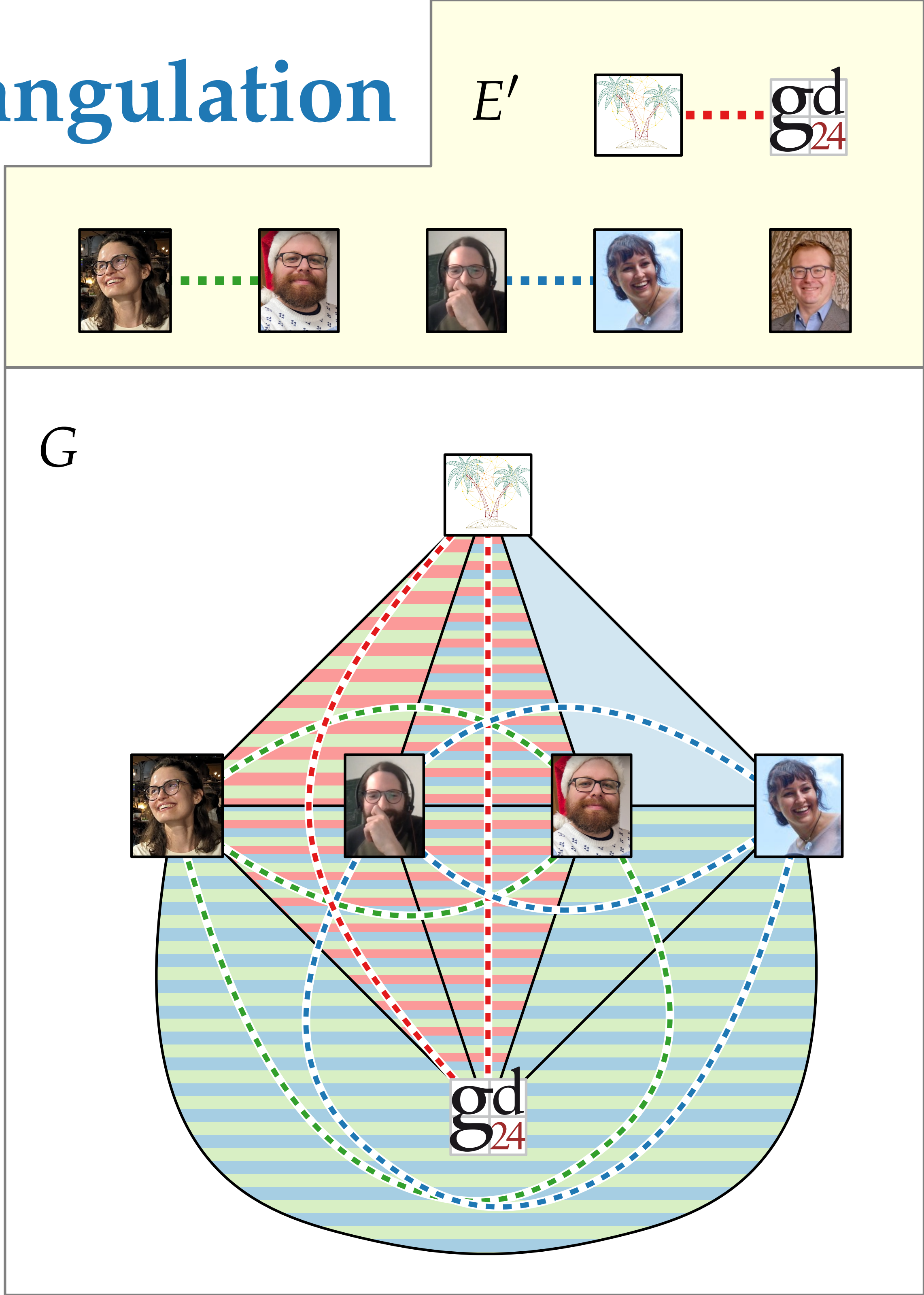

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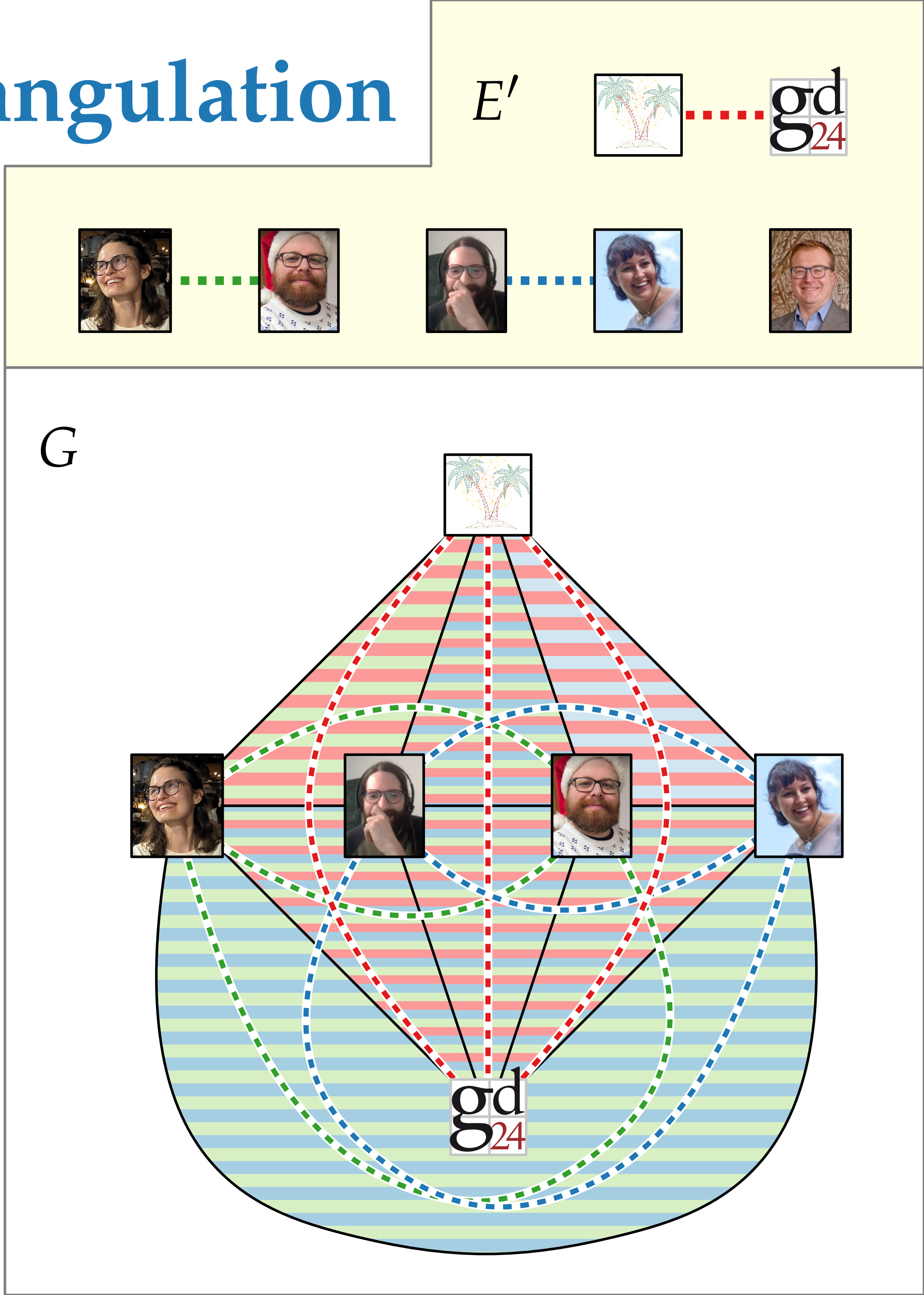

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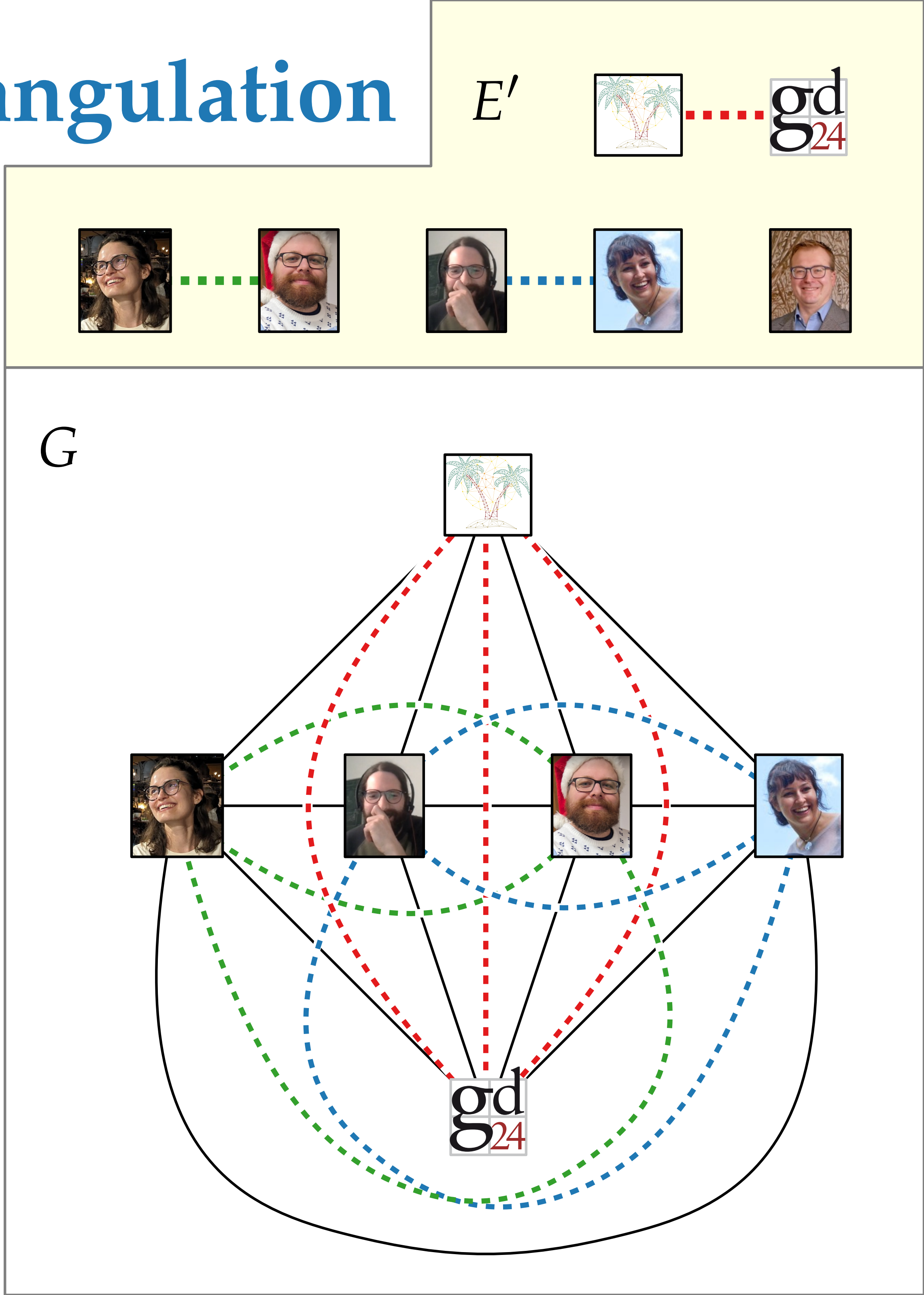

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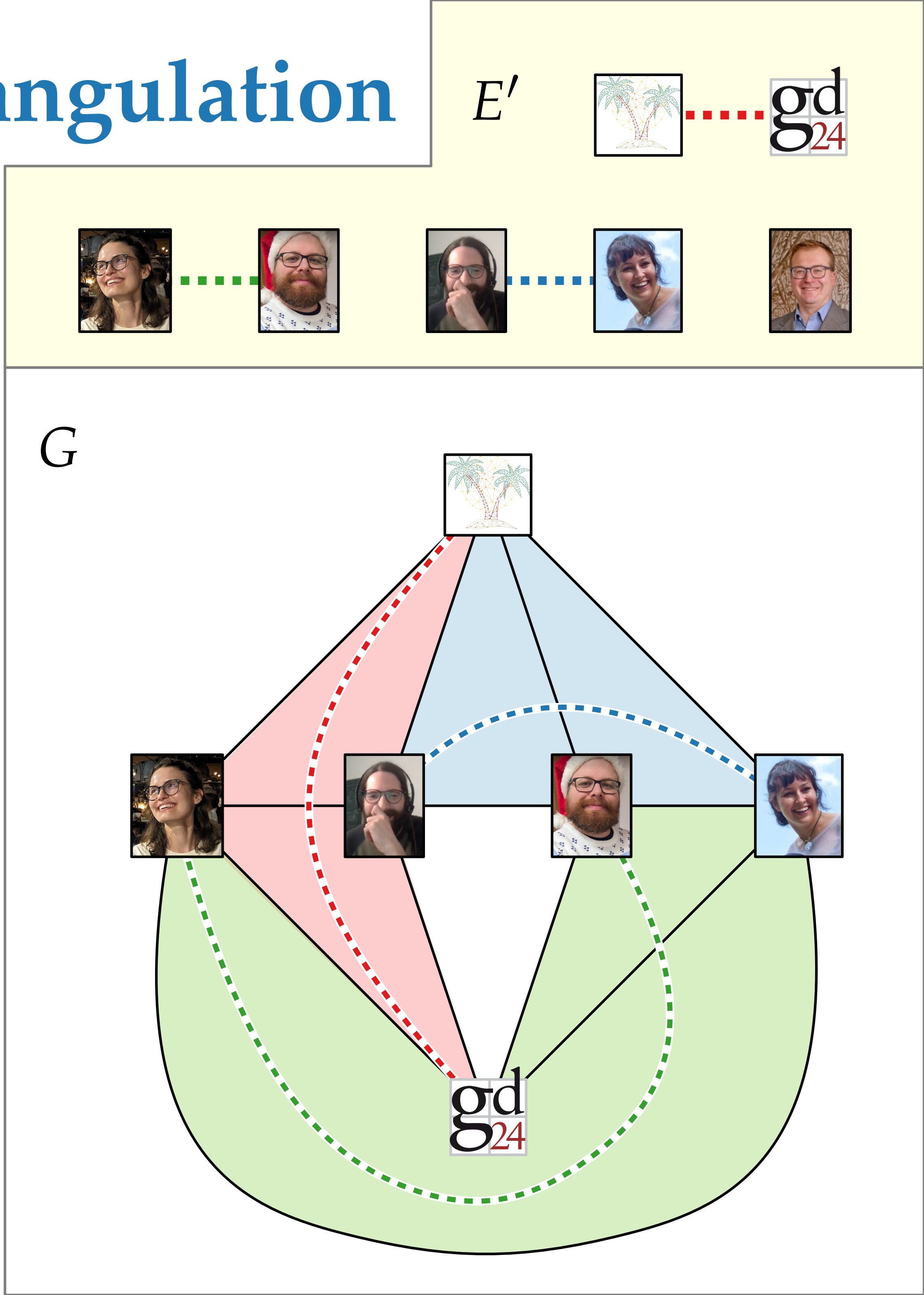

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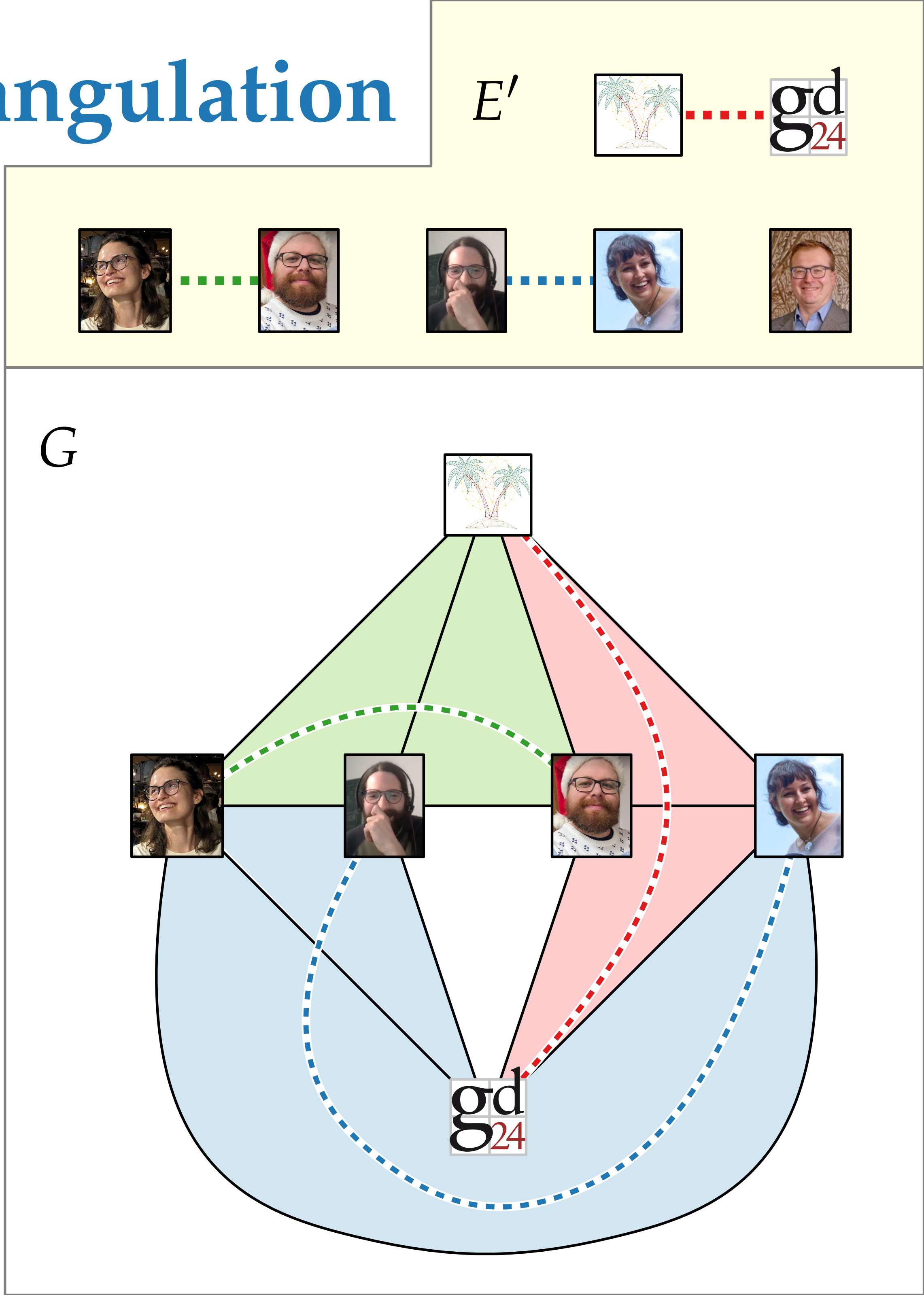

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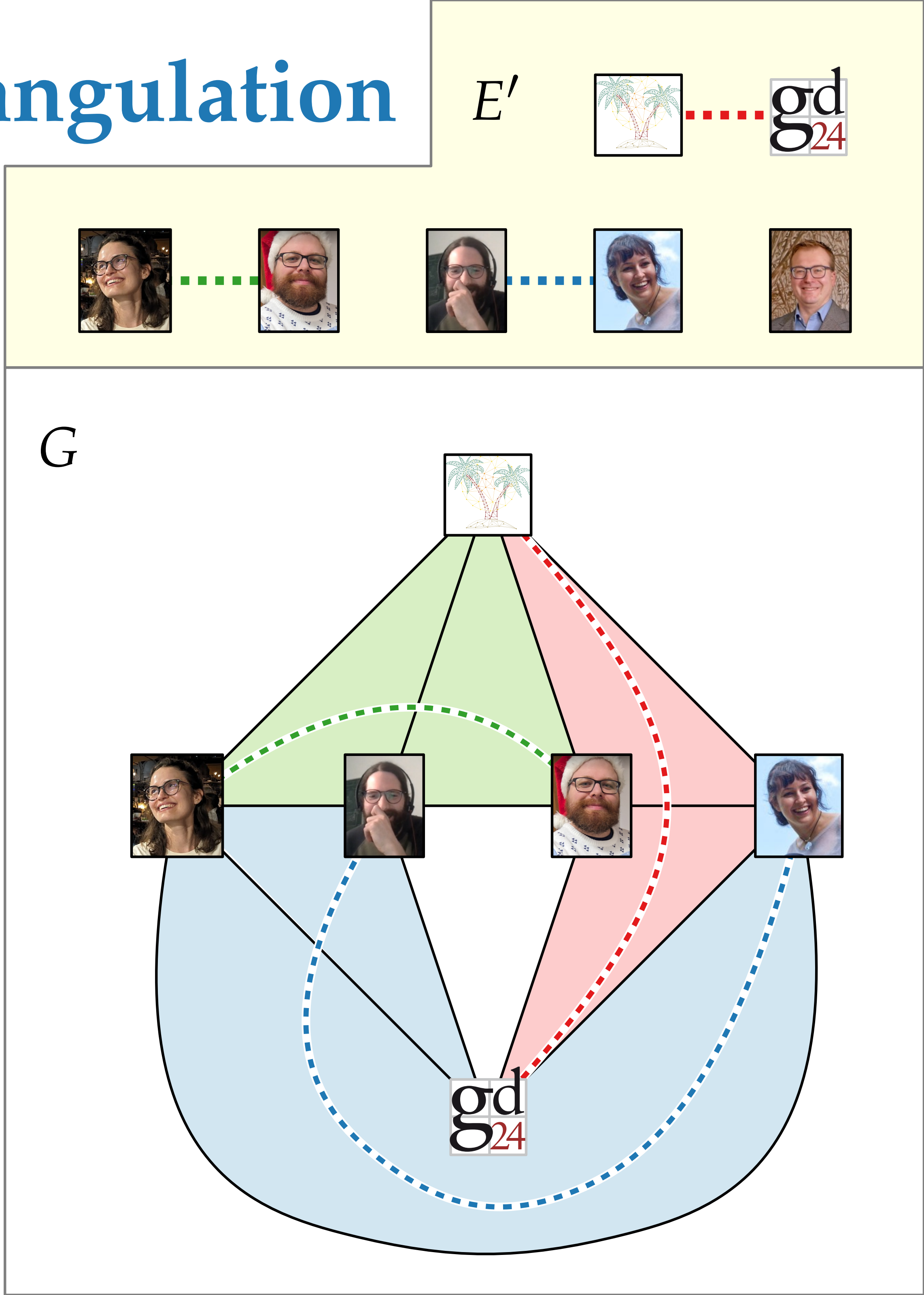
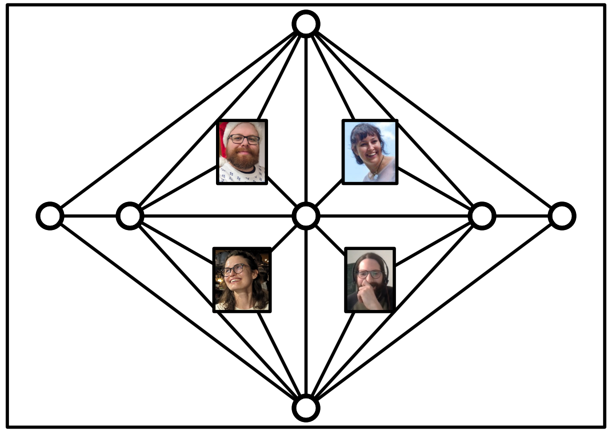
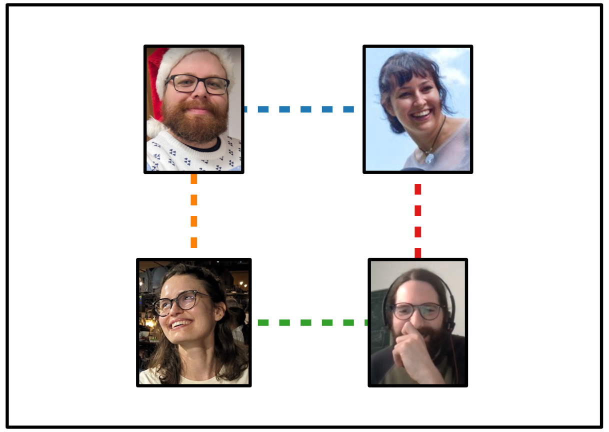


Theorem.
1-plane insertion into a plane triangulation can be solved in $O(n)$ time.

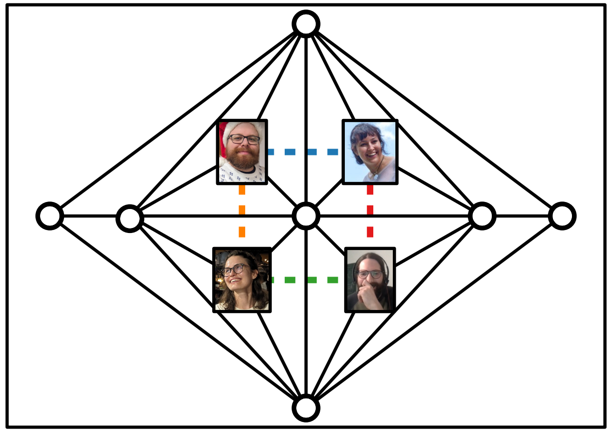
Conclusion



plane graph G

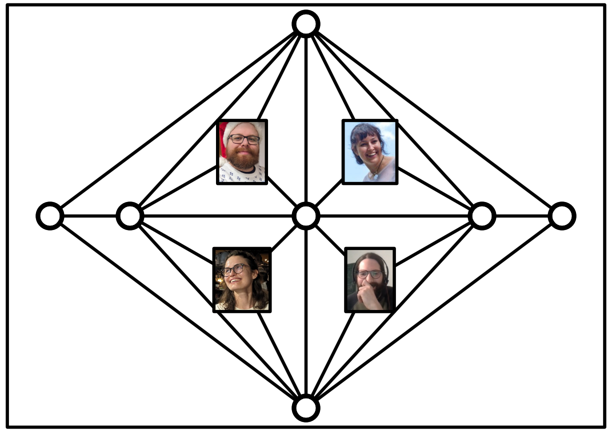


edges E' btw. vtcs in G



1-planar drawing of $G + E'$ that keeps the embedding of G

Conclusion



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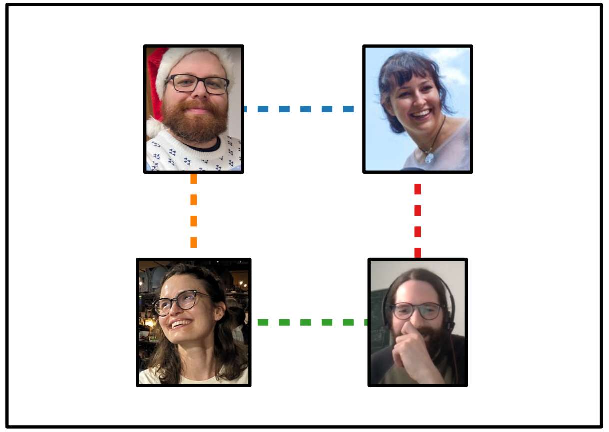

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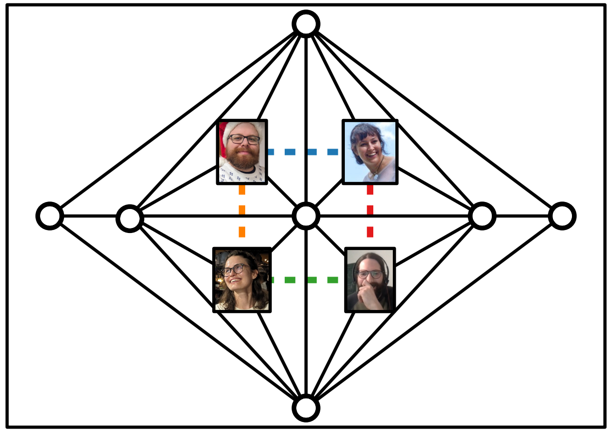
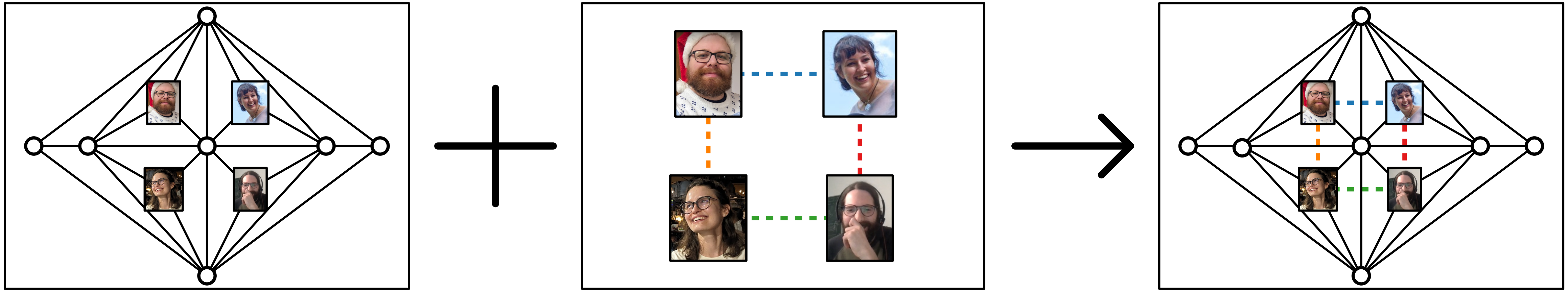


1-planar drawing of $G + E'$ that keeps the embedding of G

- G triangulated ➡ $\mathcal{O}(n)$ time

Conclusion



plane graph G | edges E' btw. vtcs in G | **1-planar** drawing of $G + E'$ that keeps the embedding of G

- G triangulated ➡ $\mathcal{O}(n)$ time
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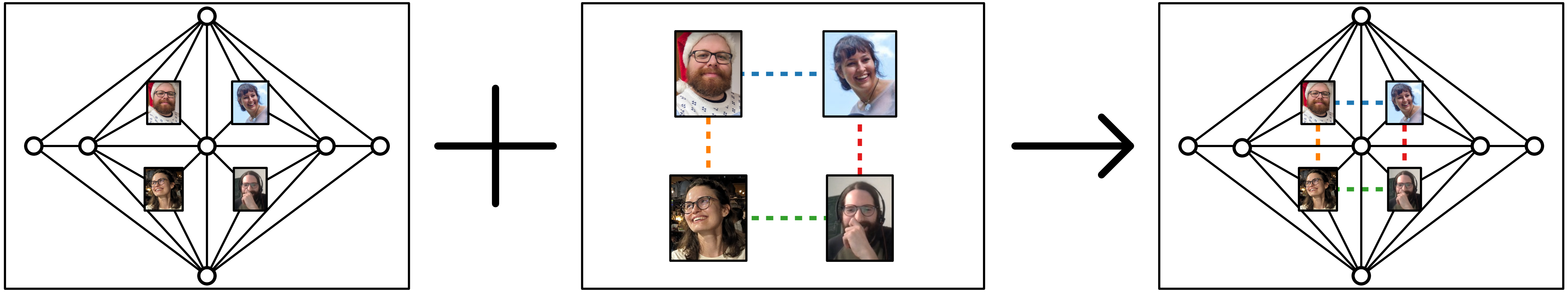
Conclusion



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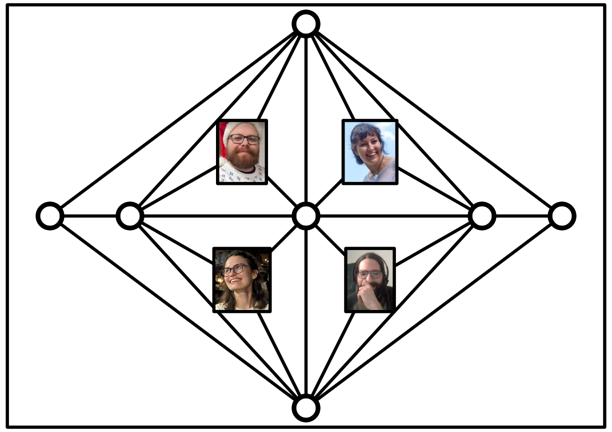
Open Problems

Conclusion



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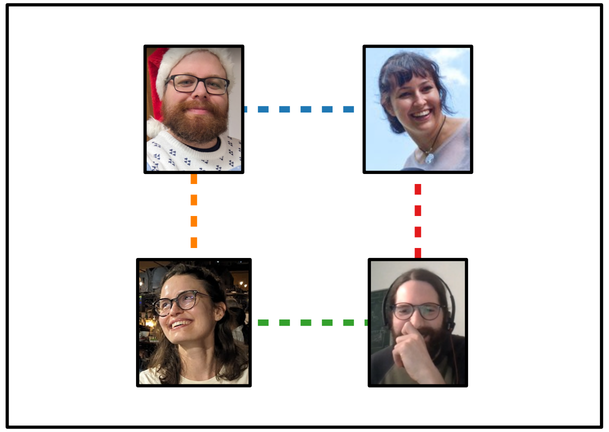

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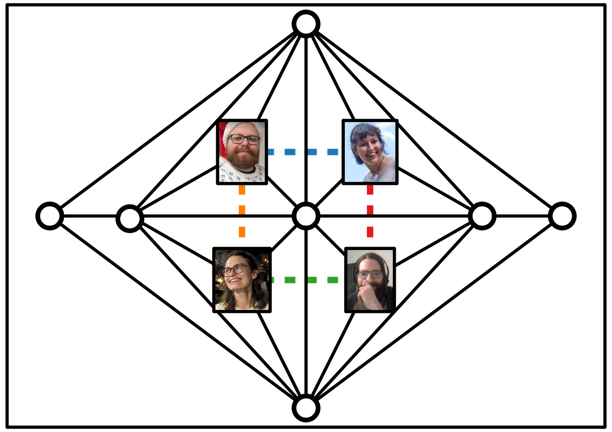

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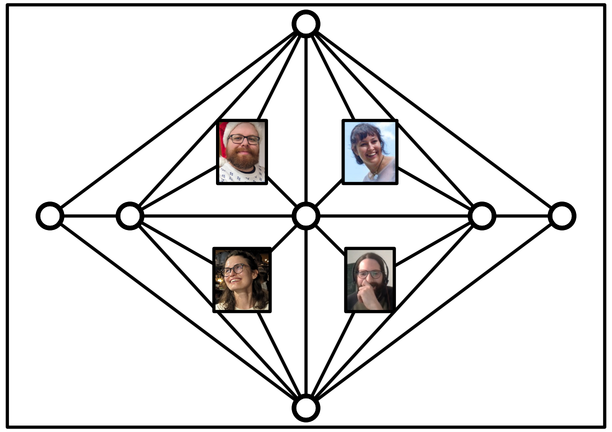
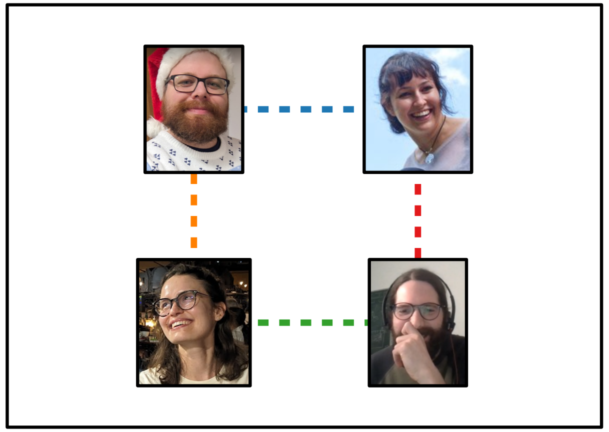
Conclusion



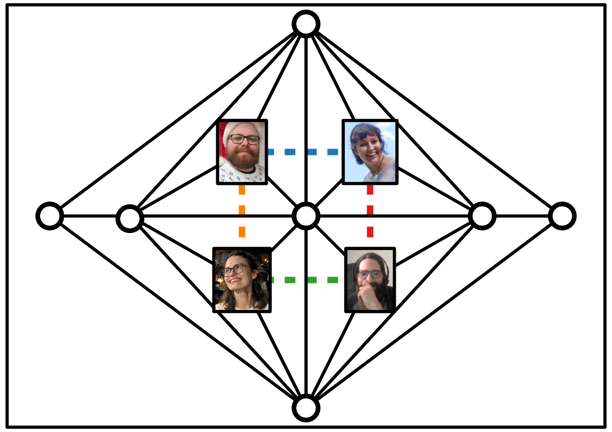
plane graph G

+



edges E' btw. vtcs in G

→



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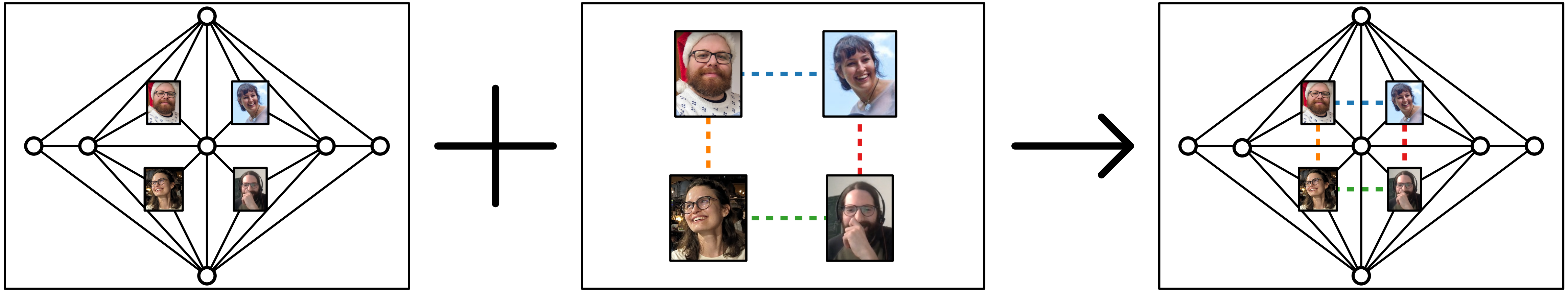

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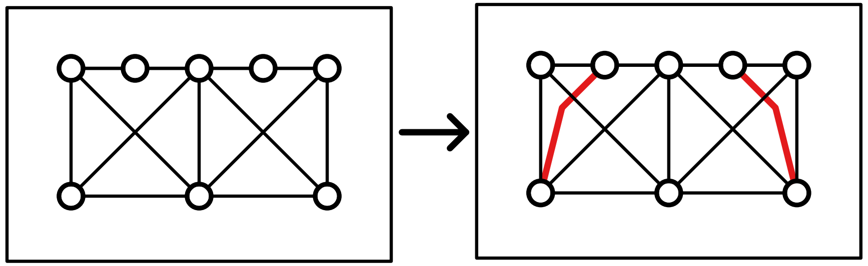
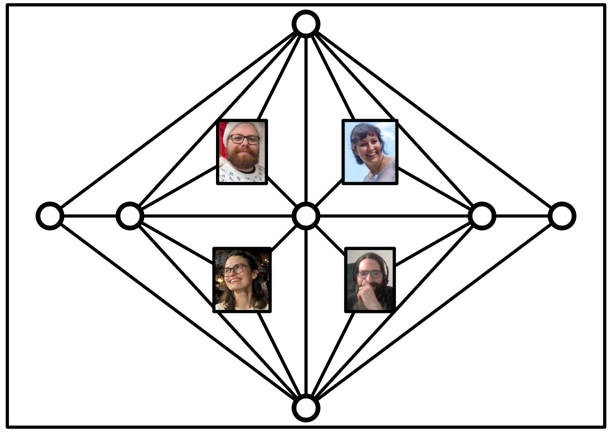
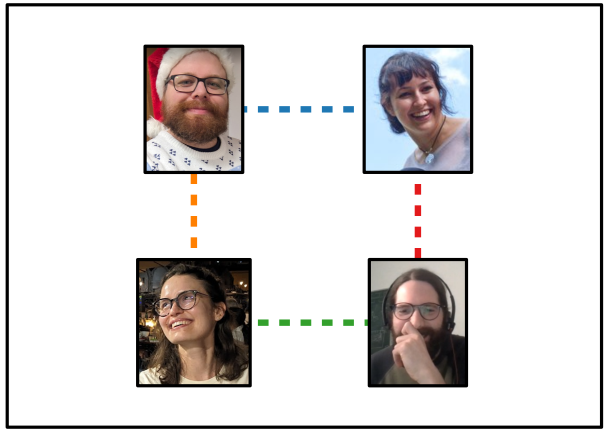


RAC 1-bend RAC

Conclusion



plane graph G



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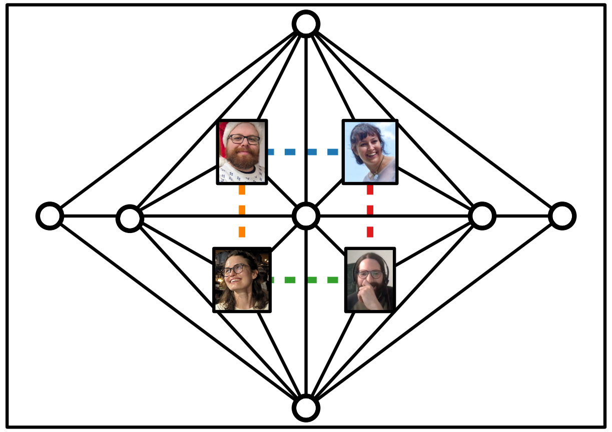

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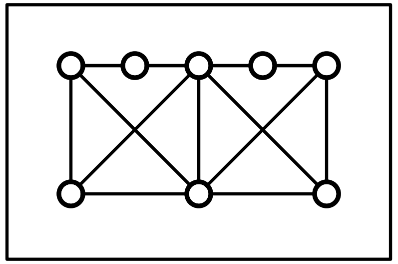
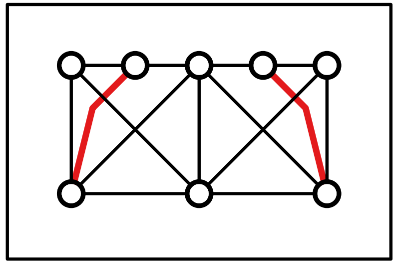
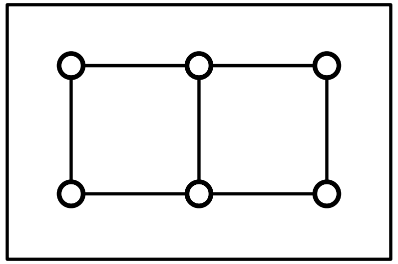


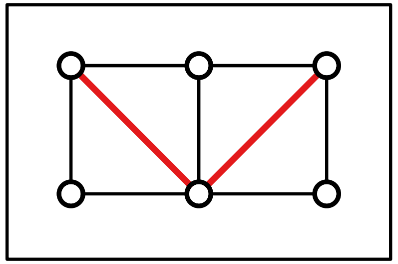
RAC



1-bend RAC

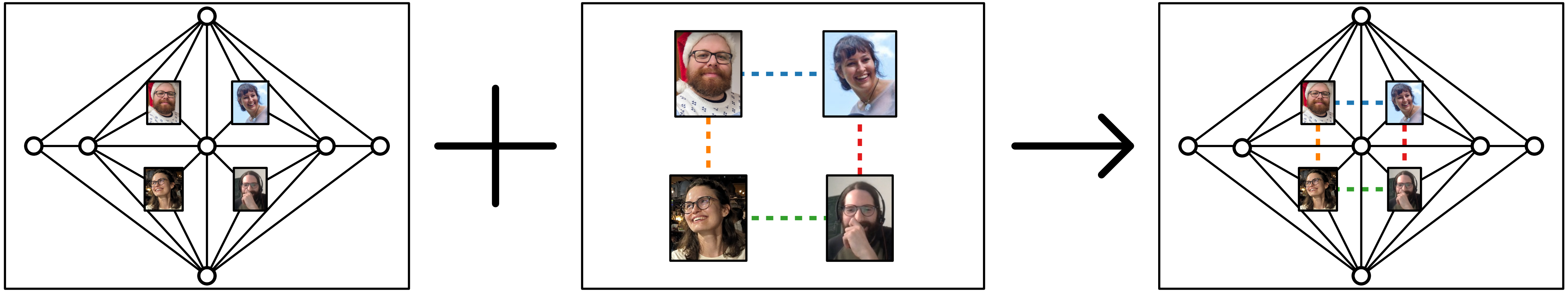


orthogonal



octilinear

Conclusion



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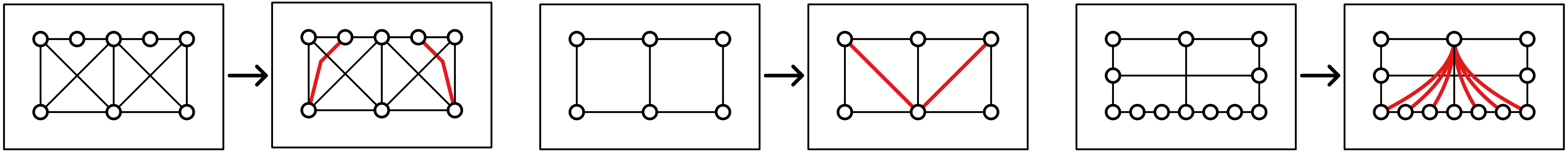

RAC → 1-bend RAC; orthogonal → octilinear; 1-planar → fan-planar